Language Reference
Compatible functions ........................................... 194
Multiple function declarations ............................. 195
Function storage class specifiers ........................... 196
The static storage class specifier .......................... 196
The extern storage class specifier ........................ 196
Function specifiers ............................................. 196
The inline function specifier ................................ 197
Function return type specifiers ............................. 201
Function return values ....................................... 202
Function declarators .......................................... 202
Parameter declarations ....................................... 203
Function attributes .......................................... 206
The alias function attribute ................................ 207
The always_inline function attribute ..................... 208
The const function attribute ................................ 208
The constructor and destructor function attributes ..... 209
The format function attribute ............................... 209
The format_arg function attribute ........................ 210
The noninline function attribute ........................... 210
The noreturn function attribute ............................ 211
The pure function attribute ................................ 211
The section function attribute ............................. 211
The weak function attribute ................................ 212
The main() function .......................................... 212
Function calls .................................................. 213
Pass by value .................................................. 214
Pass by reference .............................................. 215
Default arguments in C++ functions ...................... 216
Restrictions on default arguments ........................... 217
Evaluation of default arguments ............................ 218
Pointers to functions ......................................... 219
Nested functions .............................................. 219

Chapter 9. Namespaces (C++ only) ............................ 221
Defining namespaces ......................................... 221
Declaring namespaces ....................................... 221
Creating a namespace alias .................................. 221
Creating an alias for a nested namespace ................. 222
Extending namespaces ....................................... 222
Namespaces and overloading ................................. 223
Unnamed namespaces ....................................... 223
Namespace member definitions .............................. 225
Namespaces and friends ...................................... 225
The using directive ........................................... 226
The using declaration and namespaces ..................... 226
Explicit access .................................................. 227

Chapter 10. Overloading (C++ only) ............................ 229
Overloading functions ....................................... 229
Restrictions on overloaded functions ....................... 230
Overloading operators ....................................... 231
Overloading unary operators ................................ 234
Overloading increment and decrement operators ......... 235
Overloading binary operators ................................ 236
Overloading assignments ..................................... 237
Overloading function calls ................................... 238
Overloading subscripting ..................................... 239
Overloading class member access ............................ 240
Overload resolution .......................................... 241
Implicit conversion sequences ............................... 242
Resolving addresses of overloaded functions ............... 243

Chapter 11. Classes (C++ only) ................................. 245
Declaring class types ......................................... 246
Using class objects ........................................... 246
Classes and structures ....................................... 248
Scope of class names ......................................... 249
Incomplete class declarations ............................... 250
Nested classes .................................................. 250
Local classes ................................................... 252
Local type names .............................................. 253

Chapter 12. Class members and friends (C++ only) ............ 255
Class member lists ............................................ 255
Data members ................................................... 256
Member functions ............................................. 257
  Inline member functions .................................. 257
  Constant and volatile member functions ................. 258
  Virtual member functions ................................ 258
  Special member functions ................................ 258
Member scope .................................................. 259
Pointers to members ......................................... 260
The this pointer .............................................. 261
Static members ................................................. 265
  Using the class access operators with static members .... 265
  Static data members ....................................... 266
  Static member functions .................................. 268
Member access ................................................. 269
Friends .......................................................... 272
  Friend scope .............................................. 274
  Friend access ............................................ 276

Chapter 13. Inheritance (C++ only) ............................. 277
Derivation ...................................................... 279
Inherited member access ..................................... 282
Protected members .......................................... 282
Access control of base class members ...................... 283
The using declaration and class members ................. 284
  Overloading member functions from base and derived classes .............. 285
  Changing the access of a class member ........................................ 287
Multiple inheritance ......................................... 288
  Virtual base classes ...................................... 289
  Multiple access .......................................... 290
  Ambiguous base classes .................................. 291
Virtual functions ............................................. 295
  Ambiguous virtual function calls .................................. 298
  Virtual function access .................................... 300
Abstract classes .............................................. 300

Chapter 14. Special member functions
  (C++ only) .................................................. 303
Overview of constructors and destructors .................. 303
Constructors ................................................... 305
  Default constructors ....................................... 305
Explicit initialization with constructors .................... 306
Initializing base classes and members ...................... 308
### Chapter 15. Templates (C++ only) 329

- Template parameters 330
  - Type template parameters 330
  - Non-type template parameters 330
  - Template template parameters 331
- Default arguments for template parameters 331
- Template arguments 332
  - Template type arguments 332
  - Template non-type arguments 333
  - Template template arguments 335
- Class templates 336
  - Class template declarations and definitions 337
  - Static data members and templates 338
  - Member functions of class templates 338
  - Friends and templates 339
- Function templates 340
  - Template argument deduction 341
  - Overloading function templates 346
  - Partial ordering of function templates 347
- Template instantiation 348
  - Implicit instantiation 348
  - Explicit instantiation 350
- Template specialization 351
  - Explicit specialization 351
  - Partial specialization 356
- Name binding and dependent names 358
- The typename keyword 359
- The template keyword as qualifier 360

### Chapter 16. Exception handling (C++ only) 363

- try blocks 363
  - Nested try blocks 364
- catch blocks 365
  - Function try block handlers 366
  - Arguments of catch blocks 369
  - Matching between exceptions thrown and caught 369
  - Order of catching 370
- throw expressions 371
  - Rethrowing an exception 371
- Stack unwinding 373
- Exception specifications 374
- Special exception handling functions 377
  - The unexpected() function 377
  - The terminate() function 378
- The set_unexpected() and set_terminate() functions 379

### Appendix A. The IBM XL C language extensions 403

- C99 features as extensions to C89 403
- Extensions for Unicode support 405
- Extensions for GNU C compatibility 405
- Extensions for VMX support 407

### Appendix B. The IBM XL C++ language extensions 409

- General IBM extensions 409
- Extensions for C99 compatibility 410
- Extensions for Unicode support 410
- Extensions for GNU C compatibility 412
- Extensions for GNU C++ compatibility 413
- Extensions for VMX support 413

### Appendix C. Vector data types and literals 415

### Index 417

### Notices 425

Programming interface information 426
Trademarks and service marks 427
Industry standards 427
About this document

This document describes the syntax, semantics, and IBM® implementation of the C and C++ programming languages. Although the XL C and XL C++ compilers conform to the specifications maintained by the ISO standards for the C and C++ programming languages, the compilers also incorporate many extensions to the core languages. These extensions have been implemented with the aims of enhancing usability in specific operating environments, assuring compatibility with other compilers, and supporting new hardware capabilities. For example, on UNIX® platforms, many language constructs have been added for compatibility with the GNU Compiler Collection (GCC), to maximize portability between the two development environments.

Who should read this document

This document is a reference for users who already have experience programming applications in C or C++. Users new to C or C++ can still use this document to find information on the language and features unique to XL C/C++; however, this reference does not aim to teach programming concepts nor to promote specific programming practices.

Supported language standards

The C and C++ languages described in this reference are based on the following standards:

- The C language described in Information Technology – Programming languages – C, ISO/IEC 9899:1990, henceforth referred to as C89. This was the first ISO C standard.
- The C language described in Information Technology – Programming languages – C, ISO/IEC 9899:1999, henceforth referred to as C99. This is an update to the C89 standard.
- The C++ language described in Information Technology – Programming languages – C++, ISO/IEC 14882:1998, the first formal definition of the language, henceforth referred to as C++98.
- The C++ language described in Information Technology – Programming languages – C++, ISO/IEC 14882:2003(E), which is the current meaning of the term Standard C++.

The XL C compiler supports all of the language features specified in the C99 standard, and the C language described in this reference is consistent with the C99 standard. Note that the standard also specifies features in the runtime library. These features may not be supported in the current runtime library and operating environment.

The XL C++ compiler supports all of the language features specified in Standard C++, and the C++ language described in this reference is consistent with Standard C++.

This document uses the term K&R C to refer to the C language plus the generally accepted extensions produced by Brian Kernighan and Dennis Ritchie that were in use prior to the ISO standardization of C.
Other standards and specifications

XL C/C++ is designed to support the following additional standards and specifications. You can refer to these standards for precise definitions of some of the features found in this document.

- **Information technology — Programming Languages — Universal Multiple-Octet Coded Character Set (UCS), ISO/IEC 10646:2003**

- **The Unicode Standard, Version 4.0**, the Unicode Consortium. The online version is available at [www.unicode.org](http://www.unicode.org)

- **Information Technology — Programming Languages — Extensions for the programming language C to support new character data types, ISO/IEC DTR 19769**. This draft technical report has been accepted by the C standards committee, and is available at [http://www.open-std.org/JTC1/SC22/WG14/www/docs/n1040.pdf](http://www.open-std.org/JTC1/SC22/WG14/www/docs/n1040.pdf)


Language levels and language extensions

We refer to the following language specifications as “base language levels” in order to introduce the notion of an extension to a base. In this context the base language levels refer to the following specifications:

- Standard C++
- C++98
- C99
- C89

In addition to the features supported by the base levels, XL C/C++ contains language extensions that enhance usability and facilitate porting programs to different platforms, including:

- Extensions to C++ to support C99 standard features
- Extensions related to GNU C and C++
- Extensions supporting the AltiVec Programming Interface

You can control the language level to be used for compilation through several mechanisms, including:

- various invocation commands
- the `-qlanglvl` option
- the `#pragma langlvl` directive

With a few exceptions, almost all of the language extensions are supported when you compile using the basic invocation commands `xlc` (for C) and `xlC` or `xlc++` (for C++).

The default language level for the `xlc` invocation command is **extc89**, which includes most of the IBM extensions described in this document, as well as most of the features introduced by the C99 standard. In addition, in the default configuration file which is shipped with the compiler, the stanza for the `xlc` command enables some compiler options that control C99 features not included in **extc89**. For a complete listing of the C extensions and methods for enabling them, see Appendix A, “The IBM XL C language extensions,” on page 403.
The default language level for the xC and xlc++ invocation commands is **extended**, which includes most of the IBM extensions described in this document, as well as many C99 features. For a complete listing of the C++ extensions and methods for enabling them, see Appendix B, “The IBM XL C++ language extensions,” on page 409.

For information on the various methods for controlling the language level for compilation, see “How to choose a compiler invocation” [qlanglvl] and #pragma [langlvl] in the **XL C/C++ Compiler Reference**.

### Extensions to C++ to support C99 standard features

The Standard C++ language specification does not include many of the features specified in the C99 language standard. To promote compatibility and portability between C99 and C++, the XL C++ compiler enables many of the C99 features that are supported by the XL C compiler. Since these features extend Standard C++, they are considered extensions to the base language. In this reference, unless the text is marked to indicate that a feature is supported in C or C99 only, C99 features also apply to C++. A complete list of C99 features supported in XL C++ is also provided in Appendix B, “The IBM XL C++ language extensions,” on page 409.

### Extensions related to GNU C and C++

Certain language extensions that correspond to GNU C and C++ features are implemented to facilitate portability. These include extensions to C89, C99, C++98, and Standard C++. Throughout this document, the text indicates the IBM extensions that have been implemented for compatibility with GNU C and C++; a complete list of these is also provided in Appendix A, “The IBM XL C language extensions,” on page 403 and Appendix B, “The IBM XL C++ language extensions,” on page 409.

### Extensions supporting the Altivec Programming Interface

XL C/C++ supports Altivec vector types when [VMX support is enabled]. These language extensions exploit the SIMD and parallel processing capabilities of the PowerPC® processor, and facilitate the associated optimization techniques. The IBM implementation of the Altivec Programming Interface specification is an extended implementation, which, for the most part, matches the syntax and semantics of the GNU C implementation. A list of the IBM extensions to the Altivec Programming Interface is also provided in Appendix A, “The IBM XL C language extensions,” on page 403 and Appendix B, “The IBM XL C++ language extensions,” on page 409. A complete list of vector data types and literals is provided in Appendix C, “Vector data types and literals,” on page 415.

### How to use this document

Unless indicated otherwise, all of the text in this reference pertains to both C and C++ languages. Where there are differences between languages, these are indicated through qualifying text and other graphical elements (see “Qualifying elements (icons and bracket separators)” on page xi) for the conventions used.

While this document covers both standard and implementation-specific features, it does not include the following topics:

- Standard C and C++ library functions and headers. For information on the standard C and C++ libraries, refer to your Linux operating system documentation.
Constructs for writing multi-threaded programs, including OpenMP directives and functions and POSIX Pthread functions. For reference documentation on OpenMP constructs, see the IBM XL C/C++ Advanced Edition V8.0 for Linux Compiler Reference for documentation on Pthreads library functions, refer to your Linux documentation.

Compiler pragmas, predefined macros, and built-in functions. These are described in the IBM XL C/C++ Advanced Edition V8.0 for Linux Compiler Reference.

How this document is organized

This document is organized to loosely follow the structure of the ISO standard language specifications and topics are grouped into similar headings.

Chapters 1 through 8 discuss language elements that are common to both C and C++, including lexical elements, data types, declarations, declarators, type conversions, expressions, operators, statements, and functions. Throughout these chapters, both standard features and extensions are discussed. Chapters 9 through 16 discuss standard C++ features exclusively, including classes, overloading, inheritance, templates, and exception handling. Chapter 17 discusses directives to the preprocessor.

The appendices provide summary lists of all the extended features supported by each language.

Conventions used in this document

The following sections discuss the conventions used in this document:

- “Typographical conventions”
- “Qualifying elements (icons and bracket separators)” on page xi
- “How to read syntax diagrams” on page xi
- “Examples” on page xiii

Typographical conventions

The following table explains the typographical conventions used in this document.

<table>
<thead>
<tr>
<th>Typeface</th>
<th>Indicates</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>bold</strong></td>
<td>Commands, executable names, compiler options and pragma directives.</td>
<td>You can also use the <code>#pragma comment</code> directive to place comments into an object module.</td>
</tr>
<tr>
<td><em>italics</em></td>
<td>Parameters or variables whose actual names or values are to be supplied by the user. <em>Italics</em> are also used to introduce new terms.</td>
<td>The <em>attribute name</em> can be specified with or without leading and trailing double underscore characters.</td>
</tr>
<tr>
<td>monospace</td>
<td>Programming keywords and library functions, compiler built-in functions, file and directory names, examples of program code, command strings, or user-defined names.</td>
<td>Case and default label statements only appear in <code>switch</code> statements.</td>
</tr>
</tbody>
</table>
Qualifying elements (icons and bracket separators)

This document uses marked bracket separators to delineate large blocks of text and icons to delineate small segments of text as follows:

<table>
<thead>
<tr>
<th>Bracket separator text</th>
<th>Icon</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>C only</td>
<td>![C icon]</td>
<td>The text describes a feature that is supported in the C language only; or describes behavior that is specific to the C language.</td>
</tr>
<tr>
<td>C++ only</td>
<td>![C++]</td>
<td>The text describes a feature that is supported in the C++ language only; or describes behavior that is specific to the C++ language.</td>
</tr>
<tr>
<td>IBM extension</td>
<td>![IBM]</td>
<td>The text describes a feature that is an IBM XL C/C++ compiler extension to the standard language specifications.</td>
</tr>
</tbody>
</table>

How to read syntax diagrams

- Read the syntax diagrams from left to right, from top to bottom, following the path of the line.
  - The symbol indicates the beginning of a command, directive, or statement.
  - The symbol indicates that the command, directive, or statement syntax is continued on the next line.
  - The symbol indicates that a command, directive, or statement is continued from the previous line.
  - The symbol indicates the end of a command, directive, or statement.
  - Diagrams of syntactical units other than complete commands, directives, or statements start with the symbol and end with the symbol.
- Required items appear on the horizontal line (the main path).

- Optional items are shown below the main path.

- If you can choose from two or more items, they are shown vertically, in a stack.
  - If you must choose one of the items, one item of the stack is shown on the main path.
  - If choosing one of the items is optional, the entire stack is shown below the main path.

The item that is the default is shown above the main path.
• An arrow returning to the left above the main line indicates an item that can be repeated.

A repeat arrow above a stack indicates that you can make more than one choice from the stacked items, or repeat a single choice.

• Keywords are shown in nonitalic letters and should be entered exactly as shown (for example, extern). Variables are shown in italicized lowercase letters (for example, identifier). They represent user-supplied names or values.

• If punctuation marks, parentheses, arithmetic operators, or other such symbols are shown, you must enter them as part of the syntax.

The following syntax diagram example shows the syntax for the **#pragma comment** directive.

1. This is the start of the syntax diagram.
2. The symbol # must appear first.
3. The keyword pragma must appear following the # symbol.
4. The name of the pragma comment must appear following the keyword pragma.
5. An opening parenthesis must be present.
6. The comment type must be entered only as one of the types indicated: compiler, date, timestamp, copyright, or user.
7. A comma must appear between the comment type copyright or user, and an optional character string.
8. A character string must follow the comma. The character string must be enclosed in double quotation marks.
9. A closing parenthesis is required.
10. This is the end of the syntax diagram.

The following examples of the **#pragma comment** directive are syntactically correct according to the diagram shown above:

---

**XL C/C++ Language Reference**
Examples

The examples in this document, except where otherwise noted, are coded in a simple style that does not try to conserve storage, check for errors, achieve fast performance, or demonstrate all possible methods to achieve a specific result.

Related information

IBM XL C/C++ publications

XL XL C/C++ provides product documentation in the following formats:

- Readme files

Readme files contain late-breaking information, including changes and corrections to the product documentation. Readme files are located by default in the /opt/ibmcmp/vacpp/8.0/ directory and in the root directory of the installation CD.

- Installable man pages

Man pages are provided for the compiler invocations and all command-line utilities provided with the product. Instructions for installing and accessing the man pages are provided in the IBM XL C/C++ Advanced Edition V8.0 for Linux Installation Guide.

- Information center

The information center of searchable HTML files can be launched on a network and accessed remotely or locally. Instructions for installing and accessing the information center are provided in the IBM XL C/C++ Advanced Edition V8.0 for Linux Installation Guide. The information center is also viewable on the Web at: http://publib.boulder.ibm.com/infocenter/lnxpcomp/index.jsp.

- PDF documents

PDF documents are located by default in the /opt/ibmcmp/vacpp/8.0/doc/language/pdf/ directory, where language is one of en_US, ja_JP, or zh_CN. The PDFs are also available on the Web at: www.ibm.com/software/awdtools/xlcpp/library.

In addition to this document, the following files comprise the full set of XL C/C++ product manuals:

Table 3. XL C/C++ PDF files

<table>
<thead>
<tr>
<th>Document title</th>
<th>PDF file name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>IBM XL C/C++ Advanced Edition V8.0 for Linux</td>
<td>install.pdf</td>
<td>Contains information for installing XL C/C++ and configuring your environment for basic compilation and program execution.</td>
</tr>
<tr>
<td>Getting Started Guide, SC09-8015-00</td>
<td>getstart.pdf</td>
<td>Contains an introduction to the XL C/C++ product, with information on setting up and configuring your environment, compiling and linking programs, and troubleshooting compilation errors.</td>
</tr>
</tbody>
</table>
Table 3. XL C/C++ PDF files (continued)

<table>
<thead>
<tr>
<th>Document title</th>
<th>PDF file name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>IBM XL C/C++ Advanced Edition V8.0 for Linux Compiler Reference, SC09-8013-00</td>
<td>compiler.pdf</td>
<td>Contains information about the various compiler options, pragmas, macros, environment variables, and built-in functions, including those used for parallel processing.</td>
</tr>
<tr>
<td>IBM XL C/C++ Advanced Edition V8.0 for Linux Programming Guide, SC09-8014-00</td>
<td>proguide.pdf</td>
<td>Contains information on advanced programming topics, such as application porting, interlanguage calls with Fortran code, library development, application optimization and parallelization, and the XL C/C++ high-performance libraries.</td>
</tr>
</tbody>
</table>

These PDF files are viewable and printable from Adobe Reader. If you do not have the Adobe Reader installed, you can download it from www.adobe.com.

Additional documentation
More documentation related to XL C/C++, including redbooks, whitepapers, tutorials, and other articles, is available on the Web at:

www.ibm.com/software/awdtools/xlcpp/library

Related publications
You might want to consult the following publications, which are also referenced throughout this document:

- Using the GNU Compiler Collection (GCC), available at http://gcc.gnu.org/onlinedocs/

Technical support
Additional technical support is available from the XL C/C++ Support page. This page provides a portal with search capabilities to a large selection of technical support FAQs and other support documents. You can find the XL C/C++ Support page on the Web at:

www.ibm.com/software/awdtools/xlcpp/support

If you cannot find what you need, you can e-mail:

compinfo@ca.ibm.com

For the latest information about XL C/C++, visit the product information site at:

www.ibm.com/software/awdtools/xlcpp

How to send your comments
Your feedback is important in helping to provide accurate and high-quality information. If you have any comments about this document or any other XL C/C++ documentation, send your comments by e-mail to:

compinfo@ca.ibm.com
Be sure to include the name of the document, the part number of the document, the version of XL C/C++, and, if applicable, the specific location of the text you are commenting on (for example, a page number or table number).
Chapter 1. Scope and linkage

Scope is the largest region of program text in which a name can potentially be used without qualification to refer to an entity; that is, the largest region in which the name potentially is valid. Broadly speaking, scope is the general context used to differentiate the meanings of entity names. The rules for scope combined with those for name resolution enable the compiler to determine whether a reference to an identifier is legal at a given point in a file.

The scope of a declaration and the visibility of an identifier are related but distinct concepts. Scope is the mechanism by which it is possible to limit the visibility of declarations in a program. The visibility of an identifier is the region of program text from which the object associated with the identifier can be legally accessed. Scope can exceed visibility, but visibility cannot exceed scope. Scope exceeds visibility when a duplicate identifier is used in an inner declarative region, thereby hiding the object declared in the outer declarative region. The original identifier cannot be used to access the first object until the scope of the duplicate identifier (the lifetime of the second object) has ended.

Thus, the scope of an identifier is interrelated with the storage duration of the identified object, which is the length of time that an object remains in an identified region of storage. The lifetime of the object is influenced by its storage duration, which in turn is affected by the scope of the object identifier.

Linkage refers to the use or availability of a name across multiple translation units or within a single translation unit. The term translation unit refers to a source code file plus all the header and other source files that are included after preprocessing with the #include directive, minus any source lines skipped because of conditional preprocessing directives. Linkage allows the correct association of each instance of an identifier with one particular object or function.

Scope and linkage are distinguishable in that scope is for the benefit of the compiler, whereas linkage is for the benefit of the linker. During the translation of a source file to object code, the compiler keeps track of the identifiers that have external linkage and eventually stores them in a table within the object file. The linker is thereby able to determine which names have external linkage, but is unaware of those with internal or no linkage.

The distinctions between the different types of scopes are discussed in “Scope” on page 2. The different types of linkages are discussed in “Program linkage” on page 8.

Related information

- “Storage class specifiers” on page 43
- Chapter 9, “Namespaces (C++ only),” on page 221
Scope

The scope of an identifier is the largest region of the program text in which the identifier can potentially be used to refer to its object. In C++, the object being referred to must be unique. However, the name to access the object, the identifier itself, can be reused. The meaning of the identifier depends upon the context in which the identifier is used. Scope is the general context used to distinguish the meanings of names.

The scope of an identifier is possibly noncontiguous. One of the ways that breakage occurs is when the same name is reused to declare a different entity, thereby creating a contained declarative region (inner) and a containing declarative region (outer). Thus, point of declaration is a factor affecting scope. Exploiting the possibility of a noncontiguous scope is the basis for the technique called information hiding.

The concept of scope that exists in C was expanded and refined in C++. The following table shows the kinds of scopes and the minor differences in terminology.

<table>
<thead>
<tr>
<th>C</th>
<th>C++</th>
</tr>
</thead>
<tbody>
<tr>
<td>block</td>
<td>local</td>
</tr>
<tr>
<td>function</td>
<td>function</td>
</tr>
<tr>
<td>Function prototype</td>
<td>Function prototype</td>
</tr>
<tr>
<td>file</td>
<td>global namespace</td>
</tr>
<tr>
<td></td>
<td>namespace</td>
</tr>
<tr>
<td></td>
<td>class</td>
</tr>
</tbody>
</table>

In all declarations, the identifier is in scope before the initializer. The following example demonstrates this:

```cpp
int x;
void f() {
    int x = x;
}
```

The x declared in function f() has local scope, not global scope.

Related information
- [Chapter 9, “Namespaces (C++ only),”](#) on page 221

Block/local scope

A name has local scope or block scope if it is declared in a block. A name with local scope can be used in that block and in blocks enclosed within that block, but the name must be declared before it is used. When the block is exited, the names declared in the block are no longer available.

Parameter names for a function have the scope of the outermost block of that function. Also, if the function is declared and not defined, these parameter names have function prototype scope.
When one block is nested inside another, the variables from the outer block are usually visible in the nested block. However, if the declaration of a variable in a nested block has the same name as a variable that is declared in an enclosing block, the declaration in the nested block hides the variable that was declared in the enclosing block. The original declaration is restored when program control returns to the outer block. This is called block visibility.

Name resolution in a local scope begins in the immediate scope in which the name is used and continues outward with each enclosing scope. The order in which scopes are searched during name resolution causes the phenomenon of information hiding. A declaration in an enclosing scope is hidden by a declaration of the same identifier in a nested scope.

Related information
• “Block statements” on page 170

Function scope
The only type of identifier with function scope is a label name. A label is implicitly declared by its appearance in the program text and is visible throughout the function that declares it.

A label can be used in a goto statement before the actual label is seen.

Related information
• “Labeled statements” on page 167

Function prototype scope
In a function declaration (also called a function prototype) or in any function declarator—except the declarator of a function definition—parameter names have function prototype scope. Function prototype scope terminates at the end of the nearest enclosing function declarator.

Related information
• “Function declarations” on page 192

File/global scope

C only

A name has file scope if the identifier’s declaration appears outside of any block. A name with file scope and internal linkage is visible from the point where it is declared to the end of the translation unit.

End of C only

C++ only

Global scope or global namespace scope is the outermost namespace scope of a program, in which objects, functions, types and templates can be defined. A name has global namespace scope if the identifier’s declaration appears outside of all blocks, namespaces, and classes.

A name with global namespace scope and internal linkage is visible from the point where it is declared to the end of the translation unit.
A name with global (namespace) scope is also accessible for the initialization of global variables. If that name is declared extern, it is also visible at link time in all object files being linked.

A user-defined namespace can be nested within the global scope using namespace definitions, and each user-defined namespace is a different scope, distinct from the global scope.

Related information
- Chapter 9, “Namespaces (C++ only),” on page 221
- “Internal linkage” on page 8

Examples of scope in C

The following example declares the variable x on line 1, which is different from the x it declares on line 2. The declared variable on line 2 has function prototype scope and is visible only up to the closing parenthesis of the prototype declaration. The variable x declared on line 1 resumes visibility after the end of the prototype declaration.

```c
1 int x = 4; /* variable x defined with file scope */
2 long myfunc(int x, long y); /* variable x has function */
3     /* prototype scope */
4 int main(void)
5 {
6     /* ... */
7 }
```

The following program illustrates blocks, nesting, and scope. The example shows two kinds of scope: file and block. The main function prints the values 1, 2, 3, 0, 3, 2, 1 on separate lines. Each instance of i represents a different variable.
Class scope (C++ only)

A name declared within a member function hides a declaration of the same name whose scope extends to or past the end of the member function’s class.

When the scope of a declaration extends to or past the end of a class definition, the regions defined by the member definitions of that class are included in the scope of the class. Members defined lexically outside of the class are also in this scope. In addition, the scope of the declaration includes any portion of the declarator following the identifier in the member definitions.

The name of a class member has class scope and can only be used in the following cases:

- In a member function of that class
- In a member function of a class derived from that class
- After the . (dot) operator applied to an instance of that class
- After the . (dot) operator applied to an instance of a class derived from that class, as long as the derived class does not hide the name
- After the -> (arrow) operator applied to a pointer to an instance of that class
- After the -> (arrow) operator applied to a pointer to an instance of a class derived from that class, as long as the derived class does not hide the name
- After the :: (scope resolution) operator applied to the name of a class
- After the :: (scope resolution) operator applied to a class derived from that class

Related information

- Chapter 11, “Classes (C++ only),” on page 245
- “Scope of class names” on page 249
Namespaces of identifiers

Namespaces are the various syntactic contexts within which an identifier can be used. Within the same context and the same scope, an identifier must uniquely identify an entity. Note that the term namespace as used here applies to C as well as C++ and does not refer to the C++ namespace language feature. The compiler sets up namespaces to distinguish among identifiers referring to different kinds of entities. Identical identifiers in different namespaces do not interfere with each other, even if they are in the same scope.

The same identifier can declare different objects as long as each identifier is unique within its namespace. The syntactic context of an identifier within a program lets the compiler resolve its namespace without ambiguity.

Within each of the following four namespaces, the identifiers must be unique.

- Tags of these types must be unique within a single scope:
  - Enumerations
  - Structures and unions
- Members of structures, unions, and classes must be unique within a single structure, union, or class type.
- Statement labels have function scope and must be unique within a function.
- All other ordinary identifiers must be unique within a single scope:
  - C function names (C++ function names can be overloaded)
  - Variable names
  - Names of function parameters
  - Enumeration constants
  - typedef names.

You can redefine identifiers in the same namespace but within enclosed program blocks.

Structure tags, structure members, variable names, and statement labels are in four different namespaces. No name conflict occurs among the items named student in the following example:

```c
int get_item()
{
    struct student  /* structure tag */
    {
        char name[20]; /* this structure member may not be named student */
        int section;
        int id;
    } sam;  /* this structure variable should not be named student */

goto student;
student;:  /* null statement label */
return 0;

student fred;  /* legal struct declaration in C++ */
}```
The compiler interprets each occurrence of student by its context in the program. For example, when student appears after the keyword struct, it is a structure tag. The name student may not be used for a structure member of struct student. When student appears after the goto statement, the compiler passes control to the null statement label. In other contexts, the identifier student refers to the structure variable.

Name hiding (C++ only)

If a class name or enumeration name is in scope and not hidden, it is visible. A class name or enumeration name can be hidden by an explicit declaration of that same name — as an object, function, or enumerator — in a nested declarative region or derived class. The class name or enumeration name is hidden wherever the object, function, or enumerator name is visible. This process is referred to as name hiding.

In a member function definition, the declaration of a local name hides the declaration of a member of the class with the same name. The declaration of a member in a derived class hides the declaration of a member of a base class of the same name.

Suppose a name x is a member of namespace A, and suppose that the members of namespace A are visible in a namespace B because of a using declaration. A declaration of an object named x in namespace B will hide A::x. The following example demonstrates this:

```cpp
#include <iostream>
#include <typeinfo>
using namespace std;

namespace A {
    char x;
};

namespace B {
    using namespace A;
    int x;
};

int main() {
    cout << typeid(B::x).name() << endl;
}
```

The following is the output of the above example:

```
int
```

The declaration of the integer x in namespace B hides the character x introduced by the using declaration.

Related information

- Chapter 11, “Classes (C++ only),” on page 245
- “Member functions” on page 257
- “Member scope” on page 259
- Chapter 9, “Namespaces (C++ only),” on page 221
Program linkage

Linkage determines whether identifiers that have identical names refer to the same object, function, or other entity, even if those identifiers appear in different translation units. The linkage of an identifier depends on how it was declared. There are three types of linkages:

- **Internal linkage**: identifiers can only be seen within a translation unit.
- **External linkage**: identifiers can be seen (and referred to) in other translation units.
- **No linkage**: identifiers can only be seen in the scope in which they are defined.

Linkage does not affect scoping, and normal name lookup considerations apply.

---

C++ only

You can also have linkage between C++ and non-C++ code fragments, which is called language linkage. Language linkage enables the close relationship between C++ and C by allowing C++ code to link with that written in C. All identifiers have a language linkage, which by default is C++. Language linkage must be consistent across translation units, and non-C++ language linkage implies that the identifier has external linkage.

---

End of C++ only

---

Related information

- “The static storage class specifier” on page 44
- “The extern storage class specifier” on page 46
- “Function storage class specifiers” on page 196
- “Type qualifiers” on page 68
- “Anonymous unions” on page 62

Internal linkage

The following kinds of identifiers have internal linkage:

- Objects, references, or functions explicitly declared **static**
- Objects or references declared in namespace scope (or global scope in C) with the specifiers **const** and neither explicitly declared **extern**, nor previously declared to have external linkage
- Data members of an anonymous union
- Function templates explicitly declared **static**
- Identifiers declared in the unnamed namespace

A function declared inside a block will usually have external linkage. An object declared inside a block will usually have external linkage if it is specified **extern**. If a variable that has static storage is defined outside a function, the variable has internal linkage and is available from the point where it is defined to the end of the current translation unit.

If the declaration of an identifier has the keyword **extern** and if a previous declaration of the identifier is visible at namespace or global scope, the identifier has the same linkage as the first declaration.
External linkage

C only

In global scope, identifiers for the following kinds of entities declared without the static storage class specifier have external linkage:

- An object
- A function

If an identifier in C is declared with the extern keyword and if a previous declaration of an object or function with the same identifier is visible, the identifier has the same linkage as the first declaration. For example, a variable or function that is first declared with the keyword static and later declared with the keyword extern has internal linkage. However, a variable or function that has no linkage and was later declared with a linkage specifier will have the linkage that was expressly specified.

C++ only

In namespace scope, the identifiers for the following kinds of entities have external linkage:

- A reference or an object that does not have internal linkage
- A function that does not have internal linkage
- A named class or enumeration
- An unnamed class or enumeration defined in a typedef declaration
- An enumerator of an enumeration that has external linkage
- A template, unless it is a function template with internal linkage
- A namespace, unless it is declared in an unnamed namespace

If the identifier for a class has external linkage, then, in the implementation of that class, the identifiers for the following will also have external linkage:

- A member function.
- A static data member.
- A class of class scope.
- An enumeration of class scope.

No linkage

The following kinds of identifiers have no linkage:

- Names that have neither external or internal linkage
- Names declared in local scopes (with exceptions like certain entities declared with the extern keyword)
- Identifiers that do not represent an object or a function, including labels, enumerators, typedef names that refer to entities with no linkage, type names, function parameters, and template names
You cannot use a name with no linkage to declare an entity with linkage. For example, you cannot use the name of a class or enumeration or a typedef name referring to an entity with no linkage to declare an entity with linkage. The following example demonstrates this:

```c
int main() {
    struct A {
    // extern A a1;
    typedef A myA;
    // extern myA a2;
}
```

The compiler will not allow the declaration of `a1` with external linkage. Class `A` has no linkage. The compiler will not allow the declaration of `a2` with external linkage. The typedef name `myA` has no linkage because `A` has no linkage.

**Language linkage (C++ only)**

Linkage between C++ and non-C++ code fragments is called *language linkage*. All function types, function names, and variable names have a language linkage, which by default is C++.

You can link C++ object modules to object modules produced using other source languages such as C by using a *linkage specification*.

**Linkage specification syntax**

```
extern string_literal declaration
```

The `string_literal` is used to specify the linkage associated with a particular function. String literals used in linkage specifications should be considered as case-sensitive. All platforms support the following values for `string_literal`:

"C++" Unless otherwise specified, objects and functions have this default linkage specification.

"C" Indicates linkage to a C procedure

Calling shared libraries that were written before C++ needed to be taken into account requires the `#include` directive to be within an `extern "C" {}` declaration.

```c
extern "C" {
    #include "shared.h"
}
```

The following example shows a C printing function that is called from C++.

```c
// in C++ program
extern "C" int displayfoo(const char *);
int main() {
    return displayfoo("hello");
}
```

```c
/* in C program   */
#include <stdio.h>
extern int displayfoo(const char * str) {
    while (*str) {
        putchar(*str);
    }
```

**Language linkage (C++ only)**

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The following example shows a C printing function that is called from C++.

```c
// in C++ program
extern "C" int displayfoo(const char *);
int main() {
    return displayfoo("hello");
}
```

```c
/* in C program   */
#include <stdio.h>
extern int displayfoo(const char * str) {
    while (*str) {
        putchar(*str);
```
Name mangling (C++ only)

Name mangling is the encoding of function and variable names into unique names so that linkers can separate common names in the language. Type names may also be mangled. Name mangling is commonly used to facilitate the overloading feature and visibility within different scopes. The compiler generates function names with an encoding of the types of the function arguments when the module is compiled. If a variable is in a namespace, the name of the namespace is mangled into the variable name so that the same variable name can exist in more than one namespace. The C++ compiler also mangles C variable names to identify the namespace in which the C variable resides.

The scheme for producing a mangled name differs with the object model used to compile the source code: the mangled name of an object of a class compiled using one object model will be different from that of an object of the same class compiled using a different object model. The object model is controlled by compiler option or by pragma.

Name mangling is not desirable when linking C modules with libraries or object files compiled with a C++ compiler. To prevent the C++ compiler from mangling the name of a function, you can apply the extern "C" linkage specifier to the declaration or declarations, as shown in the following example:

```c
extern "C" {
  int f1(int);
  int f2(int);
  int f3(int);
};
```

This declaration tells the compiler that references to the functions f1, f2, and f3 should not be mangled.

The extern "C" linkage specifier can also be used to prevent mangling of functions that are defined in C++ so that they can be called from C. For example,

```c
extern "C" {
  void p(int){
    /* not mangled */
  }
};
```

In multiple levels of nested extern declarations, the innermost extern specification prevails.

```c
extern "C" {
  extern "C++" {
    void func();
  }
}
```

In this example, func has C++ linkage.
Chapter 2. Lexical Elements

A *lexical element* refers to a character or groupings of characters that may legally appear in a source file. This section contains discussions of the basic lexical elements and conventions of the C and C++ programming languages:

- “Tokens”
- “Source program character set” on page 30
- “Comments” on page 35

## Tokens

Source code is treated during preprocessing and compilation as a sequence of *tokens*. A token is the smallest independent unit of meaning in a program, as defined by the compiler. There are four different types of tokens:

- **Keywords**
- **Identifiers**
- **Literals**
- **Punctuators and operators**

Adjacent identifiers, keywords, and literals must be separated with white space. Other tokens should be separated by white space to make the source code more readable. White space includes blanks, horizontal and vertical tabs, new lines, form feeds, and comments.

## Keywords

*Keywords* are identifiers reserved by the language for special use. Although you can use them for preprocessor macro names, it is considered poor programming style. Only the exact spelling of keywords is reserved. For example, *auto* is reserved but *AUTO* is not.

*Table 5. C and C++ keywords*

<table>
<thead>
<tr>
<th>auto</th>
<th>double</th>
<th>int</th>
<th>struct</th>
</tr>
</thead>
<tbody>
<tr>
<td>break</td>
<td>else</td>
<td>long</td>
<td>switch</td>
</tr>
<tr>
<td>case</td>
<td>enum</td>
<td>register</td>
<td>typedef</td>
</tr>
<tr>
<td>char</td>
<td>extern</td>
<td>return</td>
<td>union</td>
</tr>
<tr>
<td>const</td>
<td>float</td>
<td>short</td>
<td>unsigned</td>
</tr>
<tr>
<td>continue</td>
<td>for</td>
<td>signed</td>
<td>void</td>
</tr>
<tr>
<td>default</td>
<td>goto</td>
<td>sizeof</td>
<td>volatile</td>
</tr>
<tr>
<td>do</td>
<td>if</td>
<td>static</td>
<td>while</td>
</tr>
</tbody>
</table>

---

### C only

Standard C at the C99 level also reserves the following keywords:

*Table 6. C99 keywords*

<table>
<thead>
<tr>
<th>_Bool</th>
<th>_Imaginary</th>
</tr>
</thead>
<tbody>
<tr>
<td>_Complex</td>
<td>inline^2</td>
</tr>
<tr>
<td></td>
<td>restrict^3</td>
</tr>
</tbody>
</table>
Notes:

1. The keyword _Imaginary is reserved for future extension of complex number functionality. For current complex number functionality, use _Complex; see “Complex literals” on page 24 for details.

2. The keyword inline is only recognized under compilation with c99 or -qlanglvl=stdc99 or -qlanglvl=extc99 options (or equivalent pragmas) or with the option -qkeyword=inline. Note that the latter option is enabled by default for the xlC invocation command in the configuration file that is shipped with the compiler.

3. The keyword restrict is only recognized under compilation with c99 or with the -qlanglvl=stdc99 or -qlanglvl=extc99 options (or equivalent pragmas) or the -qkeyword=restrict option.

--- End of C only ---

--- C++ only ---

The C++ language also reserves the following keywords:

<table>
<thead>
<tr>
<th>asm</th>
<th>export</th>
<th>private</th>
<th>true</th>
</tr>
</thead>
<tbody>
<tr>
<td>bool</td>
<td>false</td>
<td>protected</td>
<td>try</td>
</tr>
<tr>
<td>catch</td>
<td>friend</td>
<td>public</td>
<td>typeid</td>
</tr>
<tr>
<td>class</td>
<td>inline</td>
<td>reinterpret_cast</td>
<td>typename</td>
</tr>
<tr>
<td>const_cast</td>
<td>mutable</td>
<td>static_cast</td>
<td>using</td>
</tr>
<tr>
<td>delete</td>
<td>namespace</td>
<td>template</td>
<td>virtual</td>
</tr>
<tr>
<td>dynamic_cast</td>
<td>new</td>
<td>this</td>
<td>wchar_t</td>
</tr>
<tr>
<td>explicit</td>
<td>operator</td>
<td>throw</td>
<td></td>
</tr>
</tbody>
</table>

--- End of C++ only ---

Keywords for language extensions

--- IBM extension ---

In addition to standard language keywords, XL C/C++ reserves the following keywords for use in language extensions:

| __restrict  | __real__      | pixel²      |
| __restrict__ | __imag__     | __pixel²   |
| __attribute__ | __complex__ | vector²    |
| __attribute__ | __const__    | __vector²  |
| __signed__   | __inline__   | bool (C only)² |
| __signed__   | __extension__ |            |
| __volatile__ | __label__    |            |
| typeof      | __asm        |            |
| __typeof__  | __asm__      |            |
| __align     | asm          |            |
| __alignof__ |             |            |
| __alignof__ |             |            |
Notes:

1. The `inline` keyword uses the GNU C semantics for inline functions. For details, see “Linkage of inline functions” on page 197.

2. These keywords are recognized only in a vector declaration context, when VMX support is enabled.

---

**C++ only**

XL C++ reserves the following keywords as language extensions for compatibility with C99.

<table>
<thead>
<tr>
<th>Table 9. Keywords for C++ language extensions related to C99</th>
</tr>
</thead>
<tbody>
<tr>
<td>_Complex</td>
</tr>
<tr>
<td>restrict</td>
</tr>
</tbody>
</table>

---

Notes:

1. The keyword _Imaginary is reserved for future extension of complex number functionality. For current complex number functionality, use _Complex; see “Complex literals” on page 24 for details.

---

Related information

- `-qlanglvl` and `-qkeyword` in the XL C/C++ Compiler Reference
- Appendix C, “Vector data types and literals,” on page 415

---

**Identifiers**

*Identifiers* provide names for the following language elements:

- Functions
- Objects
- Labels
- Function parameters
- Macros and macro parameters
- Type definitions
- Enumerated types and enumerators
- Structure and union names
- Classes and class members
- Templates
- Template parameters
- Namespaces

An identifier consists of an arbitrary number of letters, digits, or the underscore character in the form:
The first character in an identifier must be a letter or the _ (underscore) character; however, beginning identifiers with an underscore is considered poor programming style.

The compiler distinguishes between uppercase and lowercase letters in identifiers. For example, PROFIT and profit represent different identifiers. If you specify a lowercase a as part of an identifier name, you cannot substitute an uppercase A in its place; you must use the lowercase letter.

The universal character names for letters and digits outside of the basic source character set are allowed in C++ and at the C99 language level, under compilation with the -qlanglvl=extc99 or -qlanglvl=stdc99 options or related pragmas (for C), or the -qlanglvl=ucs (for C and C++) option.

The dollar sign can appear in identifier names when compiled using the -qdollar compiler option or at one of the extended language levels that encompasses this option.

Identifiers with two initial underscores or an initial underscore followed by an uppercase letter are reserved globally for use by the compiler.

Identifiers that begin with a single underscore are reserved as identifiers with file scope in both the ordinary and tag namespaces.

Identifiers that begin with a single underscore are reserved in the global namespace.

Although the names of system calls and library functions are not reserved words if you do not include the appropriate headers, avoid using them as identifiers. Duplication of a predefined name can lead to confusion for the maintainers of your code and can cause errors at link time or run time. If you include a library in a program, be aware of the function names in that library to avoid name duplications. You should always include the appropriate headers when using standard library functions.

The __func__ predefined identifier

The C99 predefined identifier __func__ makes a function name available for use within the function. Immediately following the opening brace of each function...
definition, __func__ is implicitly declared by the compiler. The resulting behavior is as if the following declaration had been made:

```c
static const char __func__[] = "function-name";
```

where function-name is the name of the lexically-enclosing function. The function name is not mangled.

---

### C++ only

The function name is qualified with the enclosing class name or function name. For example, if foo is a member function of class C, the predefined identifier of foo is C::foo. If foo is defined within the body of main, the predefined identifier of foo is main::C::foo.

The names of template functions or member functions reflect the instantiated type. For example, the predefined identifier for the template function foo instantiated with int, `template<class T> void foo()` is foo<int>.

---

For debugging purposes, you can explicitly use the __func__ identifier to return the name of the function in which it appears. For example:

```c
#include <stdio.h>

void myfunc(void) {
    printf("%s\n", __func__);
    printf("size of __func__ = %d", sizeof(__func__));
}

int main() {
    myfunc();
}
```

The output of the program is:

```
myfunc
size of __func__=7
```

When the assert macro is used inside a function definition, the macro adds the name of the enclosing function on the standard error stream.

**Related information**

- "Function declarations and definitions" on page [191](#)

### Assembly labels

---

### IBM extension

The compiler binds each non-static external variable and function name in the source code to a name that it generates in the object file and any assembly code that is emitted. For compatibility with GCC, the compiler implements an extension to standard C and C++ that allows you to specify the name to be used in the object file and assembly code, by applying an assembly label to the declaration of a global variable or function prototype. You can also define names that do not start with an underscore even on systems where an underscore is normally prepended to the name of a function or variable.
You can use assembly labels with member functions, and functions and variables that are declared in namespaces other than the global namespace.

Assembly label syntax

```
declarator asm ("string_literal")
```

The `string_literal` is a valid assembly name that is to be bound to the given object or function. For a label applied to a function declaration, the name must specify an existing function that is defined in any compilation unit; if no definition is available, a link-time error will occur. For a label applied to a variable declaration, no other definition is required.

The following are examples of assembly label specifications:

```c
void func3() __asm__("foo3");
int i __asm__("abc");
char c __asm__("abc3") = 'a';
```

To distinguish between overloaded functions, XL C++ mangles function names in the object file. Therefore, if you use an assembly label to map a function name, you must use the mangled name of the target function. Furthermore, you must ensure that an assembly label name that you specify for a variable does not conflict with any mangled name. Alternatively, you can prevent name mangling on a target function by declaring it as having C linkage; for more information, see "Name mangling (C++ only)” on page 11.

The following are restrictions on the use of assembly labels:

- Assembly labels cannot be specified on local or static variables.
- The same assembly label name cannot be applied to multiple identifiers (in C++, this is the name after mangling) in the same compilation unit.
- The assembly label name cannot be the same as any other global identifier name (in C++, the name after mangling) in the same compilation unit, unless the label name and identifier name are used for the same variable or function declaration.
- The assembly label cannot be specified on typedef declarations.
- An assembly label cannot be the same as a name specified on a different variable or function by a previous `#pragma map` directive. Similarly, the map name specified by a `#pragma map` directive cannot be the same as a name specified by a previous assembly label on a different variable or function.
- You cannot apply an assembly label to an identifier that has been mapped to a different name by a `#pragma map` directive on a previous declaration of that variable or function. Similarly, you cannot specify a `#pragma map` directive on an identifier that has previously been remapped by an assembly label.
- If you apply different labels to multiple declarations of the same variable or function, the first specification is honored, and all subsequent assembly labels are ignored with a warning.
- You cannot apply an assembly label to any of the following:
member variable declarations
- friend declarations
- template function and member declarations, or any declarations contained within a template
- virtual member functions
- constructors and destructors

Related information
- `#pragma map` in the XL C/C++ Compiler Reference
- “The alias function attribute” on page 207
- “Global variables in specified registers (C only)” on page 48
- “Inline assembly statements” on page 186
- `-qreserved_reg` in the XL C/C++ Compiler Reference

Literals

The term literal constant, or literal, refers to a value that occurs in a program and cannot be changed. The C language uses the term constant in place of the noun literal. The adjective literal adds to the concept of a constant the notion that we can speak of it only in terms of its value. A literal constant is nonaddressable, which means that its value is stored somewhere in memory, but we have no means of accessing that address.

Every literal has a value and a data type. The value of any literal does not change while the program runs and must be in the range of representable values for its type. The following are the available types of literals:

- Integer literals
- Boolean literals
- Floating-point literals
- Vector literals
- Character literals
- String literals

**Integer literals**

Integer literals are numbers that do not have a decimal point or an exponential part. They can be represented as:

- Decimal integer literals
- Hexadecimal integer literals
- Octal integer literals

An integer literal may have a prefix that specifies its base, or a suffix that specifies its type.

**Integer literal syntax**
The data type of an integer literal is determined by its form, value, and suffix. The following table lists the integer literals and shows the possible data types. The smallest data type that can represent the constant value is used to store the constant.

<table>
<thead>
<tr>
<th>Integer literal</th>
<th>Possible data types</th>
</tr>
</thead>
<tbody>
<tr>
<td>unsuffixed decimal</td>
<td>int, long int, unsigned int, long long int¹</td>
</tr>
<tr>
<td>unsuffixed octal or hexadecimal</td>
<td>int, unsigned int, long int, unsigned long int,</td>
</tr>
<tr>
<td></td>
<td>&gt; C long long int¹, C unsigned long int¹</td>
</tr>
<tr>
<td>decimal, octal, or hexadecimal</td>
<td></td>
</tr>
<tr>
<td>suffixed by u or U</td>
<td>unsigned int, unsigned long int, &gt; C</td>
</tr>
<tr>
<td></td>
<td>unsigned long int¹</td>
</tr>
<tr>
<td>decimal suffixed by 1 or L</td>
<td>long int, &gt; C long long int¹</td>
</tr>
<tr>
<td>octal or hexadecimal suffixed by 1 or L</td>
<td>long int, unsigned long int, &gt; C long long int¹</td>
</tr>
<tr>
<td></td>
<td>&gt; C unsigned long int¹</td>
</tr>
<tr>
<td>decimal, octal, or hexadecimal</td>
<td></td>
</tr>
<tr>
<td>suffixed by both u or U, and 1 or L</td>
<td>unsigned long int, &gt; C unsigned long int¹</td>
</tr>
<tr>
<td>decimal suffixed by 11 or LL</td>
<td>long long int</td>
</tr>
<tr>
<td>octal or hexadecimal suffixed by 11 or LL</td>
<td>long long int, unsigned long int</td>
</tr>
<tr>
<td>decimal, octal, or hexadecimal</td>
<td></td>
</tr>
<tr>
<td>suffixed by both u or U, and 11 or LL</td>
<td>unsigned long long int</td>
</tr>
</tbody>
</table>

Notes:
1. Requires compilation with -c99 or with -qlanglvl=stdc99 or -qlanglvl=extc99 or the equivalent pragmas.

Related information
- "Integral types" on page 50
- "Integral conversions" on page 108
- -qlanglvl in the XL C/C++ Compiler Reference

Decimal integer literals: A decimal integer literal contains any of the digits 0 through 9. The first digit cannot be 0. Integer literals beginning with the digit 0 are interpreted as an octal integer literal rather than as a decimal integer literal.
Decimal integer literal syntax

A plus (+) or minus (-) symbol can precede a decimal integer literal. The operator is treated as a unary operator rather than as part of the literal.

The following are examples of decimal literals:
485976
-43312211
+20
5

Hexadecimal integer literals: A hexadecimal integer literal begins with the 0 digit followed by either an x or X, followed by any combination of the digits 0 through 9 and the letters a through f or A through F. The letters A (or a) through F (or f) represent the values 10 through 15, respectively.

Hexadecimal integer literal syntax

The following are examples of hexadecimal integer literals:
0x3b24
0XF96
0x21
0x3AA
0X29b
0X4b0

Octal integer literals: An octal integer literal begins with the digit 0 and contains any of the digits 0 through 7.

Octal integer literal syntax

The following are examples of octal integer literals:
0
0125
034673
03245

Boolean literals

At the C99 level, C defines true and false as macros in the header file stdbool.h.

There are only two Boolean literals: true and false.
Related information

- “Boolean types” on page 50
- “Boolean conversions” on page 108

Floating-point literals

Floating-point literals are numbers that have a decimal point or an exponential part. They can be represented as:

- Real literals
  - Decimal floating-point literals
  - Hexadecimal floating-point literals
- Complex literals

Decimal floating-point literals: A real decimal floating-point constant consists of the following:

- An integral part
- A decimal point
- A fractional part
- An exponent part
- An optional suffix

Both the integral and fractional parts are made up of decimal digits. You can omit either the integral part or the fractional part, but not both. You can omit either the decimal point or the exponent part, but not both.

Decimal floating-point literal syntax

Exponent:

The suffix f or F indicates a type of float, and the suffix l or L indicates a type of long double. If a suffix is not specified, the floating-point constant has a type double.

A plus (+) or minus (-) symbol can precede a floating-point literal. However, it is not part of the literal; it is interpreted as a unary operator.
The following are examples of decimal floating-point literals:

<table>
<thead>
<tr>
<th>Floating-point constant</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>5.3876e4</td>
<td>53,876</td>
</tr>
<tr>
<td>4e-11</td>
<td>0.00000000004</td>
</tr>
<tr>
<td>1e+5</td>
<td>10000</td>
</tr>
<tr>
<td>7.321E-3</td>
<td>0.007321</td>
</tr>
<tr>
<td>3.2E+4</td>
<td>32000</td>
</tr>
<tr>
<td>0.5e-6</td>
<td>0.0000005</td>
</tr>
<tr>
<td>0.45</td>
<td>0.45</td>
</tr>
<tr>
<td>6.e10</td>
<td>60000000000</td>
</tr>
</tbody>
</table>

**Related information**
- “Floating-point types” on page 51
- “Floating-point conversions” on page 108
- “Unary expressions” on page 131

**Hexadecimal floating-point literals:** Real hexadecimal floating constants, which are a C99 feature, consist of the following:

- a hexadecimal prefix
- a significant part
- a binary exponent part
- an optional suffix

The significant part represents a rational number and is composed of the following:

- a sequence of hexadecimal digits (whole-number part)
- an optional fraction part

The optional fraction part is a period followed by a sequence of hexadecimal digits.

The exponent part indicates the power of 2 to which the significant part is raised, and is an optionally signed decimal integer. The type suffix is optional. The full syntax is as follows:

**Hexadecimal floating-point literal syntax**
The suffix \( f \) or \( F \) indicates a type of float, and the suffix \( l \) or \( L \) indicates a type of long double. If a suffix is not specified, the floating-point constant has a type double. You can omit either the whole-number part or the fraction part, but not both. The binary exponent part is required to avoid the ambiguity of the type suffix \( F \) being mistaken for a hexadecimal digit.

**Complex literals:** Complex literals, which are a C99 feature, are constructed in two parts: the real part, and the imaginary part.

**Complex literal syntax**

```
real part + imaginary part
```

```
Imaginary part: floating-point constant * _Complex_I
```

**Real part:**

```
floating-point constant
```

The floating-point constant can be specified as a decimal or hexadecimal floating-point constant (including optional suffixes), in any of the formats described in the previous sections.

_]Complex_I is a macro defined in the complex.h header file, representing the imaginary unit \( i \), the square root of \(-1\).

For example, the declaration:

```c
varComplex = 2.0f + 2.0f * _Complex_I;
```

initializes the complex variable `varComplex` to a value of \(2.0 + 2.0i\).

**IBM extension**

For ease of porting applications developed with GNU C, XL C/C++ also allows you to indicate the imaginary part of a complex literal with a suffix, in addition to the standard suffixes that indicate the type of the complex number (float, double, or long double).

The simplified syntax for a complex literal using the GNU suffixes is as follows:
real part:  
\[ \text{floating-point constant} \]

imaginary part:  
\[ \text{floating-point constant} - \text{imaginary-suffix} \]

floating-point constant can be specified as a decimal or hexadecimal floating-point constant (including optional suffixes), in any of the formats described in the previous sections.

imaginary-suffix is one of the suffixes i, I, j, or J, representing the imaginary unit.

For example, the declaration
\[ \text{varComplex} = 3.0f + 4.0fi; \]
initializes the complex variable varComplex to a value of 3.0 + 4.0i.

Related information
- "Complex floating-point types" on page 52

Vector literals

A vector literal is a constant expression for which the value is interpreted as a vector type. The data type of a vector literal is represented by a parenthesized vector type, and its value is represented by a parenthesized set of constant expressions that represent the vector elements. When all vector elements have the same value, the value of the literal can be represented by a single parenthesized constant expression. Vector literals allow the initialization of vector types.

Vector literal syntax

The \text{vector_type} is a supported vector type.

The \text{constant_expression} can be either of the following:
- A single expression, which will initialize all elements of the vector to the same value
- A comma-separated list of expressions, the number of which is determined by the type of the vector. The number of constant expressions must be exactly:
  4 For vector int, vector long, and vector float types.
For vector short and vector pixel types.

For vector char types.

For example, for an unsigned integer vector type, the literal could be either of the following:

```c
(vector unsigned int)(10) /* initializes all four elements to a value of 10 */
(vector unsigned int)(14, 82, 73, 700) /* initializes the first element to 14, the second element to 82, the third element to 73, and the fourth element to 700 */
```

A vector literal can be cast to another vector type. A vector literal cast does not change the bit pattern of the operand: the 128 bits representing the value remains the same before and after the cast.

A full list of vector literals is provided in [Appendix C, “Vector data types and literals,” on page 415](#).

**Related information**

- “Vector types” on page 53
- “Initialization of vectors” on page 93

---

**Character literals**

A character literal contains a sequence of characters or escape sequences enclosed in single quotation mark symbols, for example 'c'. A character literal may be prefixed with the letter L, for example L'c'. A character literal without the L prefix is an ordinary character literal or a narrow character literal. A character literal with the L prefix is a wide character literal. An ordinary character literal that contains more than one character or escape sequence (excluding single quotes ('), backslashes (\) or new-line characters) is a multicharacter literal.

- **C** The type of a narrow character literal is int. The type of a wide character literal is wchar_t. The type of a multicharacter literal is int.

- **C++** The type of a character literal that contains only one character is char, which is an integral type. The type of a wide character literal is wchar_t. The type of a multicharacter literal is int.

**Character literal syntax**

At least one character or escape sequence must appear in the character literal, and the character literal must appear on a single logical source line.

The characters can be from the source program character set. You can represent the double quotation mark symbol by itself, but to represent the single quotation mark symbol, you must use the backslash symbol followed by a single quotation mark.
Outside of the basic source character set, the universal character names for letters and digits are allowed in C++ and at the C99 language level. To enable universal character names, you must compile with the `c99` invocation command, the `-qlanglvl=extc99` or `-qlanglvl=stdc99` options or related pragmas (for C), or the `-qlanglvl=ucs` (for C and C++) option.

The following are examples of character literals:

'`a`

'`'

L'`0`

'`('  

Related information

- “Source program character set” on page 30
- “The Unicode standard” on page 32
- “Character types” on page 52

String literals

A string literal contains a sequence of characters or escape sequences enclosed in double quotation mark symbols. A string literal with the prefix L is a wide string literal. A string literal without the prefix L is an ordinary or narrow string literal.

- **C** The type of a narrow string literal is array of char. The type of a wide string literal is array of wchar_t.

- **C++** The type of a narrow string literal is array of const char. The type of a wide string literal is array of const wchar_t. Both types have static storage duration.

A null ('`\0`') character is appended to each string. For a wide string literal, the value '`\0`' of type wchar_t is appended. By convention, programs recognize the end of a string by finding the null character.

String literal syntax

Multiple spaces contained within a string literal are retained.

Use the escape sequence `\n` to represent a new-line character as part of the string. Use the escape sequence `\\` to represent a backslash character as part of the string. You can represent a single quotation mark symbol either by itself or with the escape sequence `\'`. You must use the escape sequence `\"` to represent a double quotation mark.

Outside of the basic source character set, the universal character names for letters and digits are allowed in C++ and at the C99 language level. To enable universal
character names, you must compile with the c99 invocation command, the
-qlanglvl=extc99 or -qlanglvl=stdc99 options or related pragmas (for C), or the
-qlanglvl=ucs (for C and C++) option.

The following are examples of string literals:
char titles[] = "Handel's Water Music";
char *temp_string = "abc def ghi"; /* temp_string = "abcdefghi\0" */
wchar_t *wide_string = L"longstring";

This example illustrates escape sequences in string literals:
#include <iostream> using namespace std;

int main () {
  char *s = "Hi there! \n";
  cout << *s;
  char *p = "The backslash character \\
.";
  cout << *p << endl;
  char *q = "The double quotation mark \".\n";
  cout << *q;
}

This program produces the following output:
Hi there! The backslash character \. The double quotation mark ".

To continue a string on the next line, use the line continuation character (\ symbol)
followed by optional whitespace and a new-line character (required). For example:
char *mail_addr = "Last Name  First Name  MI  Street Address \
893   City   Province   Postal code ";

In the following example, the string literal second causes a compile-time error.
char *first = "This string continues onto the next\n  line, where it ends."; /* compiles successfully. */
char *second = "The comment makes the \   /* continuation symbol
  */ invisible to the compiler."; /* compilation error. */

Note: When a string literal appears more than once in the program source, how
that string is stored depends on whether strings are read-only or writeable.
By default, the compiler considers strings to be read-only. XL C/C++ may
allocate only one location for a read-only string; all occurrences will refer to
that one location. However, that area of storage is potentially
write-protected. If strings are writeable, then each occurrence of the string
will have a separate, distinct storage location that is always modifiable. You
can use the #pragma strings directive or the -ro compiler option to change
the default storage for string literals.

Related information
• “Character types” on page 52
• “Source program character set” on page 30
• “The Unicode standard” on page 32
• -ro in the XL C/C++ Compiler Reference
• #pragma strings in the XL C/C++ Compiler Reference
• -qsourcectype in the XL C/C++ Compiler Reference
• -qmacpstr in the XL C/C++ Compiler Reference
**String concatenation:** Another way to continue a string is to have two or more consecutive strings. Adjacent string literals will be concatenated to produce a single string. For example:

```c
"hello = "there" /* is equivalent to "hello there" */
"hello" "there" /* is equivalent to "hellothere" */
```

Characters in concatenated strings remain distinct. For example, the strings "\xab" and "3" are concatenated to form "\xab3". However, the characters \xab and 3 remain distinct and are not merged to form the hexadecimal character \xab3.

If a wide string literal and a narrow string literal are adjacent, as in the following:

```
"hello " L"there"
```

the result is a wide string literal.

Following any concatenation, '\0' of type char is appended at the end of each string. For a wide string literal, '\0' of type wchar_t is appended. C++ programs find the end of a string by scanning for this value. For example:

```c
char *first = "Hello "; /* stored as "Hello \0" */
char *second = " there" ; /* stored as "there\0" */
char *third = "Hello = "there" ; /* stored as "Hello there\0" */
```

**Related information**
- "String concatenation of u-literals" on page 34

**Punctuators and operators**

A punctuator is a token that has syntactic and semantic meaning to the compiler, but the exact significance depends on the context. A punctuator can also be a token that is used in the syntax of the preprocessor.

C99 and C++ define the following tokens as punctuators, operators, or preprocessing tokens:

**Table 10. C and C++ punctuators**

<p>| | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>[]</td>
<td>()</td>
<td>{}</td>
<td>,</td>
</tr>
<tr>
<td>*</td>
<td>=</td>
<td>...</td>
<td>#</td>
</tr>
<tr>
<td>.</td>
<td>-&gt;</td>
<td>++</td>
<td>--</td>
</tr>
<tr>
<td>&amp;</td>
<td>+</td>
<td>-</td>
<td>~</td>
</tr>
<tr>
<td>/</td>
<td>%</td>
<td>&lt;&lt;</td>
<td>&gt;&gt;</td>
</tr>
<tr>
<td>&lt;</td>
<td>&gt;</td>
<td>&lt;=</td>
<td>&gt;=</td>
</tr>
<tr>
<td>^</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>*=</td>
<td>/=</td>
<td>%=</td>
<td>+=</td>
</tr>
<tr>
<td>&lt;&lt;=</td>
<td>&gt;&gt;=</td>
<td>&amp; =</td>
<td>^=</td>
</tr>
</tbody>
</table>

In addition to the C99 preprocessing tokens, operators, and punctuators, C++ allows the following tokens as punctuators:

**Table 11. C++ punctuators**

<p>| | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>:=</td>
<td>.*</td>
<td>-*</td>
<td>new</td>
</tr>
<tr>
<td>and</td>
<td>and_eq</td>
<td>bitand</td>
<td>bitor</td>
</tr>
<tr>
<td>not</td>
<td>not_eq</td>
<td>or</td>
<td>or_eq</td>
</tr>
<tr>
<td></td>
<td>xor</td>
<td>xor_eq</td>
<td></td>
</tr>
</tbody>
</table>

---

**End of C++ only**
Related information

- [Chapter 6, “Expressions and operators,” on page 115](#)

**Alternative tokens**
Both C and C++ provide the following alternative representations for some operators and punctuators. The alternative representations are also known as *digraphs*.

<table>
<thead>
<tr>
<th>Operator or punctuator</th>
<th>Alternative representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>{</td>
<td>&lt;%</td>
</tr>
<tr>
<td>}</td>
<td>%&gt;</td>
</tr>
<tr>
<td>[</td>
<td>&lt;:</td>
</tr>
<tr>
<td>]</td>
<td>:&gt;</td>
</tr>
<tr>
<td>#</td>
<td>%:</td>
</tr>
<tr>
<td>#</td>
<td>%:%:</td>
</tr>
</tbody>
</table>

In addition to the operators and punctuators listed above, C++ and C at the C99 language level provide the following alternative representations. In C, they are defined as macros in the header file `iso646.h`.

<table>
<thead>
<tr>
<th>Operator or punctuator</th>
<th>Alternative representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>&amp;&amp;</td>
<td>and</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td>^</td>
<td>xor</td>
</tr>
<tr>
<td>~</td>
<td>compl</td>
</tr>
<tr>
<td>&amp;</td>
<td>bitand</td>
</tr>
<tr>
<td>&amp;</td>
<td>=</td>
</tr>
<tr>
<td></td>
<td>=</td>
</tr>
<tr>
<td>^=</td>
<td>xor_eq</td>
</tr>
<tr>
<td>!</td>
<td>not</td>
</tr>
<tr>
<td>!=</td>
<td>not_eq</td>
</tr>
</tbody>
</table>

Related information

- [“Digraph characters” on page 34](#)

**Source program character set**

The following lists the basic source character sets that are available at both compile time and run time:

- The uppercase and lowercase letters of the English alphabet:
  
  a b c d e f g h i j k l m n o p q r s t u v w x y z
  
  A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

- The decimal digits:
  
  0 1 2 3 4 5 6 7 8 9

- The following graphic characters:
  
  ! " # $ % & ' ( ) * + , - . / : ; < = > ? [ \ ] _ { } ~
– The caret (^) character in ASCII (bitwise exclusive OR symbol).
– The split vertical bar (¦) character in ASCII.

- The space character
- The control characters representing new-line, horizontal tab, vertical tab, form feed, end of string (NULL character), alert, backspace, and carriage return.

IBM Depending on the compiler option, other specialized identifiers, such as the dollar sign ($) or characters in national character sets, may be allowed to appear in an identifier.

Related information
• “Characters in identifiers” on page 16

Multibyte characters
The compiler recognizes and supports the additional characters (the extended character set) which you can meaningfully use in string literals and character constants. The support for extended characters includes the multibyte character sets. A multibyte character is a character whose bit representation fits into one or more bytes. To instruct the compiler to recognize multibyte character sets as source input, be sure to compile with the [-qmbcs] option.

Related information
• “Character literals” on page 26
• “The Unicode standard” on page 32
• “Character types” on page 52
• [-qmbcs] in the [XL C/C++ Compiler Reference]
• “Multibyte character support” in the [XL C/C++ Compiler Reference]

Escape sequences
You can represent any member of the execution character set by an escape sequence. They are primarily used to put nonprintable characters in character and string literals. For example, you can use escape sequences to put such characters as tab, carriage return, and backspace into an output stream.

Escape character syntax

```
\x<hexadecimal_digits>
```

An escape sequence contains a backslash (\) symbol followed by one of the escape sequence characters or an octal or hexadecimal number. A hexadecimal escape sequence contains an x followed by one or more hexadecimal digits (0-9, A-F, a-f). An octal escape sequence uses up to three octal digits (0-7). The value of the hexadecimal or octal number specifies the value of the desired character or wide character.

Note: The line continuation sequence (\ followed by a new-line character) is not an escape sequence. It is used in character strings to indicate that the current line of source code continues on the next line.
The escape sequences and the characters they represent are:

<table>
<thead>
<tr>
<th>Escape sequence</th>
<th>Character represented</th>
</tr>
</thead>
<tbody>
<tr>
<td>\a</td>
<td>Alert (bell, alarm)</td>
</tr>
<tr>
<td>\b</td>
<td>Backspace</td>
</tr>
<tr>
<td>\f</td>
<td>Form feed (new page)</td>
</tr>
<tr>
<td>\n</td>
<td>New-line</td>
</tr>
<tr>
<td>\r</td>
<td>Carriage return</td>
</tr>
<tr>
<td>\t</td>
<td>Horizontal tab</td>
</tr>
<tr>
<td>\v</td>
<td>Vertical tab</td>
</tr>
<tr>
<td>'</td>
<td>Single quotation mark</td>
</tr>
<tr>
<td>&quot;</td>
<td>Double quotation mark</td>
</tr>
<tr>
<td>?</td>
<td>Question mark</td>
</tr>
<tr>
<td>&quot;</td>
<td>Backslash</td>
</tr>
</tbody>
</table>

The value of an escape sequence represents the member of the character set used at run time. Escape sequences are translated during preprocessing. For example, on a system using the ASCII character codes, the value of the escape sequence \x56 is the letter V. On a system using EBCDIC character codes, the value of the escape sequence \xE5 is the letter V.

Use escape sequences only in character constants or in string literals. An error message is issued if an escape sequence is not recognized.

In string and character sequences, when you want the backslash to represent itself (rather than the beginning of an escape sequence), you must use a \\ backslash escape sequence. For example:

```cpp
    cout << "The escape sequence \n." << endl;
```

This statement results in the following output:

```
The escape sequence \n.
```

**The Unicode standard**

The *Unicode Standard* is the specification of an encoding scheme for written characters and text. It is a universal standard that enables consistent encoding of multilingual text and allows text data to be interchanged internationally without conflict. The ISO standards for C and C++ refer to ISO/IEC 10646–1:2000, *Information Technology—Universal Multiple-Octet Coded Character Set (UCS)*. (The term *octet* is used by ISO to refer to a byte.) The ISO/IEC 10646 standard is more restrictive than the Unicode Standard in the number of encoding forms: a character set that conforms to ISO/IEC 10646 is also conformant to the Unicode Standard.

The Unicode Standard specifies a unique numeric value and name for each character and defines three encoding forms for the bit representation of the numeric value. The name/value pair creates an identity for a character. The hexadecimal value representing a character is called a *code point*. The specification also describes overall character properties, such as case, directionality, alphabetic properties, and other semantic information for each character. Modeled on ASCII, the Unicode Standard treats alphabetic characters, ideographic characters, and symbols, and allows implementation-defined character codes in reserved code point ranges. The encoding scheme of the Unicode Standard is therefore sufficiently flexible to handle all known character encoding requirements, including coverage of historical scripts from any country in the world.
C99 and C++ allow the universal character name construct defined in ISO/IEC 10646 to represent characters outside the basic source character set. Both languages permit universal character names in identifiers, character constants, and string literals. To enable universal character names, you must compile with the `c99` invocation command, the `qlanglvl=extc99` or `qlanglvl=stdc99` options or related pragmas (for C), or the `qlanglvl=ucs` option (for C and C++).

The following table shows the generic universal character name construct and how it corresponds to the ISO/IEC 10646 short name.

<table>
<thead>
<tr>
<th>Universal character name</th>
<th>ISO/IEC 10646 short name</th>
</tr>
</thead>
<tbody>
<tr>
<td>where N is a hexadecimal digit</td>
<td></td>
</tr>
<tr>
<td>\UNNNNNNNNN</td>
<td>NNNNNNNNN</td>
</tr>
<tr>
<td>\uNNNN</td>
<td>0000NNNN</td>
</tr>
</tbody>
</table>

C99 and C++ disallow the hexadecimal values representing characters in the basic character set (base source code set) and the code points reserved by ISO/IEC 10646 for control characters.

The following characters are also disallowed:
- Any character whose short identifier is less than 00A0. The exceptions are 0024 ($), 0040 (@), or 0060 (’).
- Any character whose short identifier is in the code point range D800 through DFFF inclusive.

**UTF literals**

The C Standards Committee has approved the implementation of *u-literals* and *U-literals* to support Unicode UTF-16 and UTF-32 character literals, respectively. To enable support for UTF literals in your source code, you must compile with the option `qlanglvl=extc99` enabled. The following table shows the syntax for UTF literals.

<table>
<thead>
<tr>
<th>Syntax</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>u'character'</td>
<td>Denotes a UTF-16 character.</td>
</tr>
<tr>
<td>u&quot;character-sequence&quot;</td>
<td>Denotes an array of UTF-16 characters.</td>
</tr>
<tr>
<td>U'character'</td>
<td>Denotes a UTF-32 character.</td>
</tr>
<tr>
<td>U&quot;character-sequence&quot;</td>
<td>Denotes an array of UTF-32 characters.</td>
</tr>
</tbody>
</table>

XL C/C++ implements the macros `uint_least16_t` and `uint_least32_t`, which are defined in the header file `stdint.h`, as data types for UTF-16 and UTF-32 characters. The following example defines an array of `uint_least16_t`, including the characters represented by code points 1234 and 8180:

```c
#include <stdint.h>

uint_least16_t msg[] = u"ucs characters \u1234 and \u8180 ";
```

Related information
- [qlref] in the **XL C/C++ Compiler Reference**
String concatenation of u-literals: The u-literals and U-literals follow the same concatenation rule as wide character literals: the normal character string is widened if they are present. The following shows the allowed combinations. All other combinations are invalid.

<table>
<thead>
<tr>
<th>Combination</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>u“a” u“b”</td>
<td>u“ab”</td>
</tr>
<tr>
<td>u“a” “b”</td>
<td>u“ab”</td>
</tr>
<tr>
<td>“a” u“b”</td>
<td>u“ab”</td>
</tr>
<tr>
<td>U“a” U“b”</td>
<td>U“ab”</td>
</tr>
<tr>
<td>U“a” “b”</td>
<td>U“ab”</td>
</tr>
<tr>
<td>“a” U“b”</td>
<td>U“ab”</td>
</tr>
</tbody>
</table>

Multiple concatenations are allowed, with these rules applied recursively.

Related information
- [String concatenation](#)

---

End of IBM extension

---

**Digraph characters**

You can represent unavailable characters in a source program by using a combination of two keystrokes that are called a digraph character. The preprocessor reads digraphs as tokens during the preprocessor phase.

The digraph characters are:

<table>
<thead>
<tr>
<th>%: or %%</th>
<th>#</th>
<th>number sign</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;:</td>
<td>[</td>
<td>left bracket</td>
</tr>
<tr>
<td>:&gt;</td>
<td>]</td>
<td>right bracket</td>
</tr>
<tr>
<td>&lt;%%</td>
<td>{</td>
<td>left brace</td>
</tr>
<tr>
<td>%&gt;:</td>
<td>}</td>
<td>right brace</td>
</tr>
<tr>
<td>%:%: or %%%%</td>
<td>##</td>
<td>preprocessor macro concatenation operator</td>
</tr>
</tbody>
</table>

You can create digraphs by using macro concatenation. XL C/C++ does not replace digraphs in string literals or in character literals. For example:

```c
char *s = "<%%"; // stays "<%%"
```

```c
switch (c) {
    case '<%': { /* ... */ } // stays '<%'
    case '%>': { /* ... */ } // stays '%>'
}
```

Related information
- [qdigraph](#) in the [XL C/C++ Compiler Reference](#)

**Trigraph sequences**

Some characters from the C and C++ character set are not available in all environments. You can enter these characters into a C or C++ source program using a sequence of three characters called a trigraph. The trigraph sequences are:
The preprocessor replaces trigraph sequences with the corresponding single-character representation. For example,

```
some_array??(??) = n;
```

Represents:

```
some_array[i] = n;
```

### Comments

A *comment* is text replaced during preprocessing by a single space character; the compiler therefore ignores all comments.

There are two kinds of comments:

- The /* (slash, asterisk) characters, followed by any sequence of characters (including new lines), followed by the */ characters. This kind of comment is commonly called a C-style comment.

- The // (two slashes) characters followed by any sequence of characters. A new line not immediately preceded by a backslash terminates this form of comment. This kind of comment is commonly called a single-line comment or a C++ comment. A C++ comment can span more than one physical source line if it is joined into one logical source line with line-continuation (\) characters. The backslash character can also be represented by a trigraph.

You can put comments anywhere the language allows white space. You cannot nest C-style comments inside other C-style comments. Each comment ends at the first occurrence of */.

You can also include multibyte characters; to instruct the compiler to recognize multibyte characters in the source code, compile with the -qmbcs option.

**Note:** The /* or */ characters found in a character constant or string literal do not start or end comments.

In the following program, the second printf() is a comment:

```c
#include <stdio.h>

int main(void)
{
```

---

<table>
<thead>
<tr>
<th>Trigraph</th>
<th>Single character</th>
<th>capription</th>
</tr>
</thead>
<tbody>
<tr>
<td>??=</td>
<td>#</td>
<td>pound sign</td>
</tr>
<tr>
<td>??(</td>
<td>[</td>
<td>left bracket</td>
</tr>
<tr>
<td>??)</td>
<td>]</td>
<td>right bracket</td>
</tr>
<tr>
<td>??&lt;</td>
<td>{</td>
<td>left brace</td>
</tr>
<tr>
<td>??&gt;</td>
<td>}</td>
<td>right brace</td>
</tr>
<tr>
<td>??/</td>
<td>\</td>
<td>backslash</td>
</tr>
<tr>
<td>??'</td>
<td>'</td>
<td>caret</td>
</tr>
<tr>
<td>??!</td>
<td></td>
<td></td>
</tr>
<tr>
<td>??-</td>
<td>~</td>
<td>tilde</td>
</tr>
</tbody>
</table>

---

Chapter 2. Lexical elements 35
printf("This program has a comment.\n");
/* printf("This is a comment line and will not print.\n"); */
return 0;
}

Because the second printf() is equivalent to a space, the output of this program is:
This program has a comment.

Because the comment delimiters are inside a string literal, printf() in the following program is not a comment.
#include <stdio.h>

int main(void)
{
    printf("This program does not have \
    /* NOT A COMMENT */ a comment.\n");
    return 0;
}

The output of the program is:
This program does not have
/* NOT A COMMENT */ a comment.

In the following example, the comments are highlighted:

/* A program with nested comments. */
#include <stdio.h>

int main(void)
{
    test_function();
    return 0;
}

int test_function(void)
{
    int number;
    char letter;
    /*
        number = 55;
        letter = 'A';
    */
    /*
        number = 44;
    */
    return 999;
}

In test_function, the compiler reads the first /* through to the first */. The second */ causes an error. To avoid commenting over comments already in the source code, you should use conditional compilation preprocessor directives to cause the compiler to bypass sections of a program. For example, instead of commenting out the above statements, change the source code in the following way:

/* A program with conditional compilation to avoid nested comments. */

#define TEST_FUNCTION 0
#include <stdio.h>

int main(void)
{
    test_function();
    return 0;
}
int test_function(void)
{
    int number;
    char letter;
#if TEST_FUNCTION
    number = 55;
    letter = 'A';
    /*number = 44;*/
#endif /*TEST_FUNCTION */
}

You can nest single line comments within C-style comments. For example, the following program will not output anything:

#include <stdio.h>

int main(void)
{
    /*
    printf("This line will not print.\n");
    // This is a single line comment
    // This is another single line comment
    printf("This line will also not print.\n");
    */
    return 0;
}

Note: You can also use the #pragma comment directive to place comments into an object module.

Related information
- `-qmbcs [-qlanglvl]` and `qcpluscm` in the `XL C/C++ Compiler Reference`
- "Multibyte character support" in the `XL C/C++ Compiler Reference"
Chapter 3. Data objects and declarations

This section discusses the various elements that constitute a declaration of a data object, and includes the following topics:

- "Overview of data objects and declarations"
- "Storage class specifiers" on page 43
- "Type specifiers" on page 49
- "User-defined types" on page 55
- "Type qualifiers" on page 68
- "Type attributes" on page 74

Topics are sequenced to loosely follow the order in which elements appear in a declaration. The discussion of the additional elements of data declarations is also continued in Chapter 4, “Declarators,” on page 79.

Overview of data objects and declarations

The following sections introduce some fundamental concepts regarding data objects and data declarations that will be used throughout this reference.

Overview of data objects

A data object is a region of storage that contains a value or group of values. Each value can be accessed using its identifier or a more complex expression that refers to the object. In addition, each object has a unique data type. The data type of an object determines the storage allocation for that object and the interpretation of the values during subsequent access. It is also used in any type checking operations. Both the identifier and data type of an object are established in the object declaration.

An instance of a class type is commonly called a class object. The individual class members are also called objects. The set of all member objects comprises a class object.

Data types are often grouped into type categories that overlap, such as:

Fundamental types versus derived types

Fundamental data types are also known as "basic", "fundamental" or "built-in" to the language. These include integers, floating-point numbers, and characters. Derived types, also known as "compound" types in Standard C++, are created from the set of basic types, and include arrays, pointers, structures, unions, and enumerations. All C++ classes are considered compound types.

Built-in types versus user-defined types

Built-in data types include all of the fundamental types, plus types that refer to the addresses of basic types, such as arrays and pointers. User-defined types are created by the user from the set of basic types, in typedef, structure, union, and enumeration definitions. C++ classes are considered user-defined types.

Scalar types versus aggregate types

Scalar types represent a single data value, while aggregate types represent
multiple values, of the same type or of different types. Scalars include the
arithmetic types and pointers. Aggregate types include arrays and
structures. C++ classes are considered aggregate types.

The following matrix lists the supported data types and their classification into
fundamental, derived, scalar, and aggregate types.

Table 13. C/C++ data types

<table>
<thead>
<tr>
<th>Data object</th>
<th>Basic</th>
<th>Compound</th>
<th>Built-in</th>
<th>User-defined</th>
<th>Scalar</th>
<th>Aggregate</th>
</tr>
</thead>
<tbody>
<tr>
<td>integer types</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>floating-point</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>types(^3)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>character types</td>
<td></td>
<td></td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Booleans</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>void type</td>
<td>+(^2)</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>pointers</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>arrays</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>structures</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>unions</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>enumerations</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>see note(^3)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>C++ classes</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
<tr>
<td>typedef types</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td>+</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Notes:
1. Although complex floating-point types are represented internally as an array
of two elements, they behave in the same way as real floating-pointing types in
terms of alignment and arithmetic operations, and can therefore be considered
scalar types.
2. The void type is really an incomplete type, as discussed in “Incomplete types.”
   Nevertheless, Standard C++ defines it as a fundamental type.
3. The C standard does not classify enumerations as either scalar or
   aggregate. Standard C++ classifies enumerations as scalars.

Related information
- [Chapter 11, “Classes (C++ only),” on page 245](#)

Incomplete types
The following are incomplete types:
- The void type
- Arrays of unknown size
- Arrays of elements that are of incomplete type
- Structure, union, or enumerations that have no definition
- Pointers to class types that are declared but not defined
- Classes that are declared but not defined
However, if an array size is specified by [ * ], indicating a variable length array, the size is considered as having been specified, and the array type is then considered a complete type. For more information, see “Variable length arrays” on page 87.

The following examples illustrate incomplete types:

```c
void *incomplete_ptr;
struct dimension linear; /* no previous definition of dimension */
```

**Related information**

- “The void type” on page 53
- “Incomplete class declarations” on page 250

**Compatible types**

C only

In C, compatible types are defined as:

- two types that can be used together without modification (as in an assignment expression)
- two types that can be substituted one for the other without modification

When two compatible types are combined, the result is a composite type. Determining the resultant composite type for two compatible types is similar to following the usual binary conversions of integral types when they are combined with some arithmetic operators.

Obviously, two types that are identical are compatible; their composite type is the same type. Less obvious are the rules governing type compatibility of non-identical types, user-defined types, type-qualified types, and so on. “Type specifiers” on page 49 discusses compatibility for basic and user-defined types in C.

C++ only

A separate notion of type compatibility as distinct from being of the same type does not exist in C++. Generally speaking, type checking in C++ is stricter than in C: identical types are required in situations where C would only require compatible types.

**Related information**

- “Compatibility of arrays” on page 88
- “Compatibility of pointers (C only)” on page 85
- “Compatible functions” on page 194

**Overview of data declarations and definitions**

A declaration establishes the names and characteristics of data objects used in a program. A definition allocates storage for data objects, and associates an identifier with that object. When you declare or define a type, no storage is allocated.
The following table shows examples of declarations and definitions. The identifiers declared in the first column do not allocate storage; they refer to a corresponding definition. The identifiers declared in the second column allocate storage; they are both declarations and definitions.

<table>
<thead>
<tr>
<th>Declarations</th>
<th>Declarations and definitions</th>
</tr>
</thead>
<tbody>
<tr>
<td>extern double pi;</td>
<td>double pi = 3.14159265;</td>
</tr>
<tr>
<td>struct payroll;</td>
<td>struct payroll {</td>
</tr>
<tr>
<td></td>
<td>char *name;</td>
</tr>
<tr>
<td></td>
<td>float salary;</td>
</tr>
<tr>
<td></td>
<td>} employee;</td>
</tr>
</tbody>
</table>

**Note:** The C99 standard no longer requires that all declarations appear at the beginning of a function before the first statement. As in C++, you can mix declarations with other statements in your code.

Declarations determine the following properties of data objects and their identifiers:

- **Scope**, which describes the region of program text in which an identifier can be used to access its object
- **Visibility**, which describes the region of program text from which legal access can be made to the identifier’s object
- **Duration**, which defines the period during which the identifiers have real, physical objects allocated in memory
- **Linkage**, which describes the correct association of an identifier to one particular object
- **Type**, which determines how much memory is allocated to an object and how the bit patterns found in the storage allocation of that object should be interpreted by the program

The elements of a declaration for a data object are as follows:

- **Storage class specifiers**, which specify storage duration and linkage
- **Type specifiers**, which specify data types
- **Type qualifiers**, which specify the mutability of data values
- **Declarators**, which introduce and include identifiers
- **Initializers**, which initialize storage with initial values

---

**IBM extension**

In addition, for compatibility with GCC, XL C/C++ allows you to use **attributes** to modify the properties of data objects. **Type attributes**, which can be used to modify the definition of user-defined types, are described in “Type attributes” on page 74. **Variable attributes**, which can be used to modify the declaration of variables, are described in “Variable attributes” on page 101.

---

End of IBM extension

All declarations have the form:
Data declaration syntax

```
<table>
<thead>
<tr>
<th>storage_classSpecifier</th>
<th>typeSpecifier</th>
</tr>
</thead>
<tbody>
<tr>
<td>,</td>
<td></td>
</tr>
<tr>
<td>declarator</td>
<td>initializer</td>
</tr>
<tr>
<td>;</td>
<td></td>
</tr>
</tbody>
</table>
```

Related information
• [“Function declarations and definitions” on page 191](#)

**Tentative definitions**

A *tentative definition* is any external data declaration that has no storage class specifier and no initializer. A tentative definition becomes a full definition if the end of the translation unit is reached and no definition has appeared with an initializer for the identifier. In this situation, the compiler reserves uninitialized space for the object defined.

The following statements show normal definitions and tentative definitions.

```
int i1 = 10;  /* definition, external linkage */
static int i2 = 20; /* definition, internal linkage */
extern int i3 = 30; /* definition, external linkage */
int i4;        /* tentative definition, external linkage */
static int i5; /* tentative definition, internal linkage */
int i1;        /* valid tentative definition */
int i2;        /* not legal, linkage disagreement with previous */
int i3;        /* valid tentative definition */
int i4;        /* valid tentative definition */
int i5;        /* not legal, linkage disagreement with previous */
```

--- End of C only ---

### C++ only

C++ does not support the concept of a tentative definition: an external data declaration without a storage class specifier is always a definition.

--- End of C++ only ---

**Storage class specifiers**

A storage class specifier is used to refine the declaration of a variable, a function, and parameters. Storage classes determine whether:

- The object has internal, external, or no linkage
- The object is to be stored in memory or in a register, if available
- The object receives the default initial value of 0 or an indeterminate default initial value
The object can be referenced throughout a program or only within the function, block, or source file where the variable is defined.

The storage duration for the object is maintained throughout program run time only during the execution of the block where the object is defined.

For a variable, its default storage duration, scope, and linkage depend on where it is declared: whether inside or outside a block statement or the body of a function. When these defaults are not satisfactory, you can use a storage class specifier to explicitly set its storage class. The storage class specifiers in C and C++ are:

- `auto`
- `static`
- `extern`
- `mutable`
- `register`

Related information
- “Function storage class specifiers” on page 196
- “Initializers” on page 89

### The auto storage class specifier

The `auto` storage class specifier lets you explicitly declare a variable with automatic storage. The `auto` storage class is the default for variables declared inside a block. A variable `x` that has automatic storage is deleted when the block in which `x` was declared exits.

You can only apply the auto storage class specifier to names of variables declared in a block or to names of function parameters. However, these names by default have automatic storage. Therefore the storage class specifier `auto` is usually redundant in a data declaration.

### Storage duration of automatic variables

Objects with the auto storage class specifier have automatic storage duration. Each time a block is entered, storage for auto objects defined in that block is made available. When the block is exited, the objects are no longer available for use. An object declared with no linkage specification and without the static storage class specifier has automatic storage duration.

If an auto object is defined within a function that is recursively invoked, memory is allocated for the object at each invocation of the block.

### Linkage of automatic variables

An auto variable has block scope and no linkage.

Related information
- “Initialization and storage classes” on page 90
- “Block statements” on page 170
- “The goto statement” on page 184

### The static storage class specifier

Objects declared with the static storage class specifier have static storage duration, which means that memory for these objects is allocated when the program begins.
running and is freed when the program terminates. Static storage duration for a variable is different from file or global scope: a variable can have static duration but local scope.

- C The keyword `static` is the major mechanism in C to enforce information hiding.

- C++ C++ enforces information hiding through the namespace language feature and the access control of classes. The use of the keyword `static` to limit the scope of external variables is deprecated for declaring objects in namespace scope.

The `static` storage class specifier can be applied to the following declarations:
- Data objects
- Class members
- Anonymous unions

You cannot use the `static` storage class specifier with the following:
- Type declarations
- Function parameters

At the C99 language level, the `static` keyword can be used in the declaration of an array parameter to a function. The `static` keyword indicates that the argument passed into the function is a pointer to an array of at least the specified size. In this way, the compiler is informed that the pointer argument is never null. See “Static array indices in function parameter declarations (C only)” on page 205 for more information.

Related information
- “The static storage class specifier” on page 196
- “Static members” on page 265

**Linkage of static variables**
A declaration of an object that contains the static storage class specifier and has file scope gives the identifier internal linkage. Each instance of the particular identifier therefore represents the same object within one file only. For example, if a static variable \( x \) has been declared in function \( f \), when the program exits the scope of \( f \), \( x \) is not destroyed:

```c
#include <stdio.h>

int f(void) {
    static int x = 0;
    x++;
    return x;
}

int main(void) {
    int j;
    for (j = 0; j < 5; j++) {
```
```c
    printf("Value of f(): %d\n", f());
    return 0;
}
```

The following is the output of the above example:

Value of f(): 1  
Value of f(): 2  
Value of f(): 3  
Value of f(): 4  
Value of f(): 5

Because x is a static variable, it is not reinitialized to 0 on successive calls to f.

**Related information**

- ["Initialization and storage classes" on page 90](#)
- ["Internal linkage" on page 8](#)
- [Chapter 9, “Namespaces (C++ only),” on page 221](#)

**The extern storage class specifier**

The `extern` storage class specifier lets you declare objects that several source files can use. An `extern` declaration makes the described variable usable by the succeeding part of the current source file. This declaration does not replace the definition. The declaration is used to describe the variable that is externally defined.

An `extern` declaration can appear outside a function or at the beginning of a block. If the declaration describes a function or appears outside a function and describes an object with external linkage, the keyword `extern` is optional.

If a declaration for an identifier already exists at file scope, any `extern` declaration of the same identifier found within a block refers to that same object. If no other declaration for the identifier exists at file scope, the identifier has external linkage.

> **C++**  

C++ restricts the use of the `extern` storage class specifier to the names of objects or functions. Using the `extern` specifier with type declarations is illegal. An `extern` declaration cannot appear in class scope.

**Storage duration of external variables**

All `extern` objects have static storage duration. Memory is allocated for `extern` objects before the `main` function begins running, and is freed when the program terminates. The scope of the variable depends on the location of the declaration in the program text. If the declaration appears within a block, the variable has block scope; otherwise, it has file scope.

**Linkage of external variables**

> **C** Like the scope, the linkage of a variable declared `extern` depends on the placement of the declaration in the program text. If the variable declaration appears outside of any function definition and has been declared `static` earlier in the file, the variable has internal linkage; otherwise, it has external linkage in most cases. All object declarations that occur outside a function and that do not contain a storage class specifier declare identifiers with external linkage.
For objects in the unnamed namespace, the linkage may be external, but the name is unique, and so from the perspective of other translation units, the name effectively has internal linkage.

Related information
- “External linkage” on page 9
- “Initialization and storage classes” on page 90
- “The extern storage class specifier” on page 196
- “Class scope (C++ only)” on page 5
- Chapter 9, “Namespaces (C++ only),” on page 221

The mutable storage class specifier (C++ only)

The mutable storage class specifier is used only on a class data member to make it modifiable even though the member is part of an object declared as const. You cannot use the mutable specifier with names declared as static or const, or reference members.

In the following example:

```cpp
class A
{
    public:
        A() : x(4), y(5) {}
        mutable int x;
        int y;
};

int main()
{
    const A var2;
    var2.x = 345;
    // var2.y = 2345;
}
```

the compiler would not allow the assignment `var2.y = 2345` because `var2` has been declared as const. The compiler will allow the assignment `var2.x = 345` because `A::x` has been declared as mutable.

Related information
- “Type qualifiers” on page 68
- “References (C++ only)” on page 88

The register storage class specifier

The register storage class specifier indicates to the compiler that the object should be stored in a machine register. The register storage class specifier is typically specified for heavily used variables, such as a loop control variable, in the hopes of enhancing performance by minimizing access time. However, the compiler is not required to honor this request. Because of the limited size and number of registers available on most systems, few variables can actually be put in registers. If the compiler does not allocate a machine register for a `register` object, the object is treated as having the storage class specifier `auto`.

An object having the register storage class specifier must be defined within a block or declared as a parameter to a function.

The following restrictions apply to the register storage class specifier:
• C You cannot use pointers to reference objects that have the register storage class specifier.

• C You cannot use the register storage class specifier when declaring objects in global scope.

• C A register does not have an address. Therefore, you cannot apply the address operator (&) to a register variable.

• C++ You cannot use the register storage class specifier when declaring objects in namespace scope.

---

C++ only

Unlike C, C++ lets you take the address of an object with the register storage class. For example:

```c
register int i;
int* b = &i;    // valid in C++, but not in C
```

---

End of C++ only

Storage duration of register variables

Objects with the register storage class specifier have automatic storage duration. Each time a block is entered, storage for register objects defined in that block is made available. When the block is exited, the objects are no longer available for use.

If a register object is defined within a function that is recursively invoked, memory is allocated for the variable at each invocation of the block.

Linkage of register variables

Since a register object is treated as the equivalent to an object of the auto storage class, it has no linkage.

Related information

• “Initialization and storage classes” on page 90
• “Block/local scope” on page 2
• “References (C++ only)” on page 88

Global variables in specified registers (C only)

---

IBM extension

You can specify that a particular hardware register is dedicated to a global variable by using an assembly global register variable declaration. Stores into the reserved register are never deleted. This language extension is provided for compatibility with GNU C.

Global register variable declaration syntax

```c
register-variable_declaration asm ("register_specifier")
```

The register_specifier is a string representing a hardware register. The register name is CPU-specific. XL C supports the following register names on the PowerPC:
• r0 to r31 — general purpose registers
• f0 to f31 — floating point registers
• v0 to v31 — vector registers

Note: You cannot use registers r30 or r14 for register variable declarations, as these are reserved for use by the XL C++ runtime environment.

The following are the rules of use for register variables:
• General purpose registers can only be reserved for variables of integer or pointer type.
• Floating point registers can only be reserved for variables of float, double, or 64-bit long double type.
• Vector registers can only be reserved for variables of vector type.
• Variables of long long type cannot reserve registers.
• A global register variable cannot be initialized.
• The register dedicated for a global register variable should not be a volatile register, or the value stored into the global variable might not be preserved across a function call.
• A global register variable can only reserve a register that is not already reserved by another global register variable.
• The same global register variable cannot reserve more than one register.
• A register variable should not be used in an OpenMP clause or OpenMP parallel or work-sharing region.
• The register specified in the global register declaration is reserved for the declared variable only in the compilation unit in which the register declaration is specified. The register is not reserved in other compilation units unless you place the global register declaration in a common header file, or use the -qreserved_reg compiler option.

Related information
• “Assembly labels” on page 17
• “Inline assembly statements” on page 186
• -qreserved_reg in the XL C/C++ Compiler Reference

--- End of IBM extension ---

Type specifiers

Type specifiers indicate the type of the object being declared. The following are the available kinds of type specifiers:
• Fundamental or built-in types:
  — Arithmetic types
    - Integral types
    - Boolean types
    - Floating-point types
    - Character types
  — The void type
  — Vector types
• User-defined types
Integral types

Integral types fall into the following categories:

- Signed integer types:
  - signed char
  - short int
  - int
  - long int
  - long long int
- Unsigned integer types:
  - unsigned char
  - unsigned short int
  - unsigned int
  - unsigned long int
  - unsigned long long int

The unsigned prefix indicates that the object is a nonnegative integer. Each unsigned type provides the same size storage as its signed equivalent. For example, int reserves the same storage as unsigned int. Because a signed type reserves a sign bit, an unsigned type can hold a larger positive integer value than the equivalent signed type.

The declarator for a simple integer definition or declaration is an identifier. You can initialize a simple integer definition with an integer constant or with an expression that evaluates to a value that can be assigned to an integer.

C++ only

When the arguments in overloaded functions and overloaded operators are integer types, two integer types that both come from the same group are not treated as distinct types. For example, you cannot overload an int argument against a signed int argument.

End of C++ only

Related information

- “Integer literals” on page 19
- “Integral conversions” on page 108
- “Arithmetic conversions and promotions” on page 107
- Chapter 10, “Overloading (C++ only),” on page 229

Boolean types

A Boolean variable can be used to hold the integer values 0 or 1, or the literals true or false, which are implicitly promoted to the integers 0 and 1 whenever an arithmetic value is necessary. The Boolean type is unsigned and has the lowest ranking in its category of standard unsigned integer types; it may not be further qualified by the specifiers signed, unsigned, short, or long. In simple assignments, if the left operand is a Boolean type, then the right operand must be either an arithmetic type or a pointer.
Boolean types are a C99 feature. To declare a Boolean variable, use the _Bool type specifier. The token bool is recognized as a keyword in C only when used in a vector declaration context and VMX support is enabled.

To declare a Boolean variable in C++, use the bool type specifier. The result of the equality, relational, and logical operators is of type bool: either of the Boolean constants true or false.

You can use Boolean types make Boolean logic tests. A Boolean logic test is used to express the results of a logical operation. For example:

```c
_Bool f(int a, int b)
{
    return a==b;
}
```

If a and b have the same value, f returns true. If not, f returns false.

Floating-point types

Floating-point type specifiers fall into the following categories:

- Real floating-point types:
  - float
  - double
  - long double

- Complex floating-point types

The magnitude range of float is approximately 1.2e-38 to 3.4e38. The magnitude range of double or long double is approximately 2.2e-308 to 1.8e308. If a floating-point constant is too large or too small, the result is undefined by the language.

The declarator for a simple floating-point declaration is an identifier. Initialize a simple floating-point variable with a float constant or with a variable or expression that evaluates to an integer or floating-point number.

Note: If you do not add the f suffix to a floating-point literal, that number will be of type double. If you initialize an object of type float with an object of type double, the compiler will implicitly convert the object of type double to an object of type float.
Complex floating-point types
The complex type specifiers are:
- float _Complex
- double _Complex
- long double _Complex

The representation and alignment requirements of a complex type are the same as an array type containing two elements of the corresponding real type. The real part is equal to the first element; the imaginary part is equal to the second element.

The equality and inequality operators have the same behavior as for real types. None of the relational operators may have a complex type as an operand.

As an extension to C99 and Standard C++, complex numbers may also be operands to the unary operators ++ (increment), -- (decrement), and ~ (bitwise negation).

Related information
- “Complex literals” on page 24
- “Arithmetic conversions and promotions” on page 107
- “The __real__ and __imag__ operators” on page 144

Character types
Character types fall into the following categories:
- Narrow character types:
  - char
  - signed char
  - unsigned char
- Wide character type wchar_t

The char specifier is an integral type. The wchar_t type specifier is an integral type that has enough storage to represent a wide character literal. (A wide character literal is a character literal that is prefixed with the letter L, for example L'x' )

A char is a distinct type from signed char and unsigned char, and the three types are not compatible.

For the purposes of distinguishing overloaded functions, a C++ char is a distinct type from signed char and unsigned char.

If it does not matter if a char data object is signed or unsigned, you can declare the object as having the data type char. Otherwise, explicitly declare signed char or unsigned char to declare numeric variables that occupy a single byte. When a char (signed or unsigned) is widened to an int, its value is preserved.

By default, char behaves like an signed char. To change this default, you can use the -qchars option or the #pragma chars directive. See #pragma chars and -qchars in the XL C/C++ Compiler Reference for more information.
The void type

The void data type always represents an empty set of values. The only object that can be declared with the type specifier void is a pointer.

You cannot declare a variable of type void, but you can explicitly convert any expression to type void. The resulting expression can only be used as one of the following:
- An expression statement
- The left operand of a comma expression
- The second or third operand in a conditional expression.

Compatibility of arithmetic types (C only)

Two arithmetic types are compatible only if they are the same type.

The presence of type specifiers in various combinations for arithmetic types may or may not indicate different types. For example, the type signed int is the same as int, except when used as the types of bit fields; but char, signed char, and unsigned char are different types.

The presence of a type qualifier changes the type. That is, const int is not the same type as int, and therefore the two types are not compatible.

Vector types

IBM extension

 XL C/C++ supports Vector Multimedia Extension (VMX) technology through language extensions. XL C/C++ implements the AltiVec Programming Interface specification with an extended syntax that allows type qualifiers and storage class specifiers to precede the keyword vector (or its alternate spelling, __vector) in a declaration.

The keyword vector is recognized in a declaration context only when used as a type specifier and when VMX support is enabled. The keywords pixel and bool are recognized as valid type specifiers only when preceded by the keyword vector or __vector. To keep your source code maximally portable, avoid using vector, pixel, or bool as keywords or identifiers in your program. Use the underscore versions of the specifiers vector and pixel (__vector and __pixel) in declarations.

Most of the legal forms of the syntax are captured in the following diagram. Some variations have been omitted from the diagram for the sake of clarity: type
qualifiers such as `const` and storage class specifiers such as `static` can appear in any order within the declaration, as long as neither immediately follows the keyword `vector` (or `_vector`).

**Vector declaration syntax**

```
_type_qualifier_
storage_class_specifier_
```

Notes:
1. The `long` type specifier is deprecated in a vector context.
2. Duplicate type specifiers are ignored in a vector declaration context. In particular, `long long` is treated as `long`.
3. A `long` vector type is compatible with the corresponding `int` vector type (supported in 32-bit mode only).

All vector types are aligned on a 16-byte boundary. An aggregate that contains one or more vector types is aligned on a 16-byte boundary, and padded, if necessary, so that each member of vector type is also 16-byte aligned.

The indirection operator `*` has been extended to handle pointer to vector types. A vector pointer should point to a memory location that has 16-byte alignment. However, the compiler does not enforce this constraint. Dereferencing a vector pointer maintains the vector type and its 16-byte alignment. If a program dereferences a vector pointer that does not contain a 16-byte aligned address, the behavior is undefined.

Pointer arithmetic is defined for pointer to vector types. Given:

```
vector unsigned int *v;
```

the expression `v + 1` represents a pointer to the vector following `v`.

**Related information**
- “Vector literals” on page 25
- “Initialization of vectors” on page 93
- Appendix C, “Vector data types and literals,” on page 415
Vector type casts

Vector types can be cast to other vector types. The cast does not perform a conversion: it preserves the 128-bit pattern, but not necessarily the value. A cast between a vector type and a scalar type is not allowed.

Vector pointers and pointers to non-vector types can be cast back and forth to each other. When a pointer to a non-vector type is cast to a vector pointer, the address should be 16-byte aligned. The referenced object of the pointer to a non-vector type can be aligned on a sixteen-byte boundary by using either the __align specifier or __attribute__((aligned(16))).

Related information
- “The __align qualifier” on page 71
- “The aligned variable attribute” on page 102

User-defined types

The following are user-defined types:
- Structures and unions
- Enumerations
- typedef definitions
- C++ Classes
- C++ Elaborated type specifiers

C++ classes are discussed in Chapter 11, “Classes (C++ only),” on page 245. Elaborated type specifiers are discussed in “Scope of class names” on page 249.

Related information
- “Type attributes” on page 74

Structures and unions

A structure contains an ordered group of data objects. Unlike the elements of an array, the data objects within a structure can have varied data types. Each data object in a structure is a member or field.

A union is an object similar to a structure except that all of its members start at the same location in memory. A union variable can represent the value of only one of its members at a time.

C++ In C++, structures and unions are the same as classes except that their members and inheritance are public by default.

You can declare a structure or union type separately from the definition of variables of that type, as described in “Structure and union type definition” on page 56 and “Structure and union variable declarations” on page 60 or you can define a structure or union data type and all variables that have that type in one statement, as described in “Structure and union type and variable definitions in a single statement” on page 61.
Structures and unions are subject to alignment considerations. For a complete discussion of alignment, see “Aligning data” in the XL C/C++ Programming Guide

Related information
- “Classes and structures” on page 248

Structure and union type definition
A structure or union type definition contains the struct or union keyword followed by an optional identifier (the structure tag) and a brace-enclosed list of members.

Structure or union type definition syntax

```
struct
union

tag_identifier
{
  member_declaration;
}
```

The tag_identifier gives a name to the type. If you do not provide a tag name, you must put all variable definitions that refer to the type within the declaration of the type, as described in “Structure and union type and variable definitions in a single statement” on page 61. Similarly, you cannot use a type qualifier with a structure or union definition; type qualifiers placed in front of the struct or union keyword can only apply to variables that are declared within the type definition.

Related information
- “The aligned type attribute” on page 74
- “The packed type attribute” on page 75

Member declarations
The list of members provides a structure or union data type with a description of the values that can be stored in the structure or union. The definition of a member has the form of a standard variable declaration. The names of member variables must be distinct within a single structure or union, but the same member name may be used in another structure or union type that is defined within the same scope, and may even be the same as a variable, function, or type name.

A structure or union member may be of any type except:
- any variably modified type
- any void type
- a function
- any incomplete type

Because incomplete types are not allowed as members, a structure or union type may not contain an instance of itself as a member, but is allowed to contain a pointer to an instance of itself. As a special case, the last element of a structure with more than one member may have an incomplete array type, which is called a flexible array member, as described in “Flexible array members” on page 57.

As an extension to Standard C and C++ for compatibility with GNU C, XL C/C++ also allows zero-extent arrays as members of structures and unions, as described in “Zero-extent array members” on page 58.
A union member cannot be a class object that has a constructor, destructor, or overloaded copy assignment operator, nor can it be of reference type. A union member cannot be declared with the keyword static.

A member that does not represent a bit field can be qualified with either of the type qualifiers volatile or const. The result is an lvalue.

Structure members are assigned to memory addresses in increasing order, with the first component starting at the beginning address of the structure name itself. To allow proper alignment of components, padding bytes may appear between any consecutive members in the structure layout.

The storage allocated for a union is the storage required for the largest member of the union (plus any padding that is required so that the union will end at a natural boundary of its member having the most stringent requirements). All of a union’s components are effectively overlaid in memory: each member of a union is allocated storage starting at the beginning of the union, and only one member can occupy the storage at a time. For this reason, variably modified types may not be declared as union members.

Flexible array members: A flexible array member, which is a C99 feature, can be an element of a structure with more than one named member. A flexible array member can be used to access a variable-length object. The flexible array member must be the last element of such a structure, and it has incomplete type. It is declared with an empty index, as follows:

```c
array_identifier[ ];
```

For example, `b` is a flexible array member of `Foo`.

```c
struct Foo{
    int a;
    int b[];
};
```

Since a flexible array member has incomplete type, you cannot apply the `sizeof` operator to a flexible array.

Any structure containing a flexible array member cannot be a member of another structure or array.

---

**IBM extension**

For compatibility with GNU C, XL C/C++ extends Standard C and C++, to ease the restrictions on flexible arrays and allow the following:

- Flexible array members can be declared in any part of a structure, not just as the last member.
- Structures containing flexible array members can be members of other structures.
- Flexible array members can be statically initialized.

In the following example:

```c
struct Foo{
    int a;
    int b[];
};

struct foo foo1 = { 55, {6, 8, 10} };
struct foo foo2 = { 55, {15, 6, 14, 90} };
```
foo1 creates an array b of 3 elements, which are initialized to 6, 8, and 10; while foo2 creates an array of 4 elements, which are initialized to 15, 6, 14, and 90.

Flexible array members can only be initialized if they are contained in the outermost part of nested structures. Members of inner structures cannot be initialized.

---

Related information

- “Variable length arrays” on page 87

Zero-extent array members:

---

A zero-extent array is an array with no dimensions. Like a flexible array member, a zero-extent array can be used as a kind of template for array members whose size are determined dynamically at run time. Like a flexible array member, a zero-extent member must be the last element of a structure or union with more than one member. It must explicitly declared with zero as its dimension:

```c
array_identifier[0]
```

It can only be statically initialized with an empty set. For example:

```c
struct foo{
  int a;
  char b[0];
}; bar = { 100, { } };
```

Otherwise, it must be initialized as a dynamically-allocated array.

Zero-extent array members can only be initialized if they are contained in the outermost part of nested structures. Members of inner structures cannot be initialized.

---

Bit field members: Both C and C++ allow integer members to be stored into memory spaces smaller than the compiler would ordinarily allow. These space-saving structure members are called bit fields, and their width in bits can be explicitly declared. Bit fields are used in programs that must force a data structure to correspond to a fixed hardware representation and are unlikely to be portable.

**Bit field member declaration syntax**

```c
<type_specifier><declarator>:<constant_expression>;
```

The constant_expression is a constant integer expression that indicates the field width in bits. A bit field declaration may not use either of the type qualifiers const or volatile.

- In C99, the allowable data types for a bit field include qualified and unqualified _Bool, signed int, and unsigned int. The default integer type for a bit field is signed.
C++ extends the list of allowable types for bit fields to include any integral type or enumeration type.

The maximum bit-field length is 64 bits. For portability, do not use bit fields greater than 32 bits in size.

The following structure has three bit-field members kingdom, phylum, and genus, occupying 12, 6, and 2 bits respectively:

```c
struct taxonomy {
    int kingdom : 12;
    int phylum : 6;
    int genus : 2;
};
```

When you assign a value that is out of range to a bit field, the low-order bit pattern is preserved and the appropriate bits are assigned.

The following restrictions apply to bit fields. You cannot:

- Define an array of bit fields
- Take the address of a bit field
- Have a pointer to a bit field
- Have a reference to a bit field

If a series of bit fields does not add up to the size of an int, padding can take place. The amount of padding is determined by the alignment characteristics of the members of the structure. In some instances, bit fields can cross word boundaries.

Bit fields with a length of 0 must be unnamed. Unnamed bit fields cannot be referenced or initialized.

The following example demonstrates padding, and is valid for all implementations. Suppose that an int occupies 4 bytes. The example declares the identifier kitchen to be of type struct on_off:

```c
struct on_off {
    unsigned light : 1;
    unsigned toaster : 1; /* 4 bytes */
    int count;
    unsigned ac : 4;
    unsigned : 4;  // unnamed field
    unsigned clock : 1;
    unsigned : 0;
    unsigned flag : 1;
} kitchen;
```

The structure kitchen contains eight members totalling 16 bytes. The following table describes the storage that each member occupies:

<table>
<thead>
<tr>
<th>Member name</th>
<th>Storage occupied</th>
</tr>
</thead>
<tbody>
<tr>
<td>light</td>
<td>1 bit</td>
</tr>
<tr>
<td>toaster</td>
<td>1 bit</td>
</tr>
<tr>
<td>(padding — 30 bits)</td>
<td>To the next int boundary</td>
</tr>
<tr>
<td>count</td>
<td>The size of an int (4 bytes)</td>
</tr>
<tr>
<td>ac</td>
<td>4 bits</td>
</tr>
<tr>
<td>(unnamed field)</td>
<td>4 bits</td>
</tr>
<tr>
<td>Member name</td>
<td>Storage occupied</td>
</tr>
<tr>
<td>-------------</td>
<td>------------------</td>
</tr>
<tr>
<td>clock</td>
<td>1 bit</td>
</tr>
<tr>
<td>(padding — 23 bits)</td>
<td>To the next int boundary (unnamed field)</td>
</tr>
<tr>
<td>flag</td>
<td>1 bit</td>
</tr>
<tr>
<td>(padding — 31 bits)</td>
<td>To the next int boundary</td>
</tr>
</tbody>
</table>

Related information
- "Alignment of bit fields" in the **XL C/C++ Programming Guide**

**Structure and union variable declarations**

A structure or union **declaration** has the same form as a definition except the declaration does not have a brace-enclosed list of members. You must declare the structure or union data type before you can define a variable having that type.

**Structure or union variable declaration syntax**

```
storage_class_specifier
type_qualifier
struct
tag_identifier—declarator—;  
union
```

The **tag_identifier** indicates the previously-defined data type of the structure or union.

The keyword struct is optional in structure variable declarations.

You can declare structures or unions having any storage class. The storage class specifier and any type qualifiers for the variable must appear at the beginning of the statement. Structures or unions declared with the register storage class specifier are treated as automatic variables.

The following example defines structure type **address**:

```c
struct address {
    int street_no;
    char *street_name;
    char *city;
    char *prov;
    char *postal_code;
};
```

The following examples declare two structure variables of type **address**:

```c
struct address perm_address;
struct address temp_address;
```

Related information
- "The aligned variable attribute" on page 102
- "The _align qualifier" on page 71
- "The packed variable attribute" on page 103
- "Initialization of structures and unions" on page 94
- "Compatibility of structures, unions, and enumerations (C only)" on page 65
- "Dot operator ." on page 123
Structure and union type and variable definitions in a single statement

You can define a structure or union type and a structure or union variable in one statement, by putting a declarator and an optional initializer after the type definition. The following example defines a union data type (not named) and a union variable (named length):

```
union {
    float meters;
    double centimeters;
    long inches;
} length;
```

Note that because this example does not name the data type, `length` is the only variable that can have this data type. Putting an identifier after `struct` or `union` lets you declare additional variables of this data type later in the program.

To specify a storage class specifier for the variable or variables, you must put the storage class specifier at the beginning of the statement. For example:

```
static struct {
    int street_no;
    char *street_name;
    char *city;
    char *prov;
    char *postal_code;
} perm_address, temp_address;
```

In this case, both `perm_address` and `temp_address` are assigned static storage.

Type qualifiers can be applied to the variable or variables declared in a type definition. Both of the following examples are valid:

```
volatile struct class1 {
    char descriptor[20];
    long code;
    short complete;
} file1, file2;

struct class1 {
    char descriptor[20];
    long code;
    short complete;
} volatile file1, file2;
```

In both cases, the structures `file1` and `file2` are qualified as `volatile`.

Related information

- "Initialization of structures and unions" on page 94
- "Storage class specifiers" on page 43
- "Type qualifiers" on page 68

Access to structure and union members

Once structure or union variables have been declared, members of are referenced by specifying the variable name with the dot operator (.) or a pointer with the arrow operator (->) and the member name. For example, both of the following:

```
perm_address.prov = "Ontario";
p_perm_address -> prov = "Ontario";
```
assign the string "Ontario" to the pointer prov that is in the structure perm_address.

All references to members of structures and unions, including bit fields, must be fully qualified. In the previous example, the fourth field cannot be referenced by prov alone, but only by perm_address.prov.

**Related information**
- “Dot operator .” on page 123
- “Arrow operator ->” on page 123

**Anonymous unions**

An *anonymous union* is a union without a name. It cannot be followed by a declarator. An anonymous union is not a type; it defines an unnamed object.

The member names of an anonymous union must be distinct from other names within the scope in which the union is declared. You can use member names directly in the union scope without any additional member access syntax.

For example, in the following code fragment, you can access the data members i and cptr directly because they are in the scope containing the anonymous union. Because i and cptr are union members and have the same address, you should only use one of them at a time. The assignment to the member cptr will change the value of the member i.

```c
void f()
{
    union { int i; char* cptr; };
    /* . . . */
    i = 5;
    cptr = "string_in_union"; // overrides the value 5
}
```

**C++** An anonymous union cannot have protected or private members, and it cannot have member functions. A global or namespace anonymous union must be declared with the keyword `static`.

**Related information**
- “The static storage class specifier” on page 44
- “Member functions” on page 257

**Enumerations**

An *enumeration* is a data type consisting of a set of named values that represent integral constants, known as *enumeration constants*. An enumeration also referred to as an *enumerated type* because you must list (enumerate) each of the values in creating a name for each of them. In addition to providing a way of defining and grouping sets of integral constants, enumerations are useful for variables that have a small number of possible values.

You can declare an enumeration type separately from the definition of variables of that type, as described in “Enumeration type definition” on page 63 and “Enumeration variable declarations” on page 64 or you can define an enumeration data type and all variables that have that type in one statement, as described in “Enumeration type and variable definitions in a single statement” on page 64.

**Related information**
”Arithmetic conversions and promotions” on page 107

Enumeration type definition
An enumeration type definition contains the enum keyword followed by an optional identifier (the enumeration tag) and a brace-enclosed list of enumerators. A comma separates each enumerator in the enumerator list. C99 allows a trailing comma between the last enumerator and the closing brace. XL C++ also supports this feature, for compatibility with C99.

Enumeration definition syntax

```
enum {enumerator} ;
```

The tag_identifier gives a name to the enumeration type. If you do not provide a tag name, you must put all variable definitions that refer to the enumeration type within the declaration of the type, as described in “Enumeration type and variable definitions in a single statement” on page 64. Similarly, you cannot use a type qualifier with an enumeration definition; type qualifiers placed in front of the enum keyword can only apply to variables that are declared within the type definition.

 Enumeration members: The list of enumeration members, or enumerators, provides the data type with a set of values.

Enumeration member declaration syntax

```
identifier {enumeration_constant}
```

In C, an enumeration constant is of type int. If a constant expression is used as an initializer, the value of the expression cannot exceed the range of int (that is, INT_MIN to INT_MAX as defined in the header limits.h).

In C++, each enumeration constant has a value that can be promoted to a signed or unsigned integer value and a distinct type that does not have to be integral. You can use an enumeration constant anywhere an integer constant is allowed, or anywhere a value of the enumeration type is allowed.

The value of a constant is determined in the following way:
1. An equal sign (=) and a constant expression after the enumeration constant gives an explicit value to the constant. The identifier represents the value of the constant expression.
2. If no explicit value is assigned, the leftmost constant in the list receives the value zero (0).
3. Identifiers with no explicitly assigned values receive the integer value that is one greater than the value represented by the previous identifier.

The following data type declarations list oats, wheat, barley, corn, and rice as enumeration constants. The number under each constant shows the integer value.

```
enum grain { oats, wheat, barley, corn, rice };
/* 0 1 2 3 4 */
```
enum grain { oats=1, wheat, barley, corn, rice; /* 1 2 3 4 5 */
enum grain { oats, wheat=10, barley, corn=20, rice; /* 0 10 11 20 21 */

It is possible to associate the same integer with two different enumeration constants. For example, the following definition is valid. The identifiers suspend and hold have the same integer value.
enum status { run, clear=5, suspend, resume, hold=6; /* 0 5 6 7 6 */

Each enumeration constant must be unique within the scope in which the enumeration is defined. In the following example, the second declarations of average and poor cause compiler errors:
func()
{
    enum score { poor, average, good;  
    enum rating { below, average, above;  
    int poor;
}

Related information
• “Integral types” on page 50

Enumeration variable declarations
You must declare the enumeration data type before you can define a variable having that type.

Enumeration variable declaration syntax

The tag_identifier indicates the previously-defined data type of the enumeration.

> C++ The keyword enum is optional in enumeration variable declarations.

Related information
• “Initialization of enumerations” on page 96
• “Compatibility of structures, unions, and enumerations (C only)” on page 65

Enumeration type and variable definitions in a single statement
You can define a type and a variable in one statement by using a declarator and an optional initializer after the type definition. To specify a storage class specifier for the variable, you must put the storage class specifier at the beginning of the declaration. For example:
register enum score { poor=1, average, good } rating = good;

C++ only

C++ also lets you put the storage class immediately before the declarator list. For example:
enum score { poor=1, average, good } register rating = good;
Either of these examples is equivalent to the following two declarations:

eNum score { poor=1, average, good; }
register enum score rating = good;

Both examples define the enumeration data type score and the variable rating.
rating has the storage class specifier register, the data type enum score, and the
initial value good.

Combining a data type definition with the definitions of all variables having that
data type lets you leave the data type unnamed. For example:

eNum { Sunday, Monday, Tuesday, Wednesday, Thursday, Friday,
Saturday } weekday;

defines the variable weekday, which can be assigned any of the specified
enumeration constants. However, you can not declare any additional enumeration
variables using this set of enumeration constants.

**Compatibility of structures, unions, and enumerations (C only)**

Within a single source file, each structure or union definition creates a new type
that is neither the same as nor compatible with any other structure or union type.
However, a type specifier that is a reference to a previously defined structure or
union type is the same type. The tag associates the reference with the definition,
and effectively acts as the type name. To illustrate this, only the types of structures
j and k are compatible in this example:

```c
struct { int a; int b; } h;
struct { int a; int b; } i;
struct S { int a; int b; } j;
struct S k;
```

Compatible structures may be assigned to each other.

Structures or unions with identical members but different tags are not compatible
and cannot be assigned to each other. Structures and unions with identical
members but using different alignments are not also compatible and cannot be
assigned to each other.

Since the compiler treats enumeration variables and constants as integer types, you
can freely mix the values of different enumerated types, regardless of type
compatibility. Compatibility between an enumerated type and the integer type that
represents it is controlled by compiler options and related pragmas. For a full
discussion of the `-qenum` compiler option and related pragmas, see `-qenum` and
`#pragma enum` in the **XL C/C++ Compiler Reference**.

**Related information**

- “Arithmetic conversions and promotions” on page 107
- Chapter 11, “Classes (C++ only),” on page 245
- “Structure and union type definition” on page 56
- “Incomplete types” on page 40
Compatibility across separate source files

When the definitions for two structures, unions, or enumerations are defined in separate source files, each file can theoretically contain a different definition for an object of that type with the same name. The two declarations must be compatible, or the run time behavior of the program is undefined. Therefore, the compatibility rules are more restrictive and specific than those for compatibility within the same source file. For structure, union, and enumeration types defined in separately compiled files, the composite type is the type in the current source file.

The requirements for compatibility between two structure, union, or enumerated types declared in separate source files are as follows:

- If one is declared with a tag, the other must also be declared with the same tag.
- If both are completed types, their members must correspond exactly in number, be declared with compatible types, and have matching names.

For enumerations, corresponding members must also have the same values.

For structures and unions, the following additional requirements must be met for type compatibility:

- Corresponding members must be declared in the same order (applies to structures only).
- Corresponding bit fields must have the same widths.

typedef definitions

A typedef declaration lets you define your own identifiers that can be used in place of type specifiers such as int, float, and double. A typedef declaration does not reserve storage. The names you define using typedef are not new data types, but synonyms for the data types or combinations of data types they represent.

The namespace for a typedef name is the same as other identifiers. The exception to this rule is if the typedef name specifies a variably modified type. In this case, it has block scope.

When an object is defined using a typedef identifier, the properties of the defined object are exactly the same as if the object were defined by explicitly listing the data type associated with the identifier.

IBM extension

typedef definitions are extended to handle vector types, provided that the VMX support is enabled. A vector type can be used in a typedef definition, and the new type name can be used in the usual ways, except for declaring other vectors. In a vector declaration context, a typedef name is disallowed as a type specifier. The following example illustrates a typical usage of typedef with vector types:

```
typedef vector unsigned short vint16;
  vint16 v1;
```

End of IBM extension

Related information

- “Type names” on page 81
- “Type specifiers” on page 49
Examples of typedef definitions

The following statements define LENGTH as a synonym for int and then use this typedef to declare length, width, and height as integer variables:

```c
typedef int LENGTH;
LENGTH length, width, height;
```

The following declarations are equivalent to the above declaration:

```c
int length, width, height;
```

Similarly, typedef can be used to define a structure, union, or C++ class. For example:

```c
typedef struct {
    int scruples;
    int drams;
    int grains;
} WEIGHT;
```

The structure WEIGHT can then be used in the following declarations:

```c
WEIGHT chicken, cow, horse, whale;
```

In the following example, the type of yds is "pointer to function with no parameter specified, returning int".

```c
typedef int SCROLL();
extern SCROLL *yds;
```

In the following typedefs, the token struct is part of the type name: the type of ex1 is struct a; the type of ex2 is struct b.

```c
typedef struct a { char x; } ex1, *ptr1;
typedef struct b { char x; } ex2, *ptr2;
```

Type ex1 is compatible with the type struct a and the type of the object pointed to by ptr1. Type ex1 is not compatible with char, ex2, or struct b.

---

C++ only

In C++, a typedef name must be different from any class type name declared within the same scope. If the typedef name is the same as a class type name, it can only be so if that typedef is a synonym of the class name. This condition is not the same as in C. The following can be found in standard C headers:

```c
typedef class C { /* data and behavior */ } C;
```

A C++ class defined in a typedef without being named is given a dummy name and the typedef name for linkage. Such a class cannot have constructors or destructors. For example:

```c
typedef class {
    Trees();
} Trees;
```

Here the function Trees(); is an ordinary member function of a class whose type name is unspecified. In the above example, Trees is an alias for the unnamed class, not the class type name itself, so Trees() cannot be a constructor for that class.

---

End of C++ only

Chapter 3. Data objects and declarations  67
Type qualifiers

A type qualifier is used to refine the declaration of a variable, a function, and parameters, by specifying whether:

- The value of an object can be changed
- The value of an object must always be read from memory rather than from a register
- More than one pointer can access a modifiable memory address

XL C/C++ recognizes the following type qualifiers:

- `const`
- `volatile`
- `restrict`
- `__align`

Standard C++ refers to the type qualifiers `const` and `volatile` as *cv-qualifiers*. In both languages, the cv-qualifiers are only meaningful in expressions that are lvalues.

When the `const` and `volatile` keywords are used with pointers, the placement of the qualifier is critical in determining whether it is the pointer itself that is to be qualified, or the object to which the pointer points. For a pointer that you want to qualify as `volatile` or `const`, you must put the keyword between the * and the identifier. For example:

```c
int * volatile x; /* x is a volatile pointer to an int */
int * const y = &z; /* y is a const pointer to the int variable z */
```

For a pointer to a volatile or const data object, the type specifier and qualifier can be in any order, provided that the qualifier does not follow the * operator. For example:

```c
volatile int *x; /* x is a pointer to a volatile int */
int volatile *x; /* x is a pointer to a volatile int */
const int *y; /* y is a pointer to a const int */
int const *y; /* y is a pointer to a const int */
```

The following examples contrast the semantics of these declarations:

<table>
<thead>
<tr>
<th>Declaration</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>const int * ptr1;</code></td>
<td>Defines a pointer to a constant integer: the value pointed to cannot be changed.</td>
</tr>
<tr>
<td><code>int * const ptr2;</code></td>
<td>Defines a constant pointer to an integer: the integer can be changed, but ptr2 cannot point to anything else.</td>
</tr>
<tr>
<td><code>const int * const ptr3;</code></td>
<td>Defines a constant pointer to a constant integer: neither the value pointed to nor the pointer itself can be changed.</td>
</tr>
</tbody>
</table>

You can put more than one qualifier on a declaration: the compiler ignores duplicate type qualifiers.
A type qualifier cannot apply to user-defined types, but only to objects created from a user-defined type. Therefore, the following declaration is illegal:

```c
volatile struct omega {
    int limit;
    char code;
} group;
```

However, if a variable or variables are declared within the same definition of the type, a type qualifier can be applied to the variable or variables by placing at the beginning of the statement or before the variable declarator or declarators. Therefore:

```c
volatile struct omega {
    int limit;
    char code;
} group;
```

provides the same storage as:

```c
struct omega {
    int limit;
    char code;
} volatile group;
```

In both examples, the `volatile` qualifier only applies to the structure variable `group`.

When type qualifiers are applied to a structure, class, or union, or class variable, they also apply to the members of the structure, class or union.

**Related information**
- “Pointers” on page 82
- “Constant and volatile member functions” on page 258

### The `const` type qualifier

The `const` qualifier explicitly declares a data object as something that cannot be changed. Its value is set at initialization. You cannot use `const` data objects in expressions requiring a modifiable lvalue. For example, a `const` data object cannot appear on the lefthand side of an assignment statement.

---

```c
A const object cannot be used in constant expressions. A global const object without an explicit storage class is considered `extern` by default.
```

---

---

```c
In C++, all const declarations must have initializers, except those referencing externally defined constants. A const object can appear in a constant expression if it is an integer and it is initialized to a constant. The following example demonstrates this:
```

```c
const int k = 10;
int ary[k]; // allowed in C++, not legal in C *
```

In C++ a global `const` object without an explicit storage class is considered `static` by default, with internal linkage.
const int k = 12; /* Different meanings in C and C++ */
static const int k2 = 120; /* Same meaning in C and C++ */
extern const int k3 = 121; /* Same meaning in C and C++ */

Because its linkage is assumed to be internal, a const object can be more easily
deefined in header files in C++ than in C.

An item can be both const and volatile. In this case the item cannot be
legitimately modified by its own program but can be modified by some
asynchronous process.

Related information
- "The #define directive” on page 384
- "The this pointer” on page 261

The volatile type qualifier

The volatile qualifier declares a data object that can have its value changed in
ways outside the control or detection of the compiler (such as a variable updated
by the system clock or by another program). This prevents the compiler from
optimizing code referring to the object by storing the object’s value in a register
and re-reading it from there, rather than from memory, where it may have
changed.

Accessing any lvalue expression that is volatile-qualified produces a side effect. A
side effect means that the state of the execution environment changes.

References to an object of type "pointer to volatile" may be optimized, but no
optimization can occur to references to the object to which it points. An explicit
cast must be used to assign a value of type "pointer to volatile T" to an object of
type "pointer to T". The following shows valid uses of volatile objects.

volatile int * pvol;
int *ptr;
pvol = ptr;       /* Legal */
ptr = (int *)pvol; /* Explicit cast required */

> C A signal-handling function may store a value in a variable of type
sig_atomic_t, provided that the variable is declared volatile. This is an exception
to the rule that a signal-handling function may not access variables with static
storage duration.

An item can be both const and volatile. In this case the item cannot be
legitimately modified by its own program but can be modified by some
asynchronous process.

The restrict type qualifier

A pointer is the address of a location in memory. More than one pointer can access
the same chunk of memory and modify it during the course of a program. The
restrict (or __restrict or __restrict__) type qualifier is an indication to the
compiler that, if the memory addressed by the restrict-qualified pointer is
modified, no other pointer will access that same memory. The compiler may
choose to optimize code involving restrict-qualified pointers in a way that might
otherwise result in incorrect behavior. It is the responsibility of the programmer to
ensure that restrict-qualified pointers are used as they were intended to be used. Otherwise, undefined behavior may result.

If a particular chunk of memory is not modified, it can be aliased through more than one restricted pointer, and how an unmodified object can be aliased through two restricted pointers.

```c
void foo(int n, int * restrict a, int * restrict b, int * restrict c) {
    int i;
    for (i = 0; i < n; i++)
        a[i] = b[i] + c[i];
}
```

Assignments between restricted pointers are limited, and no distinction is made between a function call and an equivalent nested block.

```c
{
    int * restrict x;
    int * restrict y;
    x = y; // undefined
    {
        int * restrict x1 = x; // okay
        int * restrict y1 = y; // okay
        x = y1; // undefined
    }
}
```

In nested blocks containing restricted pointers, only assignments of restricted pointers from outer to inner blocks are allowed. The exception is when the block in which the restricted pointer is declared finishes execution. At that point in the program, the value of the restricted pointer can be carried out of the block in which it was declared.

**Notes:**

1. The restrict qualifier is represented by the following keywords (all have the same semantics):
   - C: The `restrict` keyword is only recognized under compilation with `c99` or with the `-qlanglvl=stdc99` or `-qlanglvl=extc99` options (or equivalent pragmas) or `-qkeyword=restrict`. The `__restrict` and `__restrict__` keywords are recognized at all language levels.
   - C++: The `restrict`, `__restrict` and `__restrict__` keywords are recognized at `-qlanglvl=extended`. The `restrict` keyword is also recognized at other language levels with `-qkeyword=restrict`.

**Related information**

- `-qlanglvl` and `-qkeyword` in the XL C/C++ Compiler Reference

**The __align qualifier**

--- IBM extension ---

The __align qualifier is a language extension that allows you to specify an explicit alignment for an aggregate or a static (or global) variable. The specified byte boundary affects the alignment of an aggregate as a whole, not that of its members. The __align qualifier can be applied to an aggregate definition nested
within another aggregate definition, but not to individual elements of an aggregate. The alignment specification is ignored for parameters and automatic variables.

A declaration takes one of the following forms:

__align qualifier syntax for simple variables

```
|type specifier__align__(--int_constant--)declarator|
```

__align qualifier syntax for structures or unions

```
|__align__(--int_constant--)struct[union][tag_identifier]|member_declaration_list->|
```

where int_constant is a positive integer value indicating the byte-alignment boundary. The legal values are 1, 2, 4, 8, or 16.

The following restrictions and limitations apply:
- The __align qualifier cannot be used where the size of the variable alignment is smaller than the size of the type alignment.
- Not all alignments may be representable in an object file.
- The __align qualifier cannot be applied to the following:
  - Individual elements within an aggregate definition.
  - Individual elements of an array.
  - Variables of incomplete type.
  - Aggregates declared but not defined.
  - Other types of declarations or definitions, such as a typedef, a function, or an enumeration.

Examples using the __align qualifier

Applying __align to static or global variables:

```c
int __align(1024) varA; /* varA is aligned on a 1024-byte boundary and padded with 1020 bytes */
main();
static int __align(512) varB; /* varB is aligned on a 512-byte boundary and padded with 508 bytes */
int __align(128) functionB(); /* An error */
typedef int __align(128) T; /* An error */
__align enum C {a, b, c}; /* An error */
```

Applying __align to align and pad aggregate tags without affecting aggregate members:

```
__align(1024) struct structA {int i; int j}; /* struct structA is aligned on a 1024-byte boundary with size including padding of 1024 bytes */
```
Applying __align to a structure or union, where the size and alignment of the aggregate using the structure or union is affected:

```
__align(128) struct S { int i; }; /* sizeof(struct S) == 128 */
struct S sarray[10]; /* sarray is aligned on 128-byte boundary with sizeof(sarray) == 1280 */
struct S __align(64) svar; /* error - alignment of variable is smaller than alignment of type */
struct S2 {struct S s1; int a;} s2; /* s2 is aligned on 128-byte boundary with sizeof(s2) == 256 */
```

Applying __align to an array:

```
AnyType __align(64) arrayA[10]; /* Only arrayA is aligned on a 64-byte boundary, and elements within that array are aligned according to the alignment of AnyType. Padding is applied after the back of the array and does not affect the size of the array member itself. */
```

Applying __align where the size of the variable alignment differs from the size of the type alignment:

```
__align(64) struct S { int i;};
struct S __align(32) s1; /* error, alignment of variable is smaller than alignment of type */
struct S __align(128) s2; /* s2 is aligned on 128-byte boundary */
struct S __align(16) s3[10]; /* error */
int __align(1) s4; /* error */
__align(1) struct S { int i;}; /* error */
```

Related information

- “The aligned variable attribute” on page 102
- “The __alignof__ operator” on page 136
- “Aligning data” in the XL C/C++ Programming Guide
Type attributes

IBM extension

Type attributes are language extensions provided to facilitate compilation of programs developed with the GNU C/C++ compilers. These language features allow you to use named attributes to specify special properties of data objects. Type attributes apply to the definitions of user-defined types, such as structures, unions, enumerations, classes. Any variables that are declared as having that type will have the attribute applied to them.

A type attribute is specified with the keyword __attribute__ followed by the attribute name and any additional arguments the attribute name requires. The __attribute__ specification is included in the definition of a user-defined type, and generally precedes the tag identifier. Although there are variations, the syntax of a type attribute is of the general form:

Type attribute syntax

```
struct __attribute__((attribute name)) {member_definition_list};
```

The attribute name can be specified with or without double underscore characters leading and trailing; however, using the double underscore reduces the likelihood of a name conflict with a macro of the same name. For unsupported attribute names, the XL C/C++ compiler issues diagnostics and ignores the attribute specification. Multiple attribute names can be specified in the same attribute specification.

The following type attributes are supported:

- The aligned type attribute
- The packed type attribute
- The transparent_union type attribute (C only)

Related information

- "Variable attributes" on page 101
- "Function attributes" on page 206

The aligned type attribute

The aligned type attribute allows you to override the default alignment mode to specify a minimum alignment value, expressed as a number of bytes, for a structure, class, union, enumeration, or other user-defined type created in a typedef declaration. The aligned attribute is typically used to increase the alignment of any variables declared of the type to which the attribute applies.

aligned type attribute syntax
The `alignment_factor` is the number of bytes, specified as a constant expression that evaluates to a positive power of 2. You can specify a value up to a maximum 1048576 bytes. If you omit the alignment factor (and its enclosing parentheses), the compiler automatically uses 16 bytes. If you specify an alignment factor greater than the maximum, the attribute specification is ignored, and the compiler simply uses the default alignment in effect.

The alignment value that you specify will be applied to all instances of the type. Also, the alignment value applies to the variable as a whole; if the variable is an aggregate, the alignment value applies to the aggregate as a whole, not to the individual members of the aggregate.

In all of the following examples, the aligned attribute is applied to the structure type A. Because `a` is declared as a variable of type `A`, it will also receive the alignment specification, as will any other instances declared of type `A`.

```
struct __attribute__((__aligned__(8))) A {};
struct __attribute__((__aligned__(8))) A {} a;
typedef struct __attribute__((__aligned__(8))) A {} a;
```

Related information
- “The __align qualifier” on page 71
- “The aligned variable attribute” on page 102
- “The __alignof operator” on page 136
- “Aligning data” in the XL C/C++ Programming Guide

The packed type attribute

The packed type attribute specifies that the minimum alignment should be used for the members of a structure, class, union, or enumeration type. For structure, class, or union types, the alignment is one byte for a member and one bit for a bit field member. For enumeration types, the alignment is the smallest size that will accommodate the range of values in the enumeration. All members of all instances of that type will use the minimum alignment.

**packed type attribute syntax**

```
struct __attribute__((__packed__)) {};
```

Unlike the aligned type attribute, the packed type attribute is not allowed in a typedef declaration.

Related information
- “The __align qualifier” on page 71
- “The packed variable attribute” on page 103
- “The __alignof operator” on page 136
- “Aligning data” in the XL C/C++ Programming Guide
The transparent_union type attribute (C only)

The transparent_union attribute applied to a union definition or a union typedef definition indicates the union can be used as a transparent union. Whenever a transparent union is the type of a function parameter and that function is called, the transparent union can accept an argument of any type that matches that of one of its members without an explicit cast. Arguments to this function parameter are passed to the transparent union, using the calling convention of the first member of the union type. Because of this, all members of the union must have the same machine representation. Transparent unions are useful in library functions that use multiple interfaces to resolve issues of compatibility.

transparent_union type attribute syntax

```
#pragma _attribute_([(transparent_union)])
```

The union must be a complete union type. The transparent_union type attribute can be applied to anonymous unions with tag names.

When the transparent_union type attribute is applied to the outer union of a nested union, the size of the inner union (that is, its largest member) is used to determine if it has the same machine representation as the other members of the outer union. For example,

```c
union __attribute__((transparent_union)) u_t {
  union u2_t {
    char a;
    short b;
    char c;
    char d;
  }
  int a;
};
```

the attribute is ignored because the first member of union u_t, which is itself a union, has a machine representation of 2 bytes, whereas the other member of union u_t is of type int, which has a machine representation of 4 bytes.

The same rationale applies to members of a union that are structures. When a member of a union to which type attribute transparent_union has been applied is a struct, the machine representation of the entire struct is considered, rather than members.

All members of the union must have the same machine representation as the first member of the union. This means that all members must be representable by the same amount memory as the first member of the union. The machine representation of the first member represents the maximum memory size for any remaining union members. For instance, if the first member of a union to which type attribute transparent_union has been applied is of type int, then all following members must be representable by at most 4 bytes. Members that are representable by 1, 2, or 4 bytes are considered valid for this transparent union.

Floating-point types (float, double, float_Complex, or double_Complex) types or vector types can be members of a transparent union, but they cannot be the first member. The restriction that all members of the transparent union have the same
machine representation as the first member still applies.

[--------------------- End of IBM extension ---------------------]
Chapter 4. Declarators

This section continues the discussion of data declarations and includes the following topics:

- “Overview of declarators”
- “Type names” on page 81
- “Pointers” on page 82
- “Arrays” on page 85
- “References (C++ only)” on page 88
- “Initializers” on page 89
- “Variable attributes” on page 101

Overview of declarators

A declarator designates a data object or function. A declarator can also include an initialization. Declarators appear in most data definitions and declarations and in some type definitions.

For data declarations, a declarator has the form:

**Declarator syntax**

```
/\   \                     /\   /
|    \                    /\   /
|      \                  /\   /
|    pointer_operator    /\   /
|                     /\   /
|    direct_declarator  /\   /
|                  /\   /
|    [ initializer ]   /\   /
\_____________________/\   /
```

**Direct declarator:**

```
<table>
<thead>
<tr>
<th>declarator_name</th>
</tr>
</thead>
<tbody>
<tr>
<td>direct_declarator- [constant_expression]</td>
</tr>
</tbody>
</table>
```

**C only**

**Pointer operator:**

```
| * type_qualifiers |
```

**Declarator name:**

```
| identifier |
```

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C++ only

Pointer operator:

```
* type_qualifiers
     & nested_name_specifier
:: type_qualifiers
```

Declarator name:

```
identifier_expression
:: nested_name_specifier
    type_name
```

The type_qualifiers represent one or a combination of const and volatile.

> C++ A nested_name_specifier is a qualified identifier expression. An identifier_expression can be a qualified or unqualified identifier.

Initializers are discussed in “Initializers” on page 89.

The following are known as derived declarator types, and are therefore discussed in this section:

- Points
- Arrays
- References (C++ only)

> IBM In addition, for compatibility with GNU C and C++, XL C/C++ allows you to use variable attributes to modify the properties of data objects. As they are normally specified as part of the declarator in a declaration, they are described in this section, in “Variable attributes” on page 101.

Related information

- “Type qualifiers” on page 68

Examples of declarators

The following table indicates the declarators within the declarations:

<table>
<thead>
<tr>
<th>Declaration</th>
<th>Declarator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int owner;</td>
<td>owner</td>
<td>owner is an integer data object.</td>
</tr>
<tr>
<td>int *node;</td>
<td>*node</td>
<td>node is a pointer to an integer data object.</td>
</tr>
<tr>
<td>int names[126];</td>
<td>names[126]</td>
<td>names is an array of 126 integer elements.</td>
</tr>
<tr>
<td>volatile int min;</td>
<td>min</td>
<td>min is a volatile integer.</td>
</tr>
<tr>
<td>Declaration</td>
<td>Declarator</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>------------</td>
<td>-------------</td>
</tr>
<tr>
<td>int * volatile volume;</td>
<td>* volatile volume</td>
<td>volume is a volatile pointer to an integer.</td>
</tr>
<tr>
<td>volatile int * next;</td>
<td>*next</td>
<td>next is a pointer to a volatile integer.</td>
</tr>
<tr>
<td>volatile int * sequence[5];</td>
<td>*sequence[5]</td>
<td>sequence is an array of five pointers to volatile integer data objects.</td>
</tr>
<tr>
<td>extern const volatile int clock;</td>
<td>clock</td>
<td>clock is a constant and volatile integer with static storage duration and external linkage.</td>
</tr>
</tbody>
</table>

**Related information**

- “Type qualifiers” on page 68
- “Array subscripting operator [ ]” on page 122
- “Scope resolution operator :: (C++ only)” on page 120
- “Function declarators” on page 202

**Type names**

A data type, more precisely, a *type name*, is required in several contexts as something that you must specify without declaring an object; for example, when writing an explicit cast expression or when applying the `sizeof` operator to a type. Syntactically, the name of a data type is the same as a declaration of a function or object of that type, but without the identifier.

To read or write a type name correctly, put an “imaginary” identifier within the syntax, splitting the type name into simpler components. For example, `int` is a type specifier, and it always appears to the left of the identifier in a declaration. An imaginary identifier is unnecessary in this simple case. However, `int *[5]` (an array of 5 pointers to `int`) is also the name of a type. The type specifier `int *` always appears to the left of the identifier, and the array subscripting operator always appears to the right. In this case, an imaginary identifier is helpful in distinguishing the type specifier.

As a general rule, the identifier in a declaration always appears to the left of the subscripting and function call operators, and to the right of a type specifier, type qualifier, or indirection operator. Only the subscripting, function call, and indirection operators may appear in a type name declaration. They bind according to normal operator precedence, which is that the indirection operator is of lower precedence than either the subscripting or function call operators, which have equal ranking in the order of precedence. Parentheses may be used to control the binding of the indirection operator.

It is possible to have a type name within a type name. For example, in a function type, the parameter type syntax nests within the function type name. The same rules of thumb still apply, recursively.

The following constructions illustrate applications of the type naming rules.
Table 14. Type names

<table>
<thead>
<tr>
<th>Syntax</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>int *[5]</td>
<td>array of 5 pointers to int</td>
</tr>
<tr>
<td>int (*)(5]</td>
<td>pointer to an array of 5 integers</td>
</tr>
<tr>
<td>int (<em>)(</em>)</td>
<td>pointer to an variable length array of an unspecified number of integers</td>
</tr>
<tr>
<td>int *()</td>
<td>function with no parameter specification returning a pointer to int</td>
</tr>
<tr>
<td>int (*)(void)</td>
<td>function with no parameters returning an int</td>
</tr>
<tr>
<td>int (*const [])(unsigned int, ...)</td>
<td>array of an unspecified number of constant pointers to functions returning an int. Each function takes one parameter of type unsigned int and an unspecified number of other parameters.</td>
</tr>
</tbody>
</table>

The compiler turns any function designator into a pointer to the function. This behavior simplifies the syntax of function calls.

```c
int foo(float); /* foo is a function designator */
int (*p)(float); /* p is a pointer to a function */
p=&foo;         /* legal, but redundant */
p=foo;          /* legal because the compiler turns foo into a function pointer */
```

In C++, the keywords typename and class, which are interchangeable, indicate the name of the type.

Related information
- “Operator precedence and associativity” on page 162
- “Examples of expressions and precedence” on page 165
- “The typename keyword” on page 359
- “Parenthesized expressions ( )” on page 119

Pointers

A pointer type variable holds the address of a data object or a function. A pointer can refer to an object of any one data type; it cannot refer to a bit field or a reference.

Some common uses for pointers are:
- To access dynamic data structures such as linked lists, trees, and queues.
- To access elements of an array or members of a structure or C++ class.
- To access an array of characters as a string.
- To pass the address of a variable to a function. (In C++, you can also use a reference to do this.) By referencing a variable through its address, a function can change the contents of that variable.

Note that the placement of the type qualifiers volatile and const affects the semantics of a pointer declaration. If either of the qualifiers appears before the *, the declarator describes a pointer to a type-qualified object. If either of the qualifiers appears between the * and the identifier, the declarator describes a type-qualified pointer.
The following table provides examples of pointer declarations.

<table>
<thead>
<tr>
<th>Declaration</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>long *pcoat;</td>
<td>pcoat is a pointer to an object having type long</td>
</tr>
<tr>
<td>extern short * const pvolt;</td>
<td>pvolt is a constant pointer to an object having type short</td>
</tr>
<tr>
<td>extern int volatile *pnut;</td>
<td>pnut is a pointer to an int object having the volatile qualifier</td>
</tr>
<tr>
<td>float * volatile psoup;</td>
<td>psoup is a volatile pointer to an object having type float</td>
</tr>
<tr>
<td>enum bird *pfowl;</td>
<td>pfowl is a pointer to an enumeration object of type bird</td>
</tr>
<tr>
<td>char (*pvish)(void);</td>
<td>pvish is a pointer to a function that takes no parameters and returns a char</td>
</tr>
</tbody>
</table>

**Related information**
- “Type qualifiers” on page 68
- “Initialization of pointers” on page 96
- “Compatibility of pointers (C only)” on page 85
- “Pointer conversions” on page 110
- “Address operator &” on page 134
- “Indirection operator *” on page 136
- “Pointers to functions” on page 219

**Pointer arithmetic**

You can perform a limited number of arithmetic operations on pointers. These operations are:
- Increment and decrement
- Addition and subtraction
- Comparison
- Assignment

The increment (++) operator increases the value of a pointer by the size of the data object the pointer refers to. For example, if the pointer refers to the second element in an array, the ++ makes the pointer refer to the third element in the array.

The decrement (--) operator decreases the value of a pointer by the size of the data object the pointer refers to. For example, if the pointer refers to the second element in an array, the -- makes the pointer refer to the first element in the array.

You can add an integer to a pointer but you cannot add a pointer to a pointer.

If the pointer p points to the first element in an array, the following expression causes the pointer to point to the third element in the same array:

\[ p = p + 2; \]
If you have two pointers that point to the same array, you can subtract one pointer from the other. This operation yields the number of elements in the array that separate the two addresses that the pointers refer to.

You can compare two pointers with the following operators: ==, !=, <, >, <=, and >=.

Pointer comparisons are defined only when the pointers point to elements of the same array. Pointer comparisons using the == and != operators can be performed even when the pointers point to elements of different arrays.

You can assign to a pointer the address of a data object, the value of another compatible pointer or the NULL pointer.

Related information
- “Increment operator ++” on page 132
- “Arrays” on page 85
- “Decrement operator --” on page 133
- Chapter 6, “Expressions and operators,” on page 115

Type-based aliasing
The compiler follows the type-based aliasing rule in the C and C++ standards when the -qalias=ansi option is in effect (which it is by default). This rule, also known as the ANSI aliasing rule, states that a pointer can only be dereferenced to an object of the same type or a compatible type. 1 The common coding practice of casting a pointer to an incompatible type and then dereferencing it violates this rule. (Note that char pointers are an exception to this rule.)

The compiler uses the type-based aliasing information to perform optimizations to the generated code. Contravening the type-based aliasing rule can lead to unexpected behavior, as demonstrated in the following example:

---

1. The C Standard states that an object shall have its stored value accessed only by an lvalue that has one of the following types:
   - the declared type of the object,
   - a qualified version of the declared type of the object,
   - a type that is the signed or unsigned type corresponding to the declared type of the object,
   - a type that is the signed or unsigned type corresponding to a qualified version of the declared type of the object,
   - an aggregate or union type that includes one of the aforementioned types among its members (including, recursively, a member of a subaggregate or contained union), or
   - a character type

The C++ standard states that if a program attempts to access the stored value of an object through an lvalue of other than one of the following types, the behavior is undefined:
   - the dynamic type of the object,
   - a cv-qualified version of the dynamic type of the object,
   - a type that is the signed or unsigned type corresponding to the dynamic type of the object,
   - a type that is the signed or unsigned type corresponding to a cv-qualified version of the dynamic type of the object,
   - an aggregate or union type that includes one of the aforementioned types among its members (including, recursively, a member of a subaggregate or contained union),
   - a type that is a (possible cv-qualified) base class type of the dynamic type of the object,
   - a char or unsigned char type.
int *p;
double d = 0.0;

int *faa(double *g); /* cast operator inside the function */

void foo(double f) {
    p = faa(&f); /* turning &f into a int ptr */
    f += 1.0; /* compiler may discard this statement */
    printf("f=%x\n", *p);
}

int *faa(double *g) { return (int*)g; } /* questionable cast; */
                          /* the function can be in */
                          /* another translation unit */

int main() {
    foo(d);
}

In the above printf statement, *p cannot be dereferenced to a double under the ANSI aliasing rule. The compiler determines that the result of f += 1.0; is never used subsequently. Thus, the optimizer may discard the statement from the generated code. If you compile the above example with optimization enabled, the printf statement may output 0 (zero).

Related information
- “The reinterpret_cast operator (C++ only)” on page 127
- -qalias=ansi in the XL C/C++ Compiler Reference

Compatibility of pointers (C only)

Two pointer types with the same type qualifiers are compatible if they point to objects of compatible types. The composite type for two compatible pointer types is the similarly qualified pointer to the composite type.

The following example shows compatible declarations for the assignment operation:

    float subtotal;
    float *sub_ptr;
    /* ... */
    sub_ptr = &subtotal;
    printf("The subtotal is %f\n", *sub_ptr);

The next example shows incompatible declarations for the assignment operation:

    double league;
    int *minor;
    /* ... */
    minor = &league; /* error */

Arrays

An array is a collection of objects of the same data type, allocated contiguously in memory. Individual objects in an array, called elements, are accessed by their position in the array. The subscripting operator ([ ]) provides the mechanics for creating an index to array elements. This form of access is called indexing or subscripting. An array facilitates the coding of repetitive tasks by allowing the statements executed on each element to be put into a loop that iterates through each element in the array.
The C and C++ languages provide limited built-in support for an array type: reading and writing individual elements. Assignment of one array to another, the comparison of two arrays for equality, returning self-knowledge of size are not supported by either language.

The type of an array is derived from the type of its elements, in what is called array type derivation. If array objects are of incomplete type, the array type is also considered incomplete. Array elements may not be of type void or of function type. However, arrays of pointers to functions are allowed. Array elements may not be of reference type or of an abstract class type.

The array declarator contains an identifier followed by an optional subscript declarator. An identifier preceded by an asterisk (*) is an array of pointers.

**Array subscript declarator syntax**

```
constant_expression
```

The constant_expression is a constant integer expression, indicating the size of the array, which must be positive.

If the declaration appears in block or function scope, a nonconstant expression can be specified for the array subscript declarator, and the array is considered a variable-length array, as described in “Variable length arrays” on page 87.

The subscript declarator describes the number of dimensions in the array and the number of elements in each dimension. Each bracketed expression, or subscript, describes a different dimension and must be a constant expression.

The following example defines a one-dimensional array that contains four elements having type char:

```c
char
list[4];
```

The first subscript of each dimension is 0. The array list contains the elements:

`list[0]`  
`list[1]`  
`list[2]`  
`list[3]`

The following example defines a two-dimensional array that contains six elements of type int:

```c
int
roster[3][2];
```

Multidimensional arrays are stored in row-major order. When elements are referred to in order of increasing storage location, the last subscript varies the fastest. For example, the elements of array roster are stored in the order:

`roster[0][0]`  
`roster[0][1]`  
`roster[1][0]`  
`roster[1][1]`  
`roster[2][0]`  
`roster[2][1]`
In storage, the elements of `roster` would be stored as:

\[
\begin{array}{ccc}
\hline
\text{roster}[0][0] & \text{roster}[0][1] & \text{roster}[1][0] \\
\hline
\end{array}
\]

You can leave the first (and only the first) set of subscript brackets empty in:
- Array definitions that contain initializations
- `extern` declarations
- Parameter declarations

In array definitions that leave the first set of subscript brackets empty, the initializer determines the number of elements in the first dimension. In a one-dimensional array, the number of initialized elements becomes the total number of elements. In a multidimensional array, the initializer is compared to the subscript declarator to determine the number of elements in the first dimension.

Related information
- “Array subscriptsing operator [ ]” on page 122
- “Initialization of arrays” on page 97

**Variable length arrays**

A variable length array, which is a C99 feature, is an array of automatic storage duration whose length is determined at run time.

**Variable length array declarator syntax**

\[
\text{array_identifier}[: \text{expression} : \text{type-qualifiers}] \\
\]

If the size of the array is indicated by `*` instead of an expression, the variable length array is considered to be of unspecified size. Such arrays are considered complete types, but can only be used in declarations of function prototype scope.

A variable length array and a pointer to a variable length array are considered *variably modified types*. Declarations of variably modified types must be at either block scope or function prototype scope. Array objects declared with the `extern` storage class specifier cannot be of variable length array type. Array objects declared with the `static` storage class specifier can be a pointer to a variable length array, but not an actual variable length array. The identifiers declared with a variably modified type must be ordinary identifiers and therefore cannot be members of structures or unions. A variable length array cannot be initialized.

A variable length array can be the operand of a `sizeof` expression. In this case, the operand is evaluated at run time, and the size is neither an integer constant nor a constant expression, even though the size of each instance of a variable array does not change during its lifetime.
A variable length array can be used in a typedef expression. The typedef name will have only block scope. The length of the array is fixed when the typedef name is defined, not each time it is used.

A function parameter can be a variable length array. The necessary size expressions must be provided in the function definition. The compiler evaluates the size expression of a variably modified parameter on entry to the function. For a function declared with a variable length array as a parameter, as in the following,

```c
void f(int x, int a[][x]);
```

the size of the variable length array argument must match that of the function definition.

**IBM C++** The C++ extension does not include support for references to a variable length array type; neither may a function parameter be a reference to a variable length array type.

**Related information**
- “Flexible array members” on page 57

**Compatibility of arrays**

Two array types that are similarly qualified are compatible if the types of their elements are compatible. For example,

```c
char ex1[25];
const char ex2[25];
```

are not compatible.

The composite type of two compatible array types is an array with the composite element type. The sizes of both original types must be equivalent if they are known. If the size of only one of the original array types is known, then the composite type has that size. For example:

```c
char ex3[];
char ex4[42];
```

The composite type of `ex3` and `ex4` is `char[42]`. If one of the original types is a variable length array, the composite type is that type.

**Related information**
- “External linkage” on page 9

**References (C++ only)**

A reference is an alias or an alternative name for an object. All operations applied to a reference act on the object to which the reference refers. The address of a reference is the address of the aliased object.

A reference type is defined by placing the reference modifier & after the type specifier. You must initialize all references except function parameters when they are defined. Once defined, a reference cannot be reassigned because it is an alias to its target. What happens when you try to reassign a reference turns out to be the assignment of a new value to the target.
Because arguments of a function are passed by value, a function call does not modify the actual values of the arguments. If a function needs to modify the actual value of an argument or needs to return more than one value, the argument must be passed by reference (as opposed to being passed by value). Passing arguments by reference can be done using either references or pointers. Unlike C, C++ does not force you to use pointers if you want to pass arguments by reference. The syntax of using a reference is somewhat simpler than that of using a pointer. Passing an object by reference enables the function to change the object being referred to without creating a copy of the object within the scope of the function. Only the address of the actual original object is put on the stack, not the entire object.

For example:
```cpp
int f(int&);  
int main()  
{  
    extern int i;  
    f(i);  
}  
```

You cannot tell from the function call f(i) that the argument is being passed by reference.

References to NULL are not allowed.

**Related information**

- [“Initialization of references (C++ only)” on page 100](#)
- [“Pointers” on page 82](#)
- [“Reference conversions (C++ only)” on page 112](#)
- [“Address operator &” on page 134](#)
- [“Pass by reference” on page 215](#)

### Initializers

An *initializer* is an optional part of a data declaration that specifies an initial value of a data object. The initializers that are legal for a particular declaration depend on the type and storage class of the object to be initialized.

The initializer consists of the = symbol followed by an initial *expression* or a brace-enclosed list of initial expressions separated by commas. Individual expressions must be separated by commas, and groups of expressions can be enclosed in braces and separated by commas. Braces ({ }) are optional if the initializer for a character string is a string literal. The number of initializers must not be greater than the number of elements to be initialized. The initial expression evaluates to the first value of the data object.

To assign a value to an arithmetic or pointer type, use the simple initializer: = *expression*. For example, the following data definition uses the initializer = 3 to set the initial value of group to 3:

```cpp
int group = 3;  
```

You initialize a variable of character type with a character literal (consisting of one character) or with an expression that evaluates to an integer.
You can initialize variables at namespace scope with nonconstant expressions. You cannot initialize variables at global scope with nonconstant expressions.

"Initialization and storage classes" discusses the rules for initialization according to the storage class of variables.

"Designated initializers for aggregate types (C only)" on page 91 describes designated initializers, which are a C99 feature that can be used to initialize arrays, structures, and unions.

The following sections discuss initializations for derived types:
- "Initialization of vectors" on page 93
- "Initialization of structures and unions" on page 94
- "Initialization of pointers" on page 96
- "Initialization of arrays" on page 97
- "Initialization of references (C++ only)" on page 100

Related information
- "Using class objects" on page 246

**Initialization and storage classes**

**Initialization of automatic variables**
You can initialize any auto variable except function parameters. If you do not explicitly initialize an automatic object, its value is indeterminate. If you provide an initial value, the expression representing the initial value can be any valid C or C++ expression. The object is then set to that initial value each time the program block that contains the object's definition is entered.

Note that if you use the goto statement to jump into the middle of a block, automatic variables within that block are not initialized.

Related information
- "The auto storage class specifier" on page 44

**Initialization of static variables**
You initialize a static object with a constant expression, or an expression that reduces to the address of a previously declared extern or static object, possibly modified by a constant expression. If you do not explicitly initialize a static (or external) variable, it will have a value of zero of the appropriate type, unless it is a pointer, in which case it will be initialized to NULL.

A static variable in a block is initialized only one time, prior to program execution, whereas an auto variable that has an initializer is initialized every time it comes into existence.

A static object of class type will use the default constructor if you do not initialize it. Automatic and register variables that are not initialized will have undefined values.

Related information
- "The static storage class specifier" on page 44
Initialization of external variables
You can initialize any object with the extern storage class specifier at global scope in C or at namespace scope in C++. The initializer for an extern object must either:
- Appear as part of the definition and the initial value must be described by a constant expression; or
- Reduce to the address of a previously declared object with static storage duration. You may modify this object with pointer arithmetic. (In other words, you may modify the object by adding or subtracting an integral constant expression.)

If you do not explicitly initialize an extern variable, its initial value is zero of the appropriate type. Initialization of an extern object is completed by the time the program starts running.

Related information
- “The extern storage class specifier” on page 46

Initialization of register variables
You can initialize any register object except function parameters. If you do not initialize an automatic object, its value is indeterminate. If you provide an initial value, the expression representing the initial value can be any valid C or C++ expression. The object is then set to that initial value each time the program block that contains the object’s definition is entered.

Related information
- “The register storage class specifier” on page 47

Designated initializers for aggregate types (C only)
Designated initializers, a C99 feature, are supported for aggregate types, including arrays, structures, and unions. A designated initializer, or designator, points out a particular element to be initialized. A designator list is a comma-separated list of one or more designators. A designator list followed by an equal sign constitutes a designation.

Designated initializers allow for the following flexibility:
- Elements within an aggregate can be initialized in any order.
- The initializer list can omit elements that are declared anywhere in the aggregate, rather than only at the end. Elements that are omitted are initialized as if they are static objects: arithmetic types are initialized to 0; pointers are initialized to NULL.
- Where inconsistent or incomplete bracketing of initializers for multi-dimensional arrays or nested aggregates may be difficult to understand, designators can more clearly identify the element or member to be initialized.

Designator list syntax for structures and unions

```
{<member>=<expression>}
```
Designator list syntax for arrays

In the following example, the designator is .any_member and the designated initializer is .any_member = 13:
union { /* ... */ } caw = { .any_member = 13 };

The following example shows how the second and third members b and c of structure variable klm are initialized with designated initializers:
struct xyz {
    int a;
    int b;
    int c;
} klm = { .a = 99, .c = 100 };

In the following example, the third and second elements of the one-dimensional array aa are initialized to 3 and 6, respectively:

The following example initializes the first four and last four elements, while omitting the middle four:

The omitted four elements of grid are initialized to zero:

<table>
<thead>
<tr>
<th>Element</th>
<th>Value</th>
<th>Element</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>grid[0] [0]</td>
<td>8</td>
<td>grid[1] [2]</td>
<td>0</td>
</tr>
<tr>
<td>grid[0] [1]</td>
<td>6</td>
<td>grid[1] [3]</td>
<td>0</td>
</tr>
<tr>
<td>grid[0] [2]</td>
<td>4</td>
<td>grid[2] [0]</td>
<td>9</td>
</tr>
<tr>
<td>grid[1] [0]</td>
<td>0</td>
<td>grid[2] [2]</td>
<td>1</td>
</tr>
</tbody>
</table>

Designated initializers can be combined with regular initializers, as in the following example:
int a[10] = {2, 4, [8]=9, 10}

In this example, a[0] is initialized to 2, a[1] is initialized to 4, a[2] to a[7] are initialized to 0, and a[9] is initialized to 10.

In the following example, a single designator is used to "allocate" space from both ends of an array:
int a[MAX] = {
    1, 3, 5, 7, 9, [MAX-5] = 8, 6, 4, 2, 0
};

The designated initializer, [MAX-5] = 8, means that the array element at subscript MAX-5 should be initialized to the value 8. If MAX is 15, a[5] through a[9] will be initialized to zero. If MAX is 7, a[2] through a[4] will first have the values 5, 7, and 9, respectively, which are overridden by the values 8, 6, and 4. In other words, if MAX is 7, the initialization would be the same as if the declaration had been written:

int a[MAX] = {
    1, 3, 8, 6, 4, 2, 0
};

You can also use designators to represent members of nested structures. For example:

struct a {
    struct b {
        int c;
        int d;
    } e;
    float f;
} g = {.e.c = 3};

initializes member c of structure variable e, which is a member of structure variable g, to the value of 3.

Related information
- "Initialization of structures and unions" on page 94
- "Initialization of arrays" on page 97

Initialization of vectors

IBM extension

A vector type is initialized by a vector literal or any expression having the same vector type. For example:

vector unsigned int v1;
vector unsigned int v2 = (vector unsigned int){10};
v1 = v2;

XL C/C++ extends the AltiVec specification to allow a vector type to be initialized by an initializer list. This feature is an extension for compatibility with GNU C.

Vector initializer list syntax

```
vector_type identifier = { initializer };
```

The number of values in a braced initializer list must be less than or equal to the number of elements of the vector type. Any uninitialized element will be initialized to zero.

The following are examples of vector initialization using initializer lists:

vector unsigned int v1 = {1}; // initialize the first 4 bytes of v1 with 1
// and the remaining 12 bytes with zeros

vector unsigned int v2 = {1,2}; // initialize the first 8 bytes of v2 with 1 and 2
vector unsigned int v3 = {1,2,3,4}; // equivalent to the vector literal (vector unsigned int) {1,2,3,4}

Unlike vector literals, the values in the initializer list do not have to be constant expressions unless the initialized vector variable has static duration. Thus, the following is legal:

```c
int i=1;
int foo() { return 2; }
int main()
{
    vector unsigned int v1 = {i, foo()};
    return 0;
}
```

Initialization of structures and unions

An initializer for a structure is a brace-enclosed comma-separated list of values, and for a union, a brace-enclosed single value. The initializer is preceded by an equal sign (=).

C99 and C++ allow the initializer for an automatic member variable of a union or structure type to be a constant or non-constant expression.

The initializer for a static member variable of a union or structure type must be a constant expression or string literal. See “Static data members” on page 266 for more information.

There are two ways to specify initializers for structures and unions:

- With C89-style initializers, structure members must be initialized in the order declared, and only the first member of a union can be initialized.

- Using designated initializers, a C99 feature which allows you to name members to be initialized, structure members can be initialized in any order, and any (single) member of a union can be initialized. Designated initializers are described in detail in “Designated initializers for aggregate types (C only)” on page 91.

Using C89-style initialization, the following example shows how you would initialize the first union member birthday of the union variable people:

```c
union {
    char birthday[9];
    int age;
    float weight;
} people = {"23/07/57"};
```

Using a designated initializer in the same example, the following initializes the second union member age:

```c
union {
    char birthday[9];
    int age;
    float weight;
} people = {.age = 14};
```
The following definition shows a completely initialized structure:

```c
struct address {
    int street_no;
    char *street_name;
    char *city;
    char *prov;
    char *postal_code;
};
```

```c
static struct address perm_address = {
    3, "Savona Dr.", "Dundas", "Ontario", "L4B 2A1"};
```

The values of `perm_address` are:

<table>
<thead>
<tr>
<th>Member</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>perm_address.street_no</td>
<td>3</td>
</tr>
<tr>
<td>perm_address.street_name</td>
<td>address of string &quot;Savona Dr.&quot;</td>
</tr>
<tr>
<td>perm_address.city</td>
<td>address of string &quot;Dundas&quot;</td>
</tr>
<tr>
<td>perm_address.prov</td>
<td>address of string &quot;Ontario&quot;</td>
</tr>
<tr>
<td>perm_address.postal_code</td>
<td>address of string &quot;L4B 2A1&quot;</td>
</tr>
</tbody>
</table>

Unnamed structure or union members do not participate in initialization and have indeterminate value after initialization. Therefore, in the following example, the bit field is not initialized, and the initializer 3 is applied to member `b`:

```c
struct {
    int a;
    int :10;  // bit field
    int b;
} w = { 2, 3 };  // initializer 3 applied to b
```

You do not have to initialize all members of a structure or union; the initial value of uninitialized structure members depends on the storage class associated with the structure or union variable. In a structure declared as static, any members that are not initialized are implicitly initialized to zero of the appropriate type; the members of a structure with automatic storage have no default initialization. The default initializer for a union with static storage is the default for the first component; a union with automatic storage has no default initialization.

The following definition shows a partially initialized structure:

```c
struct address {
    int street_no;
    char *street_name;
    char *city;
    char *prov;
    char *postal_code;
};
```

```c
struct address temp_address = {
    44, "Knyvet Ave.", "Hamilton", "Ontario" };  // initializer 44 applied to street_no
```

The values of `temp_address` are:

<table>
<thead>
<tr>
<th>Member</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>temp_address.street_no</td>
<td>44</td>
</tr>
<tr>
<td>temp_address.street_name</td>
<td>address of string &quot;Knyvet Ave.&quot;</td>
</tr>
</tbody>
</table>

Chapter 4. Declarators 95
<table>
<thead>
<tr>
<th>Member</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>temp_address.city</code></td>
<td>address of string &quot;Hamilton&quot;</td>
</tr>
<tr>
<td><code>temp_address.prov</code></td>
<td>address of string &quot;Ontario&quot;</td>
</tr>
<tr>
<td><code>temp_address.postal_code</code></td>
<td>Depends on the storage class of the <code>temp_address</code> variable; if it is static, the value would be NULL.</td>
</tr>
</tbody>
</table>

### C only

To initialize only the third and fourth members of the `temp_address` variable, you could use a designated initializer list, as follows:

```c
struct address {
    int street_no;
    char *street_name;
    char *city;
    char *prov;
    char *postal_code;
};
struct address temp_address =
{ .city = "Hamilton", .prov = "Ontario" };
```

### Related information
- [“Structure and union variable declarations” on page 60](#)
- [“Explicit initialization with constructors” on page 306](#)
- [“Assignment expressions” on page 158](#)

### Initialization of enumerations

The initializer for an enumeration variable contains the `=` symbol followed by an expression `enumeration_constant`.

```c++
// In C++, the initializer must have the same type as the associated enumeration type.
enum grain { oats, wheat, barley, corn, rice };
enum grain g_food = barley;
```

### Related information
- [“Enumeration variable declarations” on page 64](#)

### Initialization of pointers

The initializer is an `=` (equal sign) followed by the expression that represents the address that the pointer is to contain. The following example defines the variables `time` and `speed` as having type `double` and `amount` as having type `pointer` to a `double`. The pointer `amount` is initialized to point to `total`:

```c
double total, speed, *amount = &total;
```

The compiler converts an unsubscripted array name to a pointer to the first element in the array. You can assign the address of the first element of an array to
a pointer by specifying the name of the array. The following two sets of definitions are equivalent. Both define the pointer student and initialize student to the address of the first element in section:

```c
int section[80];
int *student = section;
```

is equivalent to:

```c
int section[80];
int *student = &section[0];
```

You can assign the address of the first character in a string constant to a pointer by specifying the string constant in the initializer. The following example defines the pointer variable string and the string constant "abcd". The pointer string is initialized to point to the character a in the string "abcd".

```c
char *string = "abcd";
```

The following example defines weekdays as an array of pointers to string constants. Each element points to a different string. The pointer weekdays[2], for example, points to the string "Tuesday".

```c
static char *weekdays[ ] =
{
"Sunday", "Monday", "Tuesday", "Wednesday",
"Thursday", "Friday", "Saturday"
};
```

A pointer can also be initialized to null using any integer constant expression that evaluates to 0, for example `char * a = 0;`. Such a pointer is a null pointer. It does not point to any object.

**Related information**
- "Pointers" on page 82

**Initialization of arrays**

The initializer for an array is a comma-separated list of constant expressions enclosed in braces (`{  }`). The initializer is preceded by an equal sign (`=`). You do not need to initialize all elements in an array. If an array is partially initialized, elements that are not initialized receive the value 0 of the appropriate type. The same applies to elements of arrays with static storage duration. (All file-scope variables and function-scope variables declared with the static keyword have static storage duration.)

There are two ways to specify initializers for arrays:
- With C89-style initializers, array elements must be initialized in subscript order.
- Using designated initializers, which allow you to specify the values of the subscript elements to be initialized, array elements can be initialized in any order. Designated initializers are described in detail in "Designated initializers for aggregate types (C only)" on page 91.

Using C89-style initializers, the following definition shows a completely initialized one-dimensional array:

```c
static int number[3] = { 5, 7, 2 };
```
The array number contains the following values: `number[0]` is 5, `number[1]` is 7; `number[2]` is 2. When you have an expression in the subscript declarator defining the number of elements (in this case 3), you cannot have more initializers than the number of elements in the array.

The following definition shows a partially initialized one-dimensional array:
```c
static int number1[3] = { 5, 7 };
```

The values of `number1[0]` and `number1[1]` are the same as in the previous definition, but `number1[2]` is 0.

```
C only

The following definition shows how you can use designated initializers to skip over elements of the array that you don’t want to initialize explicitly:
```c
```

The array number contains the following values: `number[0]` is 5; `number[1]` is implicitly initialized to 0; `number[2]` is 7.

```
End of C only
```

Instead of an expression in the subscript declarator defining the number of elements, the following one-dimensional array definition defines one element for each initializer specified:
```c
static int item[ ] = { 1, 2, 3, 4, 5 };
```

The compiler gives `item` the five initialized elements, because no size was specified and there are five initializers.

**Initialization of character arrays**

You can initialize a one-dimensional character array by specifying:

- A brace-enclosed comma-separated list of constants, each of which can be contained in a character
- A string constant (braces surrounding the constant are optional)

Initializing a string constant places the null character (\0) at the end of the string if there is room or if the array dimensions are not specified.

The following definitions show character array initializations:
```c
static char name1[ ] = { 'J', 'a', 'n' };
static char name2[ ] = { "Jan" };
static char name3[4] = "Jan";
```

These definitions create the following elements:

<table>
<thead>
<tr>
<th>Element</th>
<th>Value</th>
<th>Element</th>
<th>Value</th>
<th>Element</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>name1[0]</code></td>
<td>J</td>
<td><code>name2[0]</code></td>
<td>J</td>
<td><code>name3[0]</code></td>
<td>J</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>name2[3]</code></td>
<td>\0</td>
<td><code>name3[3]</code></td>
<td>\0</td>
</tr>
</tbody>
</table>
Note that the following definition would result in the null character being lost:

```c
static char name3[3]="Jan";
```

> **C++** When you initialize an array of characters with a string, the number of characters in the string — including the terminating '\0' — must not exceed the number of elements in the array.

### Initialization of multidimensional arrays

You can initialize a multidimensional array using any of the following techniques:

- Listing the values of all elements you want to initialize, in the order that the compiler assigns the values. The compiler assigns values by increasing the subscript of the last dimension fastest. This form of a multidimensional array initialization looks like a one-dimensional array initialization. The following definition completely initializes the array `month_days`:

  ```c
  static month_days[2][12] =
  {
    31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31, 31,
    31, 29, 31, 30, 31, 30, 31, 30, 31, 30, 31, 31
  };
  ```

- Using braces to group the values of the elements you want initialized. You can put braces around each element, or around any nesting level of elements. The following definition contains two elements in the first dimension (you can consider these elements as rows). The initialization contains braces around each of these two elements:

  ```c
  static int month_days[2][12] =
  {
    { 31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31, 31 },
    { 31, 29, 31, 30, 31, 30, 31, 30, 31, 30, 31, 31 }
  };
  ```

- Using nested braces to initialize dimensions and elements in a dimension selectively. In the following example, only the first eight elements of the array `grid` are explicitly initialized. The remaining four elements that are not explicitly initialized are automatically initialized to zero.

  ```c
  static short grid[3][4] = {8, 6, 4, 1, 9, 3, 1, 1};
  ```

The initial values of `grid` are:

<table>
<thead>
<tr>
<th>Element</th>
<th>Value</th>
<th>Element</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>grid[0][0]</td>
<td>8</td>
<td>grid[1][2]</td>
<td>1</td>
</tr>
<tr>
<td>grid[0][1]</td>
<td>6</td>
<td>grid[1][3]</td>
<td>1</td>
</tr>
<tr>
<td>grid[0][2]</td>
<td>4</td>
<td>grid[2][0]</td>
<td>0</td>
</tr>
<tr>
<td>grid[0][3]</td>
<td>1</td>
<td>grid[2][1]</td>
<td>0</td>
</tr>
<tr>
<td>grid[1][0]</td>
<td>9</td>
<td>grid[2][2]</td>
<td>0</td>
</tr>
<tr>
<td>grid[1][1]</td>
<td>3</td>
<td>grid[2][3]</td>
<td>0</td>
</tr>
</tbody>
</table>

> **C only**

- Using designated initializers. The following example uses designated initializers to explicitly initialize only the last four elements of the array. The first eight elements that are not explicitly initialized are automatically initialized to zero.
static short grid[3][4] = { grid[0][0] = 8, grid[0][1] = 6, grid[0][2] = 4, grid[0][3] = 1};

The initial values of grid are:

<table>
<thead>
<tr>
<th>Element</th>
<th>Value</th>
<th>Element</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>grid[0][0]</td>
<td>0</td>
<td>grid[1][2]</td>
<td>0</td>
</tr>
<tr>
<td>grid[0][1]</td>
<td>0</td>
<td>grid[1][3]</td>
<td>0</td>
</tr>
<tr>
<td>grid[0][2]</td>
<td>0</td>
<td>grid[2][0]</td>
<td>8</td>
</tr>
<tr>
<td>grid[0][3]</td>
<td>0</td>
<td>grid[2][1]</td>
<td>6</td>
</tr>
<tr>
<td>grid[1][0]</td>
<td>0</td>
<td>grid[2][2]</td>
<td>4</td>
</tr>
<tr>
<td>grid[1][1]</td>
<td>0</td>
<td>grid[2][3]</td>
<td>1</td>
</tr>
</tbody>
</table>

Related information

- “Arrays” on page 85
- “Designated initializers for aggregate types (C only)” on page 91

Initialization of references (C++ only)

The object that you use to initialize a reference must be of the same type as the reference, or it must be of a type that is convertible to the reference type. If you initialize a reference to a constant using an object that requires conversion, a temporary object is created. In the following example, a temporary object of type float is created:

```cpp
int i;
const float& f = i; // reference to a constant float
```

When you initialize a reference with an object, you bind that reference to that object.

Attempting to initialize a nonconstant reference with an object that requires a conversion is an error.

Once a reference has been initialized, it cannot be modified to refer to another object. For example:

```cpp
int num1 = 10;
int num2 = 20;

int &RefOne = num1; // valid
int &RefOne = num2; // error, two definitions of RefOne
RefOne = num2; // assign num2 to num1
int &RefTwo = num2; // error, uninitialized reference
int &RefTwo = num2; // valid
```

Note that the initialization of a reference is not the same as an assignment to a reference. Initialization operates on the actual reference by initializing the reference with the object it is an alias for. Assignment operates through the reference on the object referred to.

A reference can be declared without an initializer:
• When it is used in an parameter declaration
• In the declaration of a return type for a function call
• In the declaration of class member within its class declaration
• When the extern specifier is explicitly used

You cannot have references to any of the following:
• Other references
• Bit fields
• Arrays of references
• Pointers to references

**Direct binding**
Suppose a reference \( r \) of type \( T \) is initialized by an expression \( e \) of type \( U \).

The reference \( r \) is *bound directly* to \( e \) if the following statements are true:
• Expression \( e \) is an lvalue
• \( T \) is the same type as \( U \), or \( T \) is a base class of \( U \)
• \( T \) has the same, or more, \texttt{const} or \texttt{volatile} qualifiers than \( U \)

The reference \( r \) is also bound directly to \( e \) if \( e \) can be implicitly converted to a type such that the previous list of statements is true.

**Related information**
• “References (C++ only)” on page 88
• “Pass by reference” on page 215

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**Variable attributes**

---

**IBM extension**

Variable attributes are language extensions provided to facilitate the compilation of programs developed with the GNU C/C++ compilers. These language features allow you to use named attributes to specify special properties of data objects. Variable attributes apply to the declarations of simple variables, aggregates, and member variables of aggregates.

A variable attribute is specified with the keyword \texttt{__attribute__} followed by the attribute name and any additional arguments the attribute name requires. A variable \texttt{__attribute__} specification is included in the declaration of a variable, and can be placed before or after the declarator. Although there are variations, the syntax generally takes either of the following forms:

**Variable attribute syntax: post-declarator**

```
-declarator__attribute__((attribute name))
```
Variable attribute syntax: pre-declarator

```
(type specifier) __attribute__ (attribute name) ...
```

The `attribute name` can be specified with or without leading and trailing double underscore characters; however, using the double underscore reduces the likelihood of a name conflict with a macro of the same name. For unsupported attribute names, the XL C/C++ compiler issues diagnostics and ignores the attribute specification. Multiple attribute names can be specified in the same attribute specification.

In a comma-separated list of declarators on a single declaration line, if a variable attribute appears before all the declarators, it applies to all declarators in the declaration. If the attribute appears after a declarator, it only applies to the immediately preceding declarator. For example:

```c
struct A {
    int b __attribute__((aligned)); /* typical placement of variable attribute */
    int __attribute__((aligned)) c = 10; /* variable attribute can also be placed here */
    int d, e, f __attribute__((aligned)); /* attribute applies to f only */
    int g __attribute__((aligned)), h, i; /* attribute applies to g only */
    int __attribute__((aligned)) j, k, l; /* attribute applies to j, k, and l */
};
```

The following variable attributes are supported:

- The aligned variable attribute
- The packed variable attribute
- The init_priority variable attribute (C++ only)
- The mode variable attribute
- The section variable attribute
- The weak variable attribute

Related information

- "Type attributes" on page 74
- "Function attributes" on page 206

The aligned variable attribute

The aligned variable attribute allows you to override the default alignment mode to specify a minimum alignment value, expressed as a number of bytes, for any of the following:

- a non-aggregate variable
• an aggregate variable (such as a structure, class, or union)
• selected member variables

The attribute is typically used to increase the alignment of the given variable.

**aligned variable attribute syntax**

[Code]

The **alignment_factor** is the number of bytes, specified as a constant expression that evaluates to a positive power of 2. You can specify a value up to a maximum of 1048576 bytes. If you omit the alignment factor (and its enclosing parentheses) the compiler automatically uses 16 bytes. If you specify an alignment factor greater than the maximum, the attribute specification is ignored, and the compiler simply uses the default alignment in effect.

When you apply the **aligned** attribute to a bit field structure member variable, the attribute specification is applied to the bit field **container**. If the default alignment of the container is greater than the alignment factor, the default alignment is used.

In the following example, the structures **first_address** and **second_address** are set to an alignment of 16 bytes:

```c
struct address |
  int street_no;
  char *street_name;
  char *city;
  char *prov;
  char *postal_code;
} first_address __attribute__((__aligned__(16))))

struct address second_address __attribute__((__aligned__(16))))
```

In the following example, only the members **first_address.prov** and **first_address.postal_code** are set to an alignment of 16 bytes:

```c
struct address |
  int street_no;
  char *street_name;
  char *city;
  char *prov __attribute__((__aligned__(16))))
  char *postal_code __attribute__((__aligned__(16))))
} first_address ;
```

**Related information**

• “The __align qualifier” on page 71
• “Aligning data” in the *XL C/C++ Programming Guide*
• “The __alignof__ operator” on page 136
• “The aligned type attribute” on page 74

**The packed variable attribute**

The variable attribute **packed** allows you to override the default alignment mode, to reduce the alignment for all members of an aggregate, or selected members of an aggregate to the smallest possible alignment: one byte for a member and one bit for a bit field member.
The init_priority variable attribute (C++ only)

The variable attribute init_priority is an extension to C++ that allows you to control the initialization order of static objects defined in namespace scope across multiple compilation units.

init_priority variable attribute syntax

```
__attribute__((init_priority(relative_priority)))
```

The relative_priority is a constant integral expression between 101 and 65535, inclusive. A lower number indicates a higher priority.

Related information

- “Initializing static objects in libraries” in the XL C/C++ Programming Guide

The mode variable attribute

The variable attribute mode allows you to override the type specifier in a variable declaration, to specify the size of a particular integral type.

mode variable attribute syntax

```
__attribute__((mode(byte)))
```

The valid argument for the mode is any of the of the following type specifiers that indicates a specific width:

- byte means a 1-byte integer type
- word means a 4-byte integer type
- pointer means 4-byte integer type in 32-bit mode and an 8-byte integer type in 64-bit mode

The section variable attribute

The section variable attribute specifies the section in the object file in which the compiler should place its generated code. The language feature provides the ability to control the section in which a variable should appear.
The section name specifies a named section as a string literal, maximum length of 16 characters, not counting spaces. Spaces in the string are ignored.

The section variable attribute can be applied to a declaration or definition of the following types of variables:
- initialized or static global or namespace variables
- static local variables
- uninitialized global or namespace variables
- static structure or class member variables

A section attribute applied to a local variable with automatic storage duration is ignored with a warning because such variables are stored on the stack.

A section attribute applied to a structure member is ignored with a warning. A section attribute applied to an uninitialized global variable is ignored without a warning; the symbols for uninitialized global variables are always placed in the common section.

When multiple section attributes are applied to a variable declaration, the last specification prevails. The section indicated in the prevailing variable declaration should match that of the variable definition because a variable definition cannot be overwritten. Each defined variable can reside in only one section.

A named section can be used for multiple variables, but not for both variables and functions in the same compilation unit.

Related information
- “The section function attribute” on page 211

The weak variable attribute

The weak variable attribute causes the symbol resulting from the variable declaration to appear in the object file as a weak symbol, rather than a global one. The language feature provides the programmer writing library functions with a way to allow variable definitions in user code to override the library declaration without causing duplicate name errors.

weak variable attribute syntax

Related information
- #pragma weak in the XL C/C++ Compiler Reference
- “The weak function attribute” on page 212
End of IBM extension
Chapter 5. Type conversions

An expression of a given type is *implicitly converted* in the following situations:

- The expression is used as an operand of an arithmetic or logical operation.
- The expression is used as a condition in an if statement or an iteration statement (such as a for loop). The expression will be converted to a Boolean (or an integer in C89).
- The expression is used in a switch statement. The expression will be converted to an integral type.
- The expression is used as an initialization. This includes the following:
  - An assignment is made to an lvalue that has a different type than the assigned value.
  - A function is provided an argument value that has a different type than the parameter.
  - The value specified in the return statement of a function has a different type from the defined return type for the function.

You can perform *explicit type* conversions using a cast expression, as described in “Cast expressions” on page 144. The following sections discuss the conversions that are allowed by either implicit or explicit conversion, and the rules governing type promotions:

- “Arithmetic conversions and promotions”
- “Integral-to-rvalue conversions” on page 110
- “Pointer conversions” on page 110
- “Reference conversions (C++ only)” on page 112
- “Qualification conversions (C++ only)” on page 112
- “Function argument conversions” on page 112

Related information

- “User-defined conversions” on page 320
- “Conversion by constructor” on page 322
- “Conversion functions” on page 324
- “The switch statement” on page 173
- “The if statement” on page 172
- “The return statement” on page 183

Arithmetic conversions and promotions

The following sections discuss the rules for the standard conversions for arithmetic types:

- “Integral conversions” on page 108
- “Floating-point conversions” on page 108
- “Boolean conversions” on page 108

If two different types are operands in an expression, they are subject to the rules of the *usual arithmetic conversions*, as described in “Integral and floating-point promotions” on page 109.
Integral conversions

**Unsigned integer to unsigned integer or signed integer to signed integer**
If the types are identical, there is no change. If the types are of a different size, and the value can be represented by the new type, the value is not changed; if the value cannot be represented by the new type, truncation or sign shifting will occur.

**Signed integer to unsigned integer**
The resulting value is the smallest unsigned integer type congruent to the source integer. If the value cannot be represented by the new type, truncation or sign shifting will occur.

**Unsigned integer to signed integer**
If the signed type is large enough to hold the original value, there is no change. If the value can be represented by the new type, the value is not changed; if the value cannot be represented by the new type, truncation or sign shifting will occur.

**Signed and unsigned character types to integer**
If the original value can be represented by int, it is represented as int. If the value cannot be represented by int, it is promoted to unsigned int.

**Wide character type wchar_t to integer**
If the original value can be represented by int, it is represented as int. If the value cannot be represented by int, it is promoted to the smallest type that can hold: unsigned int, long, or unsigned long.

**Signed and unsigned integer bit field to integer**
If the original value can be represented by int, it is represented as int. If the value cannot be represented by int, it is promoted to unsigned int.

**Enumeration type to integer**
If the original value can be represented by int, it is represented as int. If the value cannot be represented by int, it is promoted to the smallest type that can hold: unsigned int, long, or unsigned long. Note that an enumerated type can be converted to an integral type, but an integral type cannot be converted to an enumeration.

Boolean conversions

**Boolean to integer**

- C If the Boolean value is 0, the result is an int with a value of 0. If the Boolean value is 1, the result is an int with a value of 1.
- C++ If the Boolean value is false, the result is an int with a value of 0. If the Boolean value is true, the result is an int with a value of 1.

**Scalar to Boolean**

- C If the scalar value is equal to 0, the Boolean value is 0; otherwise the Boolean value is 1.
- C++ A zero, null pointer, or null member pointer value is converted to false. All other values are converted to true.

Floating-point conversions

**Real to real or complex to complex**
If the types are identical, there is no change. If the types are of a different size, and the value can be represented by the new type, the value is not changed; if the value cannot be represented by the new type, rounding and loss of precision will occur.
Complex to real
The imaginary part of the complex value is discarded. The value of the real part is converted according to the "real to real" rule given above.

Real to complex
The value of the real part is converted according to the "real to real" rule given above. The value of the imaginary part is zero.

Integral and floating-point promotions
When different arithmetic types are used as operands in certain types of expressions, standard conversions known as usual arithmetic conversions are applied. These conversions are applied according to the rank of the arithmetic type: the operand with a type of lower rank is converted to the type of the operand with a higher rank. This is known as integral or floating point promotion.

For example, when the values of two different integral types are added together, both values are first converted to the same type: when a short int value and an int value are added together, the short int value is converted to the int type. Chapter 6, “Expressions and operators,” on page 115 provides a list of the operators and expressions that participate in the usual arithmetic conversions.

The ranking of arithmetic types, listed from highest to lowest, is as follows:

Table 16. Conversion rankings for floating-point types

<table>
<thead>
<tr>
<th>Operand type</th>
</tr>
</thead>
<tbody>
<tr>
<td>long double or long double _Complex</td>
</tr>
<tr>
<td>double or double _Complex</td>
</tr>
<tr>
<td>float or float _Complex</td>
</tr>
</tbody>
</table>

Table 17. Conversion rankings for integer types

<table>
<thead>
<tr>
<th>Operand type</th>
</tr>
</thead>
<tbody>
<tr>
<td>unsigned long long or unsigned long int</td>
</tr>
<tr>
<td>long long or long int</td>
</tr>
<tr>
<td>unsigned long int</td>
</tr>
<tr>
<td>long int¹</td>
</tr>
<tr>
<td>unsigned int¹</td>
</tr>
<tr>
<td>int and enumerated types</td>
</tr>
<tr>
<td>short int</td>
</tr>
<tr>
<td>char, signed char and unsigned char</td>
</tr>
<tr>
<td>Boolean</td>
</tr>
</tbody>
</table>

Notes:
1. If one operand has unsigned int type and the other operand has long int type but the value of the unsigned int cannot be represented in a long int, both operands are converted to unsigned long int.

Related information
- "Integral types" on page 50
- "Boolean types" on page 50
- "Floating-point types" on page 51
Lvalue-to-rvalue conversions

If an lvalue appears in a situation in which the compiler expects an rvalue, the compiler converts the lvalue to an rvalue. The following table lists exceptions to this:

<table>
<thead>
<tr>
<th>Situation before conversion</th>
<th>Resulting behavior</th>
</tr>
</thead>
<tbody>
<tr>
<td>The lvalue is a function type.</td>
<td></td>
</tr>
<tr>
<td>The lvalue is an array.</td>
<td></td>
</tr>
<tr>
<td>The type of the lvalue is an incomplete type.</td>
<td>compile-time error</td>
</tr>
<tr>
<td>The lvalue refers to an uninitialized object.</td>
<td>undefined behavior</td>
</tr>
<tr>
<td>The lvalue refers to an object not of the type of the rvalue, nor of a type derived from the type of the rvalue.</td>
<td>undefined behavior</td>
</tr>
</tbody>
</table>

The type after conversion is not qualified by either const or volatile.

Related information

- “Lvalues and rvalues” on page 115

Pointer conversions

Pointer conversions are performed when pointers are used, including pointer assignment, initialization, and comparison.

C only

Conversions that involve pointers must use an explicit type cast. The exceptions to this rule are the allowable assignment conversions for C pointers. In the following table, a const-qualified lvalue cannot be used as a left operand of the assignment.

<table>
<thead>
<tr>
<th>Left operand type</th>
<th>Permitted right operand types</th>
</tr>
</thead>
<tbody>
<tr>
<td>pointer to (object) $T$</td>
<td>• the constant 0&lt;br&gt;• a pointer to a type compatible with $T$&lt;br&gt;• a pointer to void ($void*$)</td>
</tr>
<tr>
<td>pointer to (function) $F$</td>
<td>• the constant 0&lt;br&gt;• a pointer to a function compatible with $F$</td>
</tr>
</tbody>
</table>

The referenced type of the left operand must have the same qualifiers as the right operand. An object pointer may be an incomplete type if the other pointer has type $void*$.

End of C only

Zero constant to null pointer

A constant expression that evaluates to zero is a null pointer constant. This expression can be converted to a pointer. This pointer will be a null pointer.
(pointer with a zero value), and is guaranteed not to point to any object. 

A constant expression that evaluates to zero can also be converted to the null pointer to a member.

**Array to pointer**

An lvalue or rvalue with type "array of N," where N is the type of a single element of the array, to N*. The result is a pointer to the initial element of the array. A conversion cannot be performed if the expression is used as the operand of the & (address) operator or the sizeof operator.

**Function to pointer**

An lvalue that is a function can be converted to an rvalue that is a pointer to a function of the same type, except when the expression is used as the operand of the & (address) operator, the () (function call) operator, or the sizeof operator.

**Related information**

- "Pointers" on page 82
- "Integer constant expressions" on page 118
- "Arrays" on page 85
- "Pointers to functions" on page 219
- "Pointers to members" on page 260
- "Pointer conversions" on page 294

**Conversion to void***

C pointers are not necessarily the same size as type int. Pointer arguments given to functions should be explicitly cast to ensure that the correct type expected by the function is being passed. The generic object pointer in C is void*, but there is no generic function pointer.

Any pointer to an object, optionally type-qualified, can be converted to void*, keeping the same const or volatile qualifications.

<table>
<thead>
<tr>
<th>C only</th>
</tr>
</thead>
<tbody>
<tr>
<td>The allowable assignment conversions involving void* as the left operand are shown in the following table.</td>
</tr>
</tbody>
</table>

**Table 19. Legal assignment conversions in C for void***

<table>
<thead>
<tr>
<th>Left operand type</th>
<th>Permitted right operand types</th>
</tr>
</thead>
<tbody>
<tr>
<td>(void*)</td>
<td>• The constant 0.</td>
</tr>
<tr>
<td></td>
<td>• A pointer to an object. The object may be of incomplete type.</td>
</tr>
<tr>
<td></td>
<td>• (void*)</td>
</tr>
</tbody>
</table>

| End of C only |

<table>
<thead>
<tr>
<th>C++ only</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pointers to functions cannot be converted to the type void* with a standard conversion: this can be accomplished explicitly, provided that a void* has sufficient bits to hold it.</td>
</tr>
</tbody>
</table>

| End of C++ only |
Reference conversions (C++ only)

A reference conversion can be performed wherever a reference initialization occurs, including reference initialization done in argument passing and function return values. A reference to a class can be converted to a reference to an accessible base class of that class as long as the conversion is not ambiguous. The result of the conversion is a reference to the base class subobject of the derived class object.

Reference conversion is allowed if the corresponding pointer conversion is allowed.

Related information
• “References (C++ only)” on page 88
• “Initialization of references (C++ only)” on page 100
• “Function calls” on page 213
• “Function return values” on page 202

Qualification conversions (C++ only)

An type-qualified rvalue of any type, containing zero or more const or volatile qualifications, can be converted to an rvalue of type-qualified type where the second rvalue contains more const or volatile qualifications than the first rvalue.

An rvalue of type pointer to member of a class can be converted to an rvalue of type pointer to member of a class if the second rvalue contains more const or volatile qualifications than the first rvalue.

Related information
• “Type qualifiers” on page 68

Function argument conversions

When a function is called, if a function declaration is present and includes declared argument types, the compiler performs type checking. The compiler compares the data types provided by the calling function with the data types that the called function expects and performs necessary type conversions. For example, when function funct is called, argument f is converted to a double, and argument c is converted to an int:

```c
char * funct (double d, int i);
/* ... */
int main(void)
{
  float f;
  char c;
  funct(f, c) /* f is converted to a double, c is converted to an int */
  return 0;
}
```

If no function declaration is visible when a function is called, or when an expression appears as an argument in the variable part of a prototype argument list, the compiler performs default argument promotions or converts the value of the expression before passing any arguments to the function. The automatic conversions consist of the following:
• Integral and floating-point values are promoted.
• Arrays or functions are converted to pointers.
• C++ Non-static class member functions are converted to pointers to members.

Related information
• “Integral and floating-point promotions” on page 109
• “The transparent_union type attribute (C only)” on page 76
• “Function call operator ()” on page 121
• “Function calls” on page 213
Expressions are sequences of operators, operands, and punctuators that specify a computation. The evaluation of expressions is based on the operators that the expressions contain and the context in which they are used. An expression can result in a value and can produce side effects. A side effect is a change in the state of the execution environment.

The following sections describe these types of expressions:

- “Lvalues and rvalues”
- “Primary expressions” on page 117
- “Postfix expressions” on page 121
- “Unary expressions” on page 131
- “Cast expressions” on page 144
- “Binary expressions” on page 147
- “Conditional expressions” on page 156
- “Assignment expressions” on page 158
- “Comma expressions” on page 160
- “throw expressions (C++ only)” on page 162

“Operator precedence and associativity” on page 162 provides tables listing the precedence of all the operators described in the sections on postfix, unary, and binary expressions.

C++ operators can be defined to behave differently when applied to operands of class type. This is called operator overloading, and is described in “Overloading operators” on page 231.

**Lvalues and rvalues**

An object is a region of storage that can be examined and stored into. An lvalue is an expression that refers to such an object. An lvalue does not necessarily permit modification of the object it designates. For example, a const object is an lvalue that cannot be modified. The term modifiable lvalue is used to emphasize that the lvalue allows the designated object to be changed as well as examined. The following object types are lvalues, but not modifiable lvalues:

- An array type
- An incomplete type
- A const-qualified type
- A structure or union type with one of its members qualified as a const type

Because these lvalues are not modifiable, they cannot appear on the left side of an assignment statement.

The term rvalue refers to a data value that is stored at some address in memory. An rvalue is an expression that cannot have a value assigned to it. Both a literal constant and a variable can serve as an rvalue. When an lvalue appears in a context that requires an rvalue, the lvalue is implicitly converted to an rvalue. The
reverse, however, is not true: an rvalue cannot be converted to an lvalue. Rvalues always have complete types or the void type.

C defines a function designator as an expression that has function type. A function designator is distinct from an object type or an lvalue. It can be the name of a function or the result of dereferencing a function pointer. The C language also differentiates between its treatment of a function pointer and an object pointer.

In both C and C++, certain operators require lvalues for some of their operands. The table below lists these operators and additional constraints on their usage.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Requirement</th>
</tr>
</thead>
<tbody>
<tr>
<td>&amp; (unary)</td>
<td>Operand must be an lvalue.</td>
</tr>
<tr>
<td>++ --</td>
<td>Operand must be an lvalue. This applies to both prefix and postfix forms.</td>
</tr>
<tr>
<td>*= /= %= &lt;&lt;= &gt;&gt;= &amp;= ^=</td>
<td>=</td>
</tr>
</tbody>
</table>

For example, all assignment operators evaluate their right operand and assign that value to their left operand. The left operand must be a modifiable lvalue or a reference to a modifiable object.

The address operator (&) requires an lvalue as an operand while the increment (++) and the decrement (--) operators require a modifiable lvalue as an operand. The following example shows expressions and their corresponding lvalues.

<table>
<thead>
<tr>
<th>Expression</th>
<th>Lvalue</th>
</tr>
</thead>
<tbody>
<tr>
<td>x = 42</td>
<td>x</td>
</tr>
<tr>
<td>*ptr = newvalue</td>
<td>*ptr</td>
</tr>
<tr>
<td>a++</td>
<td>a</td>
</tr>
<tr>
<td>int&amp; f()</td>
<td>The function call to f()</td>
</tr>
</tbody>
</table>

IBM extension

When compiled with the GNU C language extensions enabled, compound expressions, conditional expressions, and casts are allowed as lvalues, provided that their operands are lvalues. The use of this language extension is deprecated for C++ code.

A compound expression can be assigned if the last expression in the sequence is an lvalue. The following expressions are equivalent:

(x + 1, y) *= 42;
x + 1, (y *= 42);

The address operator can be applied to a compound expression, provided the last expression in the sequence is an lvalue. The following expressions are equivalent:

&(x + 1, y);
x + 1, &y;
A conditional expression can be a valid lvalue if its type is not void and both of its branches for true and false are valid lvalues. Casts are valid lvalues if the operand is an lvalue. The primary restriction is that you cannot take the address of an lvalue cast.

Related information
- “Arrays” on page 85
- “Lvalue-to-rvalue conversions” on page 110

Primary expressions

Primary expressions fall into the following general categories:

- Names (identifiers)
- Literals (constants)
- Integer constant expressions
- Identifier expressions (C++ only)
- Parenthesized expressions ( )
- > C++ The this pointer (described in “The this pointer” on page 261)
- > C++ Names qualified by the scope resolution operator (::)

Names

The value of a name depends on its type, which is determined by how that name is declared. The following table shows whether a name is an lvalue expression.

Table 20. Primary expressions: Names

<table>
<thead>
<tr>
<th>Name declared as</th>
<th>Evaluates to</th>
<th>Is an lvalue?</th>
</tr>
</thead>
<tbody>
<tr>
<td>Variable of arithmetic, pointer, enumeration, structure, or union type</td>
<td>An object of that type</td>
<td>yes</td>
</tr>
<tr>
<td>Enumeration constant</td>
<td>The associated integer value</td>
<td>no</td>
</tr>
<tr>
<td>Array</td>
<td>That array. In contexts subject to conversions, a pointer to the first object in the array, except where the name is used as the argument to the sizeof operator.</td>
<td>C no</td>
</tr>
<tr>
<td>Function</td>
<td>That function. In contexts subject to conversions, a pointer to that function, except where the name is used as the argument to the sizeof operator, or as the function in a function call expression.</td>
<td>C no</td>
</tr>
</tbody>
</table>

As an expression, a name may not refer to a label, typedef name, structure member, union member, structure tag, union tag, or enumeration tag. Names used for these purposes reside in a namespace that is separate from that of names used in expressions. However, some of these names may be referred to within expressions by means of special constructs: for example, the dot or arrow operators.
may be used to refer to structure and union members; typedef names may be used in casts or as an argument to the sizeof operator.

**Literals**
A literal is a numeric constant or string literal. When a literal is evaluated as an expression, its value is a constant. A lexical constant is never an lvalue. However, a string literal is an lvalue.

**Related information**
- “Literals” on page 19
- “The this pointer” on page 261

**Integer constant expressions**
An integer compile-time constant is a value that is determined during compilation and cannot be changed at run time. An integer compile-time constant expression is an expression that is composed of constants and evaluated to a constant.

An integer constant expression is an expression that is composed of only the following:
- literals
- enumerators
- const variables
- static data members of integral or enumeration types
- casts to integral types
- sizeof expressions, where the operand is not a variable length array

The sizeof operator applied to a variable length array type is evaluated at run time, and therefore is not a constant expression.

You must use an integer constant expression in the following situations:
- In the subscript declarator as the description of an array bound.
- After the keyword case in a switch statement.
- In an enumerator, as the numeric value of an enumeration constant.
- In a bit-field width specifier.
- In the preprocessor #if statement. (Enumeration constants, address constants, and sizeof cannot be specified in a preprocessor #if statement.)

**Related information**
- “The sizeof operator” on page 137

**Identifier expressions (C++ only)**
An identifier expression, or id-expression, is a restricted form of primary expression. Syntactically, an id-expression requires a higher level of complexity than a simple identifier to provide a name for all of the language elements of C++.

An id-expression can be either a qualified or unqualified identifier. It can also appear after the dot and arrow operators.
Identifier expression syntax

unqualified_id:

- identifier
- operator_function_id
- conversion_function_id
- class_name
- template_id

qualified_id:

- identifier
- operator_function_id
- template_id
- class_or_namespace::
- template
- class_or_namespace::
- template
- unqualified_id

Related information
- “Identifiers” on page 15
- Chapter 4, “Declarators,” on page 79

Parenthesized expressions ( )

Use parentheses to explicitly force the order of expression evaluation. The following expression does not use parentheses to group operands and operators. The parentheses surrounding weight, zipcode are used to form a function call. Note how the compiler groups the operands and operators in the expression according to the rules for operator precedence and associativity:

expression
+ function call
  parameters
    expression
    expression
    expression
unary minus
  * discount
  * item
  + handling
    ( weight , zipcode )

The following expression is similar to the previous expression, but it contains parentheses that change how the operands and operators are grouped:
In an expression that contains both associative and commutative operators, you can use parentheses to specify the grouping of operands with operators. The parentheses in the following expression guarantee the order of grouping operands with the operators:

\[ x = f + (g + h); \]

**Related information**
- "Operator precedence and associativity” on page 162

**Scope resolution operator  :: (C++ only)**

The :: (scope resolution) operator is used to qualify hidden names so that you can still use them. You can use the unary scope operator if a namespace scope or global scope name is hidden by an explicit declaration of the same name in a block or class. For example:

```c
int count = 0;

int main(void) {
    int count = 0;
    ::count = 1; // set global count to 1
    count = 2;   // set local count to 2
    return 0;
}
```

The declaration of `count` declared in the `main` function hides the integer named `count` declared in global namespace scope. The statement `::count = 1` accesses the variable named `count` declared in global namespace scope.

You can also use the class scope operator to qualify class names or class member names. If a class member name is hidden, you can use it by qualifying it with its class name and the class scope operator.
In the following example, the declaration of the variable X hides the class type X, but you can still use the static class member count by qualifying it with the class type X and the scope resolution operator.

```cpp
#include <iostream>
using namespace std;

class X
{
public:
    static int count;
};
int X::count = 10; // define static data member

int main ()
{
    int X = 0;       // hides class type X
    cout << X::count << endl; // use static member of class X
}
```

Related information
- “Scope of class names” on page 249
- Chapter 9, “Namespaces (C++ only),” on page 221

### Postfix expressions

Postfix operators are operators that appear after their operands. A postfix expression is a primary expression, or a primary expression that contains a postfix operator. The available postfix operators are:

- **Function call operator ( )**
- **Array subscripting operator [ ]**
- **Dot operator .**
- **Arrow operator −>**
- **Compound literals**
  - ```
  > C++ `typeid`
  > C++ `static_cast`
  > C++ `reinterpret_cast`
  > C++ `const_cast`
  > C++ `dynamic_cast`
  ```

The ranking, precedence and associativity rules for the postfix operators are summarized in Table 23 on page 163.

### Function call operator ( )

A function call is an expression containing the function name followed by the function call operator, ( ). If the function has been defined to receive parameters, the values that are to be sent into the function are listed inside the parentheses of the function call operator. The argument list can contain any number of expressions separated by commas. It can also be empty.

The type of a function call expression is the return type of the function. This type can either be a complete type, a reference type, or the type `void`. A function call expression is always an rvalue. A function call is an lvalue if and only if the type of the function is a reference.
Here are some examples of the function call operator:

stub()
overdue(account, date, amount)
notify(name, date + 5)
report(error, time, date, ++num)

The order of evaluation for function call arguments is not specified. In the following example:
method(sample1, batch.process--, batch.process);
the argument batch.process-- might be evaluated last, causing the last two arguments to be passed with the same value.

**Related information**
- “Function argument conversions” on page 112
- “Function calls” on page 213

### Array subscripting operator [ ]

A postfix expression followed by an expression in [ ] (brackets) specifies an element of an array. The expression within the brackets is referred to as a subscript. The first element of an array has the subscript zero.

By definition, the expression a[b] is equivalent to the expression *((a) + (b)), and, because addition is associative, it is also equivalent to b[a]. Between expressions a and b, one must be a pointer to a type T, and the other must have integral or enumeration type. The result of an array subscript is an lvalue. The following example demonstrates this:

```c
#include <stdio.h>

int main(void) {
    int a[3] = { 10, 20, 30 };
    printf("a[0] = %d\n", a[0]);
    printf("a[1] = %d\n", 1[a]);
    printf("a[2] = %d\n", *(2 + a));
    return 0;
}
```

The following is the output of the above example:

a[0] = 10
a[1] = 20
a[2] = 30

The above restrictions on the types of expressions required by the subscript operator, as well as the relationship between the subscript operator and pointer arithmetic, do not apply if you overload operator[] of a class.

The first element of each array has the subscript 0. The expression contract[35] refers to the 36th element in the array contract.

In a multidimensional array, you can reference each element (in the order of increasing storage locations) by incrementing the right-most subscript most frequently.

For example, the following statement gives the value 100 to each element in the array code[4][3][6]:

```c
#include <stdio.h>

int main(void) {
    int a[3] = { 10, 20, 30 };
    printf("a[0] = %d\n", a[0]);
    printf("a[1] = %d\n", 1[a]);
    printf("a[2] = %d\n", *(2 + a));
    return 0;
}
```
for (first = 0; first < 4; ++first)
{
    for (second = 0; second < 3; ++second)
    {
        for (third = 0; third < 6; ++third)
        {
            code[first][second][third] = 100;
        }
    }
}

C only

C99 allows array subscripting on arrays that are not lvalues. However, using the
taddress of a non-lvalue as an array subscript is still not allowed. The following
element is valid in C99:
struct trio{int a[3];};
struct trio f();
foo (int index)
{
    return f().a[index];
}

Related information
• “Pointers” on page 82
• “Integral types” on page 50
• “Lvalues and rvalues” on page 115
• “Arrays” on page 85
• “Overloading subscripting” on page 239
• “Pointer arithmetic” on page 83

Dot operator .
The . (dot) operator is used to access class, structure, or union members. The
member is specified by a postfix expression, followed by a . (dot) operator,
followed by a possibly qualified identifier or a pseudo-destructor name. (A
pseudo-destructor is a destructor of a nonclass type.) The postfix expression must be
an object of type class, struct or union. The name must be a member of that
object.

The value of the expression is the value of the selected member. If the postfix
expression and the name are lvalues, the expression value is also an lvalue. If the
postfix expression is type-qualified, the same type qualifiers will apply to the
designated member in the resulting expression.

Related information
• “Access to structure and union members” on page 61
• “Pseudo-destructors” on page 314

Arrow operator −>
The −> (arrow) operator is used to access class, structure or union members using a
pointer. A postfix expression, followed by an −> (arrow) operator, followed by a
possibly qualified identifier or a pseudo-destructor name, designates a member of
the object to which the pointer points. (A pseudo-destructor is a destructor of a
nonclass type.) The postfix expression must be a pointer to an object of type class,
struct or union. The name must be a member of that object.

The value of the expression is the value of the selected member. If the name is an
lvalue, the expression value is also an lvalue. If the expression is a pointer to a
qualified type, the same type-qualifiers will apply to the designated member in the
resulting expression.

Related information
• “Pointers” on page 82
• “Access to structure and union members” on page 61
• Chapter 12, “Class members and friends (C++ only),” on page 255
• “Pseudo-destructors” on page 314

Compound literals

A compound literal is a postfix expression that provides an unnamed object whose
value is given by an initializer list. The C99 language feature allows you to pass
parameters to functions without the need for temporary variables. It is useful for
specifying constants of an aggregate type (arrays, structures, and unions) when
only one instance of such types is needed.

The syntax for a compound literal resembles that of a cast expression. However, a
compound literal is an lvalue, while the result of a cast expression is not.
Furthermore, a cast can only convert to scalar types or void, whereas a compound
literal results in an object of the specified type.

Compound literal syntax

```
\texttt{\textbackslash \{<\text{type\_name}\rangle,\ldots\text{\{}\texttt{\textbackslash \{}<\text{initializer\_list}\rangle\texttt{\}\}}}
```

The type_name can be any data type, including vectors, and user-defined types. It
can be an array of unknown size, but not a variable length array. If the type is an
array of unknown size, the size is determined by the initializer list.

The following example passes a constant structure variable of type point
containing two integer members to the function drawline:

```
drawline((\text{struct \ point\})\{6,7\});
```

If the compound literal occurs outside the body of a function, the initializer list
must consist of constant expressions, and the unnamed object has static storage
duration. If the compound literal occurs within the body of a function, the
initializer list need not consist of constant expressions, and the unnamed object has
automatic storage duration.

IBM For compatibility with GNU C, a static variable can be initialized with a
compound literal of the same type, provided that all the initializers in the
initializer list are constant expressions.

Related information
The typeid operator (C++ only)

The typeid operator provides a program with the ability to retrieve the actual derived type of the object referred to by a pointer or a reference. This operator, along with the dynamic_cast operator, are provided for runtime type identification (RTTI) support in C++.

typeid operator syntax

```
    typeid(expr)       type-name
```

The typeid operator requires runtime type information (RTTI) to be generated, which must be explicitly specified at compile time through a compiler option.

The typeid operator returns an lvalue of type `const std::type_info` that represents the type of expression `expr`. You must include the standard template library header `<typeinfo>` to use the typeid operator.

If `expr` is a reference or a dereferenced pointer to a polymorphic class, typeid will return a type_info object that represents the object that the reference or pointer denotes at run time. If it is not a polymorphic class, typeid will return a type_info object that represents the type of the reference or dereferenced pointer. The following example demonstrates this:

```cpp
#include <iostream>
#include <typeinfo>
using namespace std;

struct A { virtual ~A() {} };
struct B : A {};
struct C {};
struct D : C {};

int main() {
    B bobj;
    A* ap = &bobj;
    A& ar = bobj;
    cout << "ap: " << typeid(*ap).name() << endl;
    cout << "ar: " << typeid(ar).name() << endl;

    D dobj;
    C* cp = &dobj;
    C& cr = dobj;
    cout << "cp: " << typeid(*cp).name() << endl;
    cout << "cr: " << typeid(cr).name() << endl;
}
```

The following is the output of the above example:

```
ap: B
ar: B
cp: C
cr: C
```

Classes A and B are polymorphic; classes C and D are not. Although `cp` and `cr` refer to an object of type D, typeid(*cp) and typeid(cr) return objects that represent class C.
Lvalue-to-rvalue, array-to-pointer, and function-to-pointer conversions will not be applied to `expr`. For example, the output of the following example will be `int [10]`, not `int *`:

```cpp
#include <iostream>
#include <typeinfo>
using namespace std;

int main() {
  int myArray[10];
  cout << typeid(myArray).name() << endl;
}
```

If `expr` is a class type, that class must be completely defined.

The `typeid` operator ignores top-level `const` or `volatile` qualifiers.

Related information
- “Type names” on page 81
- “The typeof operator” on page 139

### The `static_cast` operator (C++ only)

The `static_cast` operator converts a given expression to a specified type.

#### `static_cast` operator syntax

```cpp
static_cast<Type>(expression)
```

The following is an example of the `static_cast` operator.

```cpp
#include <iostream>
using namespace std;

int main() {
  int j = 41;
  int v = 4;
  float m = j/v;
  float d = static_cast<float>(j)/v;
  cout << "m = " << m << endl;
  cout << "d = " << d << endl;
}
```

The following is the output of the above example:

```
m = 10.25
```

In this example, `m = j/v;` produces an answer of type `int` because both `j` and `v` are integers. Conversely, `d = static_cast<float>(j)/v;` produces an answer of type `float`. The `static_cast` operator converts variable `j` to a type `float`. This allows the compiler to generate a division with an answer of type `float`. All `static_cast` operators resolve at compile time and do not remove any `const` or `volatile` modifiers.

Applying the `static_cast` operator to a null pointer will convert it to a null pointer value of the target type.

You can explicitly convert a pointer of a type `A` to a pointer of a type `B` if `A` is a base class of `B`. If `A` is not a base class of `B`, a compiler error will result.
You may cast an lvalue of a type A to a type B& if the following are true:
• A is a base class of B
• You are able to convert a pointer of type A to a pointer of type B
• The type B has the same or greater const or volatile qualifiers than type A
• A is not a virtual base class of B

The result is an lvalue of type B.

A pointer to member type can be explicitly converted into a different pointer to member type if both types are pointers to members of the same class. This form of explicit conversion may also take place if the pointer to member types are from separate classes, however one of the class types must be derived from the other.

Related information
• "User-defined conversions“ on page 320

The reinterpret_cast operator (C++ only)

A reinterpret_cast operator handles conversions between unrelated types.

reinterpret_cast operator syntax

\[ \text{reinterpret_cast} \rightarrow \text{Type} \rightarrow \text{expression} \]

The reinterpret_cast operator produces a value of a new type that has the same bit pattern as its argument. You cannot cast away a const or volatile qualification. You can explicitly perform the following conversions:
• A pointer to any integral type large enough to hold it
• A value of integral or enumeration type to a pointer
• A pointer to a function to a pointer to a function of a different type
• A pointer to an object to a pointer to an object of a different type
• A pointer to a member to a pointer to a member of a different class or type, if the types of the members are both function types or object types

A null pointer value is converted to the null pointer value of the destination type.

Given an lvalue expression of type T and an object x, the following two conversions are synonymous:
• reinterpret_cast<T&>(x)
• *reinterpret_cast<T*>(&x)

C++ also supports C-style casts. The two styles of explicit casts have different syntax but the same semantics, and either way of reinterpreting one type of pointer as an incompatible type of pointer is usually invalid. The reinterpret_cast operator, as well as the other named cast operators, is more easily spotted than C-style casts, and highlights the paradox of a strongly typed language that allows explicit casts.

The C++ compiler detects and quietly fixes most but not all violations. It is important to remember that even though a program compiles, its source code may not be completely correct. On some platforms, performance optimizations are
predicated on strict adherence to standard aliasing rules. Although the C++
compiler tries to help with type-based aliasing violations, it cannot detect all
possible cases.

The following example violates the aliasing rule, but will execute as expected when
compiled unoptimized in C++ or in K&R C. It will also successfully compile
optimized in C++, but will not necessarily execute as expected. The offending line
7 causes an old or uninitialized value for \( x \) to be printed.

```cpp
1 extern int y = 7.;
2
3 int main() {
4     float x;
5     int i;
6     x = y;
7     i = *(int *) &x;
8     printf("i=%d. x=%f.\n", i, x);
9 }
```

The next code example contains an incorrect cast that the compiler cannot even
detect because the cast is across two different files.

```cpp
1 /* separately compiled file 1 */
2 extern float f;
3 extern int * int_pointer_to_f = (int *) &f; /* suspicious cast */
4
5 /* separately compiled file 2 */
6 extern float f;
7 extern int * int_pointer_to_f;
8 f = 1.0;
9 int i = *int_pointer_to_f; /* no suspicious cast but wrong */
```

In line 8, there is no way for the compiler to know that \( f = 1.0 \) is storing into the
same object that \( \text{int } i = \ast \text{int_pointer_to_f} \) is loading from.

**Related information**

* "User-defined conversions" on page 320

**The const_cast operator (C++ only)**

A `const_cast` operator is used to add or remove a `const` or `volatile` modifier to or
from a type.

**const_cast operator syntax**

```
const_cast<--Type-->(--expression--)
```

`Type` and the type of `expression` may only differ with respect to their `const` and
`volatile` qualifiers. Their cast is resolved at compile time. A single `const_cast`
expression may add or remove any number of `const` or `volatile` modifiers.

The result of a `const_cast` expression is an rvalue unless `Type` is a reference type.
In this case, the result is an lvalue.

Types can not be defined within `const_cast`.

The following demonstrates the use of the `const_cast` operator:

```cpp
#include <iostream>
using namespace std;
```
void f(int* p) {
    cout << *p << endl;
}

int main(void) {
    const int a = 10;
    const int* b = &a;
    // Function f() expects int*, not const int*
    // f(b);
    int* c = const_cast<int*>(b);
    // Lvalue is const
    // *b = 20;
    // Undefined behavior
    // *c = 30;
    int a1 = 40;
    const int* b1 = &a1;
    int* c1 = const_cast<int*>(b1);
    // Integer a1, the object referred to by c1, has
    // not been declared const
    *c1 = 50;

    return 0;
}

The compiler will not allow the function call f(b). Function f() expects a pointer to an int, not a const int. The statement int* c = const_cast<int*>(b) returns a pointer c that refers to a without the const qualification of a. This process of using const_cast to remove the const qualification of an object is called casting away constness. Consequently the compiler will allow the function call f(c).

The compiler would not allow the assignment *b = 20 because b points to an object of type const int. The compiler will allow the *c = 30, but the behavior of this statement is undefined. If you cast away the constness of an object that has been explicitly declared as const, and attempt to modify it, the results are undefined.

However, if you cast away the constness of an object that has not been explicitly declared as const, you can modify it safely. In the above example, the object referred to by b1 has not been declared const, but you cannot modify this object through b1. You may cast away the constness of b1 and modify the value to which it refers.

Related information
  • “Type qualifiers” on page 68

The dynamic_cast operator (C++ only)

The dynamic_cast operator performs type conversions at run time. The dynamic_cast operator guarantees the conversion of a pointer to a base class to a pointer to a derived class, or the conversion of an lvalue referring to a base class to a reference to a derived class. A program can thereby use a class hierarchy safely. This operator and the typeid operator provide runtime type information (RTTI) support in C++.
The expression dynamic_cast<T>(v) converts the expression v to type T. Type T must be a pointer or reference to a complete class type or a pointer to void. If T is a pointer and the dynamic_cast operator fails, the operator returns a null pointer of type T. If T is a reference and the dynamic_cast operator fails, the operator throws the exception std::bad_cast. You can find this class in the standard library header <typeinfo>.

The dynamic_cast operator requires runtime type information (RTTI) to be generated, which must be explicitly specified at compile time through a compiler option.

If T is a void pointer, then dynamic_cast will return the starting address of the object pointed to by v. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual ~A() {
    }
};

struct B : A {
};

int main() {
    B bobj;
    A* ap = &bobj;
    void* vp = dynamic_cast<void*>(ap);
    cout << "Address of vp : " << vp << endl;
    cout << "Address of bobj : " << &bobj << endl;
}
```

The output of this example will be similar to the following. Both vp and &bobj will refer to the same address:

```
Address of vp : 12FF6C
Address of bobj: 12FF6C
```

The primary purpose for the dynamic_cast operator is to perform type-safe downcasts. A downcast is the conversion of a pointer or reference to a class A to pointer or reference to a class B, where class A is a base class of B. The problem with downcasts is that a pointer of type A* can and must point to any object of a class that has been derived from A. The dynamic_cast operator ensures that if you convert a pointer of class A to a pointer of a class B, the object that A points to belongs to class B or a class derived from B.

The following example demonstrates the use of the dynamic_cast operator:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "Class A" << endl; }
};

struct B : A {
    virtual void f() { cout << "Class B" << endl; }
};

struct C : A {
    virtual void f() { cout << "Class C" << endl; }
};

void f(A* arg) {
    B* bp = dynamic_cast<B*>(arg);
    bp->f();
}
```
C* cp = dynamic_cast<C*>(arg);
if (bp)
  bp->f();
else if (cp)
  cp->f();
else
  arg->f();
}

int main() {
  A aobj;
  C cobj;
  A* ap = &cobj;
  A* ap2 = &aobj;
  f(ap);
  f(ap2);
}

The following is the output of the above example:
Class C
Class A

The function f() determines whether the pointer arg points to an object of type A, B, or C. The function does this by trying to convert arg to a pointer of type B, then to a pointer of type C, with the dynamic_cast operator. If the dynamic_cast operator succeeds, it returns a pointer that points to the object denoted by arg. If dynamic_cast fails, it returns 0.

You may perform downcasts with the dynamic_cast operator only on polymorphic classes. In the above example, all the classes are polymorphic because class A has a virtual function. The dynamic_cast operator uses the runtime type information generated from polymorphic classes.

**Related information**
- “Derivation” on page 279
- “User-defined conversions” on page 320

---

**Unary expressions**

A *unary expression* contains one operand and a unary operator.

The supported standard unary operators are:
- Increment operator `++`
- Decrement operator `--`
- Unary plus operator `+`
- Unary minus operator `-`
- Logical negation operator `!`
- Bitwise negation operator `~`
- Address operator `&`
- Indirection operator `*`
- `alignof`
- `sizeof`
- `typeof`
- `Label value operator `&&`
As indicated in the descriptions of the operators, the usual arithmetic conversions are performed on the operands of most unary expressions.

**Related information**
- "Pointer arithmetic" on page 83
- "Lvalues and rvalues" on page 115
- "Arithmetic conversions and promotions" on page 107

**Increment operator ++**

The ++ (increment) operator adds 1 to the value of a scalar operand, or if the operand is a pointer, increments the operand by the size of the object to which it points. The operand receives the result of the increment operation. The operand must be a modifiable lvalue of arithmetic or pointer type.

You can put the ++ before or after the operand. If it appears before the operand, the operand is incremented. The incremented value is then used in the expression. If you put the ++ after the operand, the value of the operand is used in the expression before the operand is incremented. For example:

```
play = ++play1 + play2++;  
```

is similar to the following expressions; play2 is altered before play:

```
int temp, temp1, temp2;

temp1 = play1 + 1;
temp2 = play2;
play1 = temp1;
temp = temp1 + temp2;
play2 = play2 + 1;
play = temp;
```

The result has the same type as the operand after integral promotion.

The usual arithmetic conversions on the operand are performed.

--- IBM extension ---

The increment operator has been extended to handle complex types. The operator works in the same manner as it does on a real type, except that only the real part of the operand is incremented, and the imaginary part is unchanged.

--- End of IBM extension ---
Decrement operator --

The -- (decrement) operator subtracts 1 from the value of a scalar operand, or if the operand is a pointer, decreases the operand by the size of the object to which it points. The operand receives the result of the decrement operation. The operand must be a modifiable lvalue.

You can put the -- before or after the operand. If it appears before the operand, the operand is decremented, and the decremented value is used in the expression. If the -- appears after the operand, the current value of the operand is used in the expression and the operand is decremented.

For example:

```cpp
play = --play1 + play2--;  
```

is similar to the following expressions; play2 is altered before play:

```cpp
int temp, temp1, temp2;

temp1 = play1 - 1;
temp2 = play2;
play1 = temp1;
temp = temp1 + temp2;
play2 = play2 - 1;
play = temp;
```

The result has the same type as the operand after integral promotion, but is not an lvalue.

The usual arithmetic conversions are performed on the operand.

--- IBM extension ---

The decrement operator has been extended to handle complex types, for compatibility with GNU C. The operator works in the same manner as it does on a real type, except that only the real part of the operand is decremented, and the imaginary part is unchanged.

--- End of IBM extension ---

Unary plus operator +

The + (unary plus) operator maintains the value of the operand. The operand can have any arithmetic type or pointer type. The result is not an lvalue.

The result has the same type as the operand after integral promotion.

Note: Any plus sign in front of a constant is not part of the constant.

Unary minus operator

The - (unary minus) operator negates the value of the operand. The operand can have any arithmetic type. The result is not an lvalue.

For example, if quality has the value 100, -quality has the value -100.

The result has the same type as the operand after integral promotion.

Note: Any minus sign in front of a constant is not part of the constant.
Logical negation operator !
The ! (logical negation) operator determines whether the operand evaluates to 0 (false) or nonzero (true).

- **C** The expression yields the value 1 (true) if the operand evaluates to 0, and yields the value 0 (false) if the operand evaluates to a nonzero value.

- **C++** The expression yields the value true if the operand evaluates to false (0), and yields the value false if the operand evaluates to true (nonzero). The operand is implicitly converted to bool, and the type of the result is bool.

The following two expressions are equivalent:
```c
!right;
right == 0;
```

Related information
- "Boolean types" on page 50

Bitwise negation operator ~
The ~ (bitwise negation) operator yields the bitwise complement of the operand. In the binary representation of the result, every bit has the opposite value of the same bit in the binary representation of the operand. The operand must have an integral type. The result has the same type as the operand but is not an lvalue.

Suppose x represents the decimal value 5. The 16-bit binary representation of x is:
```
0000000000000101
```

The expression ~x yields the following result (represented here as a 16-bit binary number):
```
1111111111111010
```

Note that the ~ character can be represented by the trigraph ??-.

The 16-bit binary representation of ~0 is:
```
1111111111111111
```

---

IBM extension

The bitwise negation operator has been extended to handle complex types. With a complex type, the operator computes the complex conjugate of the operand by reversing the sign of the imaginary part.

---

Related information
- "Trigraph sequences" on page 34

Address operator &
The & (address) operator yields a pointer to its operand. The operand must be an lvalue, a function designator, or a qualified name. It cannot be a bit field, nor can it have the storage class register.

---

134 XL C/C++ Language Reference
If the operand is an lvalue or function, the resulting type is a pointer to the expression type. For example, if the expression has type `int`, the result is a pointer to an object having type `int`.

If the operand is a qualified name and the member is not static, the result is a pointer to a member of class and has the same type as the member. The result is not an lvalue.

If `p_to_y` is defined as a pointer to an `int` and `y` as an `int`, the following expression assigns the address of the variable `y` to the pointer `p_to_y`:

```
p_to_y = &y;
```

**IBM extension**

The address operator has been extended to handle vector types, provided that VMX support is enabled. The result of the address operator applied to a vector type can be stored in a pointer to a compatible vector type. The address of a vector type can be used to initialize a pointer to vector type if both sides of the initialization have compatible types. A pointer to `void` can also be initialized with the address of a vector type.

End of IBM extension

**C++ only**

The ampersand symbol `&` is used in C++ as a reference declarator in addition to being the address operator. The meanings are related but not identical.

```
int target;
int &rTarg = target; // rTarg is a reference to an integer.
void f(int*& p);    // p is a reference to a pointer
```

If you take the address of a reference, it returns the address of its target. Using the previous declarations, `&rTarg` is the same memory address as `&target`.

You may take the address of a register variable.

You can use the `&` operator with overloaded functions only in an initialization or assignment where the left side uniquely determines which version of the overloaded function is used.

End of C++ only

**IBM** The address of a label can be taken using the GNU C address operator `&&`. The label can thus be used as a value.

Related information

- "Indirection operator "+" on page 136
- "Pointers" on page 82
- "References (C++ only)" on page 88
- "Labels as values" on page 168
Indirection operator *

The * (indirection) operator determines the value referred to by the pointer-type operand. The operand cannot be a pointer to an incomplete type. If the operand points to an object, the operation yields an lvalue referring to that object. If the operand points to a function, the result is a function designator in C or, in C++, an lvalue referring to the object to which the operand points. Arrays and functions are converted to pointers.

The type of the operand determines the type of the result. For example, if the operand is a pointer to an int, the result has type int.

Do not apply the indirection operator to any pointer that contains an address that is not valid, such as NULL. The result is not defined.

If p_to_y is defined as a pointer to an int and y as an int, the expressions:

\[
p_{-to-y} = &y; \\
*p_{-to-y} = 3;
\]

cause the variable y to receive the value 3.

IBM extension

The indirection operator has been extended to handle vector types, provided that VMX support is enabled.

Related information

- "Arrays" on page 85
- "Pointers" on page 82

The __alignof__ operator

IBM extension

The __alignof__ operator is a language extension to C99 and Standard C++ that returns the number of bytes used in the alignment of its operand. The operand can be an expression or a parenthesized type identifier. If the operand is an expression representing an lvalue, the number returned by __alignof__ represents the alignment that the lvalue is known to have. The type of the expression is determined at compile time, but the expression itself is not evaluated. If the operand is a type, the number represents the alignment usually required for the type on the target platform.

The __alignof__ operator may not be applied to the following:

- An lvalue representing a bit field
- A function type
- An undefined structure or class
- An incomplete type (such as void)

__alignof__ operator syntax

\[
\text{__alignof__ (type-id)}
\]
If `type-id` is a reference or a referenced type, the result is the alignment of the referenced type. If `type-id` is an array, the result is the alignment of the array element type. If `type-id` is a fundamental type, the result is implementation-defined.

For example, on Linux, `__alignof__(long)` returns 4 in 32-bit mode, and 8 in 64-bit mode.

The operand of `__alignof__` can be a vector type, provided that VMX support is enabled. For example,

```c
vector unsigned int v1 = (vector unsigned int)(10);
vector unsigned int *pv1 = &v1;
__alignof__(v1); // vector type alignment: 16.
__alignof__(&v1); // address of vector alignment: 4.
__alignof__(*pv1); // dereferenced pointer to vector alignment: 16.
__alignof__(pv1); // pointer to vector alignment: 4.
__alignof__(vector signed char); // vector type alignment: 16.
```

When `__attribute__((aligned))` is used to increase the alignment of a variable of vector type, the value returned by the `__alignof__` operator is the alignment factor specified by `__attribute__((aligned))`.

Related information
- "The aligned variable attribute" on page 102

---

**The sizeof operator**

The `sizeof` operator yields the size in bytes of the operand, which can be an expression or the parenthesized name of a type.

**sizeof operator syntax**

```c
sizeof(expr) (type-name)
```

The result for either kind of operand is not an lvalue, but a constant integer value. The type of the result is the unsigned integral type `size_t` defined in the header file `stddef.h`.

Except in preprocessor directives, you can use a `sizeof` expression wherever an integral constant is required. One of the most common uses for the `sizeof` operator is to determine the size of objects that are referred to during storage allocation, input, and output functions.

Another use of `sizeof` is in porting code across platforms. You can use the `sizeof` operator to determine the size that a data type represents. For example:

```c
sizeof(int);
```

The `sizeof` operator applied to a type name yields the amount of memory that would be used by an object of that type, including any internal or trailing padding.

---

**IBM extension**

The operand of the `sizeof` operator can be a vector type or the result of dereferencing a pointer to vector type, provided that VMX support is enabled. In these cases, the return value of `sizeof` is always 16.
vector bool int v1;
vector bool int *pv1 = &v1;
sizeof(v1); // vector type: 16.
sizeof(&v1); // address of vector: 4.
sizeof(*pv1); // dereferenced pointer to vector: 16.
sizeof(pv1); // pointer to vector: 4.
sizeof(vector bool int); // vector type: 16.

\[ \text{End of IBM extension} \]

For compound types, results are as follows:

<table>
<thead>
<tr>
<th>Operand</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>An array</td>
<td>The result is the total number of bytes in the array. For example, in an array with 10 elements, the size is equal to 10 times the size of a single element. The compiler does not convert the array to a pointer before evaluating the expression.</td>
</tr>
<tr>
<td>A class</td>
<td>The result is always nonzero, and is equal to the number of bytes in an object of that class including any padding required for placing class objects in an array.</td>
</tr>
<tr>
<td>A reference</td>
<td>The result is the size of the referenced object.</td>
</tr>
</tbody>
</table>

The \texttt{sizeof} operator may not be applied to:
- A bit field
- A function type
- An undefined structure or class
- An incomplete type (such as void)

The \texttt{sizeof} operator applied to an expression yields the same result as if it had been applied to only the name of the type of the expression. At compile time, the compiler analyzes the expression to determine its type. None of the usual type conversions that occur in the type analysis of the expression are directly attributable to the \texttt{sizeof} operator. However, if the operand contains operators that perform conversions, the compiler does take these conversions into consideration in determining the type. For example, the second line of the following sample causes the usual arithmetic conversions to be performed. Assuming that a short uses 2 bytes of storage and an int uses 4 bytes,

```c
short x; ... sizeof (x) /* the value of sizeof operator is 2 */
short x; ... sizeof (x + 1) /* value is 4, result of addition is type int */
```

The result of the expression \(x + 1\) has type \texttt{int} and is equivalent to \texttt{sizeof(int)}. The value is also 4 if \(x\) has type \texttt{char}, \texttt{short}, or \texttt{int} or any enumeration type.

Related information
- “Type names” on page 81
- “Integer constant expressions” on page 118
- “Arrays” on page 85
- “References (C++ only)” on page 88
The **typeof** operator

The typeof operator returns the type of its argument, which can be an expression or a type. The language feature provides a way to derive the type from an expression. The alternate spelling of the keyword, __typeof__, is recommended. Given an expression e, __typeof__(e) can be used anywhere a type name is needed, for example in a declaration or in a cast.

The typeof operator is extended to accept a vector type as its operand, when VMX support is enabled.

**typeof operator syntax**

```
__typeof__(<type-name>)
```

A typeof construct itself is not an expression, but the name of a type. A typeof construct behaves like a type name defined using typedef, although the syntax resembles that of sizeof.

The following examples illustrate its basic syntax. For an expression e:

```c
int e;
__typeof__(e + 1) j; /* the same as declaring int j; */
```

Using a typeof construct is equivalent to declaring a typedef name. Given

```c
int T[2];
int i[2];
```

you can write

```c
__typeof__(i) a; /* all three constructs have the same meaning */
__typeof__(int[2]) a;
__typeof__(T) a;
```

The behavior of the code is as if you had declared `int a[2];`.

For a bit field, typeof represents the underlying type of the bit field. For example, int m:2; the typeof(m) is int. Since the bit field property is not reserved, n in typeof(m) n; is the same as int n, but not int n:2.

The typeof operator can be nested inside sizeof and itself. The following declarations of arr as an array of pointers to int are equivalent:

```c
int *arr[10]; /* traditional C declaration */
__typeof__(__typeof__(int *)[10]) a; /* equivalent declaration */
```

The typeof operator can be useful in macro definitions where expression e is a parameter. For example,

```c
#define SWAP(a,b) { __typeof__(a) temp; temp = a; a = b; b = temp;
```

**Related information**

- "Type names" on page 81
- "typedef definitions" on page 66
- "The typeid operator (C++ only)" on page 125
Label value operator &&

The label value operator && returns the address of its operand, which must be a label defined in the current function or a containing function. The value is a constant of type void* and should be used only in a computed goto statement. The language feature is an extension to C and C++, implemented to facilitate porting programs developed with GNU C.

Related information

- “Labels as values” on page 168
- “Computed goto statement” on page 185

The new operator (C++ only)

The new operator provides dynamic storage allocation.

new operator syntax

\[
\text{new} \left( \begin{array}{c}
\text{argument_list} \\
\text{new_type}
\end{array} \right)
\]

\[
\left( \begin{array}{c}
\text{initial_value}
\end{array} \right)
\]

If you prefix new with the scope resolution operator (::), the global operator new() is used. If you specify an argument_list, the overloaded new operator that corresponds to that argument_list is used. The type is an existing built-in or user-defined type. A new_type is a type that has not already been defined and can include type specifiers and declarators.

An allocation expression containing the new operator is used to find storage in free store for the object being created. The new expression returns a pointer to the object created and can be used to initialize the object. If the object is an array, a pointer to the initial element is returned.

You cannot use the new operator to allocate function types, void, or incomplete class types because these are not object types. However, you can allocate pointers to functions with the new operator. You cannot create a reference with the new operator.

When the object being created is an array, only the first dimension can be a general expression. All subsequent dimensions must be constant integral expressions. The first dimension can be a general expression even when an existing type is used. You can create an array with zero bounds with the new operator. For example:

\[
\text{char } * c = \text{new char}[0];
\]

In this case, a pointer to a unique object is returned.
An object created with operator `new()` or operator `new[]()` exists until the operator `delete()` or operator `delete[]()` is called to deallocate the object's memory. A delete operator or a destructor will not be implicitly called for an object created with a `new` that has not been explicitly deallocated before the end of the program.

If parentheses are used within a new type, parentheses should also surround the new type to prevent syntax errors.

In the following example, storage is allocated for an array of pointers to functions:

```c++
void f();
void g();

int main(void)
{
    void (**p)(), (**q)();
    // declare p and q as pointers to pointers to void functions
    p = new (void (*)(*[3]())); // p now points to an array of pointers to functions
    q = new void(*)(*[3]()); // error: error - bound as 'q = (new void) (*[3])();'
    p[0] = f; // p[0] to point to function f
    q[2] = g; // q[2] to point to function g
    p[0](); // call f()
    q[2](); // call g()
    return (0);
}
```

However, the second use of `new` causes an erroneous binding of `q = (new void)(*[3])()`.

The type of the object being created cannot contain class declarations, enumeration declarations, or const or volatile types. It can contain pointers to const or volatile objects.

For example, `const char*` is allowed, but `char* const` is not.

**Placement syntax**

Additional arguments can be supplied to `new` by using the `argument_list`, also called the placement syntax. If placement arguments are used, a declaration of operator `new()` or operator `new[]()` with these arguments must exist. For example:

```c++
#include <new>
using namespace std;

class X
{
    public:
        void* operator new(size_t, int, int){ /* ... */ }
    };
    // ...

    int main ()
    {
        X* ptr = new(1,2) X;
    }
```

The placement syntax is commonly used to invoke the global placement `new` function. The global placement `new` function initializes an object or objects at the location specified by the placement argument in the placement new expression. This location must address storage that has previously been allocated by some
other means, because the global placement new function does not itself allocate memory. In the following example, no new memory is allocated by the calls new(whole) X(8); new(seg2) X(9); or new(seg3) X(10); Instead, the constructors X(8), X(9), and X(10) are called to reinitialize the memory allocated to the buffer whole.

Because placement new does not allocate memory, you should not use delete to deallocate objects created with the placement syntax. You can only delete the entire memory pool (delete whole). In the example, you can keep the memory buffer but destroy the object stored in it by explicitly calling a destructor.

```cpp
#include <new>
class X
{
    public:
        X(int n): id(n) {}
        ~X() {}
    private:
        int id;
        // ...
};

int main()
{
    char* whole = new char[3 * sizeof(X)]; // a 3-part buffer
    X* p1 = new(whole) X(8); // fill the front
    char* seg2 = &whole[sizeof(X)]; // mark second segment
    X* p2 = new(seg2) X(9); // fill second segment
    char* seg3 = &whole[2 * sizeof(X)]; // mark third segment
    X* p3 = new(seg3) X(10); // fill third segment

    p2->~X(); // clear only middle segment, but keep the buffer
    // ...
    return 0;
}
```

The placement new syntax can also be used for passing parameters to an allocation routine rather than to a constructor.

Related information
- "Free store" on page 315
- "The delete operator (C++ only)"
- "Scope resolution operator :: (C++ only)" on page 120
- "Overview of constructors and destructors” on page 303

The delete operator (C++ only)
The delete operator destroys the object created with new by deallocating the memory associated with the object.

The delete operator has a void return type.

delete operator syntax

```
define delete object_pointer
```

The operand of delete must be a pointer returned by new, and cannot be a pointer to constant. Deleting a null pointer has no effect.
The `delete[]` operator frees storage allocated for array objects created with `new[]`. The `delete` operator frees storage allocated for individual objects created with `new`.

**delete[] operator syntax**

```
::delete[ ] array
```

The result of deleting an array object with `delete` is undefined, as is deleting an individual object with `delete[]`. The array dimensions do not need to be specified with `delete[]`.

The result of any attempt to access a deleted object or array is undefined.

If a destructor has been defined for a class, `delete` invokes that destructor. Whether a destructor exists or not, `delete` frees the storage pointed to by calling the function `operator delete()` of the class if one exists.

The global `::operator delete()` is used if:
  * The class has no `operator delete()`.
  * The object is of a nonclass type.
  * The object is deleted with the `::delete` expression.

The global `::operator delete[]()` is used if:
  * The class has no `operator delete[]()`.
  * The object is of a nonclass type.
  * The object is deleted with the `::delete[]` expression.

The default global operator `delete()` only frees storage allocated by the default global operator `new()`. The default global operator `delete[]()` only frees storage allocated for arrays by the default global operator `new[]()`.

**Related information**

- “Free store” on page 315
- “Overview of constructors and destructors” on page 303
- “The void type” on page 53

**The _Pragma preprocessing operator**

The unary operator `_Pragma`, which is a C99 feature, allows a preprocessor macro to be contained in a pragma directive.

**_Pragma operator syntax**

```
Pragma("string_literal")
```

The `string_literal` may be prefixed with `L`, making it a wide-string literal.

The string literal is destringized and tokenized. The resulting sequence of tokens is processed as if it appeared in a pragma directive. For example:

```
_Pragma( "pack(full)" )
```

would be equivalent to
The __real__ and __imag__ operators

IBM extension

XL C/C++ extends the C99 and C++ standards to support the unary operators __real__ and __imag__. These operators provide the ability to extract the real and imaginary parts of a complex type. These extensions have been implemented to ease the porting applications developed with GNU C.

__real__ and __imag__ operator syntax

\[ \text{__real__}(\var_identifier) \]

\[ \text{__imag__}(\var_identifier) \]

The var_identifier is the name of a previously declared complex variable. The __real__ operator returns the real part of the complex variable, while the __imag__ operator returns the imaginary part of the variable. If the operand of these operators is an lvalue, the resulting expression can be used in any context where lvalues are allowed. They are especially useful in initializations of complex variables, and as arguments to calls to library functions such as printf and scanf that have no format specifiers for complex types. For example:

```c
float _Complex myvar;
__imag__(myvar) = 2.0f;
__real__(myvar) = 3.0f;
```

initializes the imaginary part of the complex variable myvar to 2.0i and the real part to 3.0, and

```c
printf("myvar = %f + %f * i\n", __real__(myvar), __imag__(myvar));
```

prints:

myvar = 2.000000 + 3.000000 * i

Related information

- “Complex literals” on page 24
- “Complex floating-point types” on page 52

End of IBM extension

Cast expressions

The cast operator is used for explicit type conversions. It converts the value of an expression to a specified type.

Cast expression syntax

\[ (\text{--type--})\text{expression} \]
The result of this operation is not an lvalue. The result of this operation is an lvalue if type is a reference; in all other cases, the result is an rvalue.

The following demonstrates the use of the cast operator to dynamically create an integer array of size 10:
#include <stdlib.h>

int main(void) {
    int* myArray = (int*) malloc(10 * sizeof(int));
    free(myArray);
    return 0;
}

The malloc library function returns a void pointer that points to memory that will hold an object of the size of its argument. The statement int* myArray = (int*) malloc(10 * sizeof(int)) does the following:
- Creates a void pointer that points to memory that can hold ten integers.
- Converts that void pointer into an integer pointer with the use of the cast operator.
- Assigns that integer pointer to myArray. Because a name of an array is the same as a pointer to the initial element of the array, myArray is an array of ten integers stored in the memory created by the call to malloc().

---

C++ only

In C++ you can also use the following in cast expressions:
- Function-style casts
- C++ conversion operators, such as static_cast.

Function-style notation converts the value of expression to the type type:

expression( type )

The following example shows the same value cast with a C-style cast, the C++ function-style cast, and a C++ cast operator:
#include <iostream>
using namespace std;

int main() {
    float num = 98.76;
    int x1 = (int) num;
    int x2 = int(num);
    int x3 = static_cast<int>(num);

    cout << "x1 = " << x1 << endl;
    cout << "x2 = " << x2 << endl;
    cout << "x3 = " << x3 << endl;
}

The following is the output of the above example:

x1 = 98
x2 = 98
x3 = 98

The integer x1 is assigned a value in which num has been explicitly converted to an int with the C-style cast. The integer x2 is assigned a value that has been
converted with the function-style cast. The integer x3 is assigned a value that has been converted with the static_cast operator.

A cast is a valid lvalue if its operand is an lvalue. In the following simple assignment expression, the right-hand side is first converted to the specified type, then to the type of the inner left-hand side expression, and the result is stored. The value is converted back to the specified type, and becomes the value of the assignment. In the following example, i is of type char *.

```c
(int)i = 8    // This is equivalent to the following expression
(int)(i = (char*) (int)(8))
```

For compound assignment operation applied to a cast, the arithmetic operator of the compound assignment is performed using the type resulting from the cast, and then proceeds as in the case of simple assignment. The following expressions are equivalent. Again, i is of type char *.

```c
(int)i += 8    // This is equivalent to the following expression
(int)(i = (char*) (int)((int)i = 8))
```

For C++, the operand of a cast expression can have class type. If the operand has class type, it can be cast to any type for which the class has a user-defined conversion function. Casts can invoke a constructor, if the target type is a class, or they can invoke a conversion function, if the source type is a class. They can be ambiguous if both conditions hold.

---

End of C++ only

---

Related information

- “Type names” on page 81
- “Conversion functions” on page 324
- “Conversion by constructor” on page 322
- “Lvalues and rvalues” on page 115
- “Temporary objects” on page 320

Cast to union type (C only)

---

IBM extension

Casting to a union type is the ability to cast a union member to the same type as the union to which it belongs. Such a cast does not produce an lvalue, unlike other casts. The feature is supported as an extension to C99, implemented to facilitate porting programs developed with GNU C.

Only a type that explicitly exists as a member of a union type can be cast to that union type. The cast can use either the tag of the union type or a union type name declared in a typedef expression. The type specified must be a complete union type. An anonymous union type can be used in a cast to a union type, provided that it has a tag or type name. A bit field can be cast to a union type, provided that the union contains a bit field member of the same type, but not necessarily of the same length.

Casting to a nested union is also allowed. In the following example, the double type dd can be cast to the nested union u2_t.

```c
int main() {
    union u2_t {
        char a;
        double dd;
    }
    u2_t uu;
    uu.a = 'c';
    uu.dd = 1.2;
    printf("%c
", uu.a);
    printf("%e
", uu.dd);
}
```
short b;
int c;
union u2_t {
  double d;
} u2;

union u_t U;
double dd = 1.234;
U.u2 = (union u2_t) dd;   // Valid.
printf("U.u2 is \%f\n", U.u2);
}

The output of this example is:
U.u2 is 1.234

A union cast is also valid as a function argument, part of a constant expression for initialization, and in a compound literal statement.

Related information
- “Structures and unions” on page 55
- “The transparent_union type attribute (C only)” on page 76

End of IBM extension

---

Binary expressions

A binary expression contains two operands separated by one operator. The supported binary operators are:
- Multiplication operator *
- Division operator /
- Remainder operator %
- Addition operator +
- Subtraction operator −
- Bitwise left and right shift operators << >>
- Relational operators < <= >=
- Equality and inequality operators == !=
- Bitwise AND operator &
- Bitwise exclusive OR operator ^
- Bitwise inclusive OR operator |
- Logical AND operator &&
- Logical OR operator ||
- Pointer to member operators .* ->* (C++ only)

All binary operators have left-to-right associativity, but not all binary operators have the same precedence. The ranking and precedence rules for binary operators is summarized in Table 25 on page 164.

The order in which the operands of most binary operators are evaluated is not specified. To ensure correct results, avoid creating binary expressions that depend on the order in which the compiler evaluates the operands.

As indicated in the descriptions of the operators, the usual arithmetic conversions are performed on the operands of most binary expressions.
Related information
- "Lvalues and rvalues" on page 115
- "Arithmetic conversions and promotions" on page 107

Multiplication operator *

The * (multiplication) operator yields the product of its operands. The operands must have an arithmetic or enumeration type. The result is not an lvalue. The usual arithmetic conversions on the operands are performed.

Because the multiplication operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one multiplication operator. For example, the expression:

\[ \text{sites} \times \text{number} \times \text{cost} \]

can be interpreted in any of the following ways:

\[ (\text{sites} \times \text{number}) \times \text{cost} \]
\[ \text{sites} \times (\text{number} \times \text{cost}) \]
\[ (\text{cost} \times \text{sites}) \times \text{number} \]

Division operator /

The / (division) operator yields the algebraic quotient of its operands. If both operands are integers, any fractional part (remainder) is discarded. Throwing away the fractional part is often called truncation toward zero. The operands must have an arithmetic or enumeration type. The right operand may not be zero: the result is undefined if the right operand evaluates to 0. For example, expression 7 / 4 yields the value 1 (rather than 1.75 or 2). The result is not an lvalue.

The usual arithmetic conversions on the operands are performed.

If both operands are negative, the sign of the remainder is also negative. Otherwise, the sign of the remainder is the same as the sign of the quotient.

Remainder operator %

The % (remainder) operator yields the remainder from the division of the left operand by the right operand. For example, the expression 5 % 3 yields 2. The result is not an lvalue.

Both operands must have an integral or enumeration type. If the right operand evaluates to 0, the result is undefined. If either operand has a negative value, the result is such that the following expression always yields the value of \(a\) if \(b\) is not 0 and \(a/b\) is representable:

\[ (a/b) \times b + a \%b; \]

The usual arithmetic conversions on the operands are performed.

If both operands are negative, the sign of the remainder is also negative. Otherwise, the sign of the remainder is the same as the sign of the quotient.

Addition operator +

The + (addition) operator yields the sum of its operands. Both operands must have an arithmetic type, or one operand must be a pointer to an object type and the other operand must have an integral or enumeration type.
When both operands have an arithmetic type, the usual arithmetic conversions on the operands are performed. The result has the type produced by the conversions on the operands and is not an lvalue.

A pointer to an object in an array can be added to a value having integral type. The result is a pointer of the same type as the pointer operand. The result refers to another element in the array, offset from the original element by the amount of the integral value treated as a subscript. If the resulting pointer points to storage outside the array, other than the first location outside the array, the result is undefined. A pointer to one element past the end of an array cannot be used to access the memory content at that address. The compiler does not provide boundary checking on the pointers. For example, after the addition, `ptr` points to the third element of the array:

```
int array[5];
int *ptr;
ptr = array + 2;
```

**Related information**
- "Pointer arithmetic" on page 83
- "Pointer conversions" on page 110

**Subtraction operator –**

The `−` (subtraction) operator yields the difference of its operands. Both operands must have an arithmetic or enumeration type, or the left operand must have a pointer type and the right operand must have the same pointer type or an integral or enumeration type. You cannot subtract a pointer from an integral value.

When both operands have an arithmetic type, the usual arithmetic conversions on the operands are performed. The result has the type produced by the conversions on the operands and is not an lvalue.

When the left operand is a pointer and the right operand has an integral type, the compiler converts the value of the right to an address offset. The result is a pointer of the same type as the pointer operand.

If both operands are pointers to elements in the same array, the result is the number of objects separating the two addresses. The number is of type `ptrdiff_t`, which is defined in the header file `stddef.h`. Behavior is undefined if the pointers do not refer to objects in the same array.

**Related information**
- "Pointer arithmetic" on page 83
- "Pointer conversions" on page 110

**Bitwise left and right shift operators `<<` `>>`**

The bitwise shift operators move the bit values of a binary object. The left operand specifies the value to be shifted. The right operand specifies the number of positions that the bits in the value are to be shifted. The result is not an lvalue. Both operands have the same precedence and are left-to-right associative.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&lt;&lt;</code></td>
<td>Indicates the bits are to be shifted to the left.</td>
</tr>
<tr>
<td><code>&gt;&gt;</code></td>
<td>Indicates the bits are to be shifted to the right.</td>
</tr>
</tbody>
</table>
Each operand must have an integral or enumeration type. The compiler performs integral promotions on the operands, and then the right operand is converted to type int. The result has the same type as the left operand (after the arithmetic conversions).

The right operand should not have a negative value or a value that is greater than or equal to the width in bits of the expression being shifted. The result of bitwise shifts on such values is unpredictable.

If the right operand has the value 0, the result is the value of the left operand (after the usual arithmetic conversions).

The << operator fills vacated bits with zeros. For example, if left_op has the value 4019, the bit pattern (in 16-bit format) of left_op is:
0000111110110011

The expression left_op << 3 yields:
0111110110011000

The expression left_op >> 3 yields:
00000011110110

Relational operators < > <= >=

The relational operators compare two operands and determine the validity of a relationship. The following table describes the four relational operators:

<table>
<thead>
<tr>
<th>Operator</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;</td>
<td>Indicates whether the value of the left operand is less than the value of the right operand.</td>
</tr>
<tr>
<td>&gt;</td>
<td>Indicates whether the value of the left operand is greater than the value of the right operand.</td>
</tr>
<tr>
<td>&lt;=</td>
<td>Indicates whether the value of the left operand is less than or equal to the value of the right operand.</td>
</tr>
<tr>
<td>&gt;=</td>
<td>Indicates whether the value of the left operand is greater than or equal to the value of the right operand.</td>
</tr>
</tbody>
</table>

Both operands must have arithmetic or enumeration types or be pointers to the same type.

- The type of the result is int and has the values 1 if the specified relationship is true, and 0 if false.
- The type of the result is bool and has the values true or false.

The result is not an lvalue.

If the operands have arithmetic types, the usual arithmetic conversions on the operands are performed.
When the operands are pointers, the result is determined by the locations of the objects to which the pointers refer. If the pointers do not refer to objects in the same array, the result is not defined.

A pointer can be compared to a constant expression that evaluates to 0. You can also compare a pointer to a pointer of type `void*`. The pointer is converted to a pointer of type `void*`.

If two pointers refer to the same object, they are considered equal. If two pointers refer to nonstatic members of the same object, the pointer to the object declared later is greater, provided that they are not separated by an access specifier; otherwise the comparison is undefined. If two pointers refer to data members of the same union, they have the same address value.

If two pointers refer to elements of the same array, or to the first element beyond the last element of an array, the pointer to the element with the higher subscript value is greater.

You can only compare members of the same object with relational operators.

Relational operators have left-to-right associativity. For example, the expression:

\[
a < b <= c
\]

is interpreted as:

\[
(a < b) <= c
\]

If the value of `a` is less than the value of `b`, the first relationship yields 1 in C, or `true` in C++. The compiler then compares the value `true` (or 1) with the value of `c` (integral promotions are carried out if needed).

### Equality and inequality operators == !=

The equality operators, like the relational operators, compare two operands for the validity of a relationship. The equality operators, however, have a lower precedence than the relational operators. The following table describes the two equality operators:

<table>
<thead>
<tr>
<th>Operator</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>==</code></td>
<td>Indicates whether the value of the left operand is equal to the value of the right operand.</td>
</tr>
<tr>
<td><code>!=</code></td>
<td>Indicates whether the value of the left operand is not equal to the value of the right operand.</td>
</tr>
</tbody>
</table>

Both operands must have arithmetic or enumeration types or be pointers to the same type, or one operand must have a pointer type and the other operand must be a pointer to void or a null pointer.

- **C** The type of the result is `int` and has the values 1 if the specified relationship is true, and 0 if false.  
- **C++** The type of the result is `bool` and has the values `true` or `false`.

If the operands have arithmetic types, the usual arithmetic conversions on the operands are performed.
If the operands are pointers, the result is determined by the locations of the objects to which the pointers refer.

If one operand is a pointer and the other operand is an integer having the value 0, the == expression is true only if the pointer operand evaluates to NULL. The != operator evaluates to true if the pointer operand does not evaluate to NULL.

You can also use the equality operators to compare pointers to members that are of the same type but do not belong to the same object. The following expressions contain examples of equality and relational operators:

time < max_time == status < complete
letter != EOF

**Note:** The equality operator (==) should not be confused with the assignment (=) operator.

For example,

```
if (x == 3)
   evaluates to true (or 1) if x is equal to three. Equality tests like this should be coded with spaces between the operator and the operands to prevent unintentional assignments.

while

if (x = 3)
   is taken to be true because (x = 3) evaluates to a nonzero value (3). The expression also assigns the value 3 to x.
```

**Related information**

- “Simple assignment operator=” on page 158

**Bitwise AND operator &**

The & (bitwise AND) operator compares each bit of its first operand to the corresponding bit of the second operand. If both bits are 1’s, the corresponding bit of the result is set to 1. Otherwise, it sets the corresponding result bit to 0.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands.

Because the bitwise AND operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise AND operator.

The following example shows the values of a, b, and the result of a & b represented as 16-bit binary numbers:

<table>
<thead>
<tr>
<th>bit pattern of a</th>
<th>0000000001011100</th>
</tr>
</thead>
<tbody>
<tr>
<td>bit pattern of b</td>
<td>0000000000101110</td>
</tr>
<tr>
<td>bit pattern of a &amp; b</td>
<td>0000000000011100</td>
</tr>
</tbody>
</table>

**Note:** The bitwise AND (&) should not be confused with the logical AND. (&&) operator. For example,
1 & 4 evaluates to 0
while
1 && 4 evaluates to true

Bitwise exclusive OR operator ^
The bitwise exclusive OR operator (in EBCDIC, the ^ symbol is represented by the ^ symbol) compares each bit of its first operand to the corresponding bit of the second operand. If both bits are 1’s or both bits are 0’s, the corresponding bit of the result is set to 0. Otherwise, it sets the corresponding result bit to 1.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands and is not an lvalue.

Because the bitwise exclusive OR operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise exclusive OR operator. Note that the ^ character can be represented by the trigraph ??'.

The following example shows the values of a, b, and the result of a ^ b represented as 16-bit binary numbers:

<table>
<thead>
<tr>
<th>bit pattern of a</th>
<th>0000000001011100</th>
</tr>
</thead>
<tbody>
<tr>
<td>bit pattern of b</td>
<td>0000000000101110</td>
</tr>
<tr>
<td>bit pattern of a ^ b</td>
<td>0000000001110010</td>
</tr>
</tbody>
</table>

Related information
- “Trigraph sequences” on page 34

Bitwise inclusive OR operator I
The | (bitwise inclusive OR) operator compares the values (in binary format) of each operand and yields a value whose bit pattern shows which bits in either of the operands has the value 1. If both of the bits are 0, the result of that bit is 0; otherwise, the result is 1.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands and is not an lvalue.

Because the bitwise inclusive OR operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise inclusive OR operator. Note that the | character can be represented by the trigraph ??!.

The following example shows the values of a, b, and the result of a | b represented as 16-bit binary numbers:

<table>
<thead>
<tr>
<th>bit pattern of a</th>
<th>0000000001011100</th>
</tr>
</thead>
<tbody>
<tr>
<td>bit pattern of b</td>
<td>0000000000101110</td>
</tr>
<tr>
<td>bit pattern of a</td>
<td>b</td>
</tr>
</tbody>
</table>

Note: The bitwise OR (|) should not be confused with the logical OR (||) operator. For example,
1 | 4 evaluates to 5
while
1 || 4 evaluates to true

Related information
• "Trigraph sequences" on page 34

Logical AND operator &&

The && (logical AND) operator indicates whether both operands are true.

C

If both operands have nonzero values, the result has the value 1. Otherwise, the result has the value 0. The type of the result is int. Both operands must have a arithmetic or pointer type. The usual arithmetic conversions on each operand are performed.

C++

If both operands have values of true, the result has the value true. Otherwise, the result has the value false. Both operands are implicitly converted to bool and the result type is bool.

Unlike the & (bitwise AND) operator, the && operator guarantees left-to-right evaluation of the operands. If the left operand evaluates to 0 (or false), the right operand is not evaluated.

The following examples show how the expressions that contain the logical AND operator are evaluated:

<table>
<thead>
<tr>
<th>Expression</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>1 &amp;&amp; 0</td>
<td>false or 0</td>
</tr>
<tr>
<td>1 &amp;&amp; 4</td>
<td>true or 1</td>
</tr>
<tr>
<td>0 &amp;&amp; 0</td>
<td>false or 0</td>
</tr>
</tbody>
</table>

The following example uses the logical AND operator to avoid division by zero:
(y != 0) && (x / y)

The expression x / y is not evaluated when y != 0 evaluates to 0 (or false).

Note: The logical AND (&&) should not be confused with the bitwise AND (&) operator. For example:

1 && 4 evaluates to 1 (or true)
while
1 & 4 evaluates to 0

Logical OR operator ||

The || (logical OR) operator indicates whether either operand is true.

C

If either of the operands has a nonzero value, the result has the value 1. Otherwise, the result has the value 0. The type of the result is int. Both operands must have a arithmetic or pointer type. The usual arithmetic conversions on each operand are performed.
If either operand has a value of true, the result has the value true. Otherwise, the result has the value false. Both operands are implicitly converted to bool and the result type is bool.

Unlike the | (bitwise inclusive OR) operator, the || operator guarantees left-to-right evaluation of the operands. If the left operand has a nonzero (or true) value, the right operand is not evaluated.

The following examples show how expressions that contain the logical OR operator are evaluated:

<table>
<thead>
<tr>
<th>Expression</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td></td>
</tr>
<tr>
<td>0</td>
<td></td>
</tr>
</tbody>
</table>

The following example uses the logical OR operator to conditionally increment y:

```cpp
++x || ++y;
```

The expression ++y is not evaluated when the expression ++x evaluates to a nonzero (or true) quantity.

**Note:** The logical OR (||) should not be confused with the bitwise OR (|) operator. For example:

- 1 || 4 evaluates to 1 (or true)
- 1 | 4 evaluates to 5

**Pointer to member operators .* ->* (C++ only)**

There are two pointer to member operators: .* and ->*.

The .* operator is used to dereference pointers to class members. The first operand must be of class type. If the type of the first operand is class type T, or is a class that has been derived from class type T, the second operand must be a pointer to a member of a class type T.

The ->* operator is also used to dereference pointers to class members. The first operand must be a pointer to a class type. If the type of the first operand is a pointer to class type T, or is a pointer to a class derived from class type T, the second operand must be a pointer to a member of class type T.

The .* and ->* operators bind the second operand to the first, resulting in an object or function of the type specified by the second operand.

If the result of .* or ->* is a function, you can only use the result as the operand for the () (function call) operator. If the second operand is an lvalue, the result of .* or ->* is an lvalue.

**Related information**

- “Class member lists” on page 255
Conditional expressions

A conditional expression is a compound expression that contains a condition that is implicitly converted to type bool in C++(operand₁), an expression to be evaluated if the condition evaluates to true (operand₂), and an expression to be evaluated if the condition has the value false (operand₃).

The conditional expression contains one two-part operator. The ? symbol follows the condition, and the : symbol appears between the two action expressions. All expressions that occur between the ? and : are treated as one expression.

The first operand must have a scalar type. The type of the second and third operands must be one of the following:
- An arithmetic type
- A compatible pointer, structure, or union type
- void

The second and third operands can also be a pointer or a null pointer constant.

Two objects are compatible when they have the same type but not necessarily the same type qualifiers (volatile or const). Pointer objects are compatible if they have the same type or are pointers to void.

The first operand is evaluated, and its value determines whether the second or third operand is evaluated:
- If the value is true, the second operand is evaluated.
- If the value is false, the third operand is evaluated.

The result is the value of the second or third operand.

If the second and third expressions evaluate to arithmetic types, the usual arithmetic conversions are performed on the values. The types of the second and third operands determine the type of the result as shown in the following tables.

Conditional expressions have right-to-left associativity with respect to their first and third operands. The leftmost operand is evaluated first, and then only one of the remaining two operands is evaluated. The following expressions are equivalent:

\[
a \ ? \ b : \ c \ ? \ d : \ e \ ? \ f : \ g \\
a \ ? \ b : (c \ ? \ d : (e \ ? \ f : g))
\]

Types in conditional C expressions

<table>
<thead>
<tr>
<th>Type of one operand</th>
<th>Type of other operand</th>
<th>Type of result</th>
</tr>
</thead>
<tbody>
<tr>
<td>Arithmetic</td>
<td>Arithmetic</td>
<td>Arithmetic type after usual arithmetic conversions</td>
</tr>
<tr>
<td>Structure or union type</td>
<td>Compatible structure or union type</td>
<td>Structure or union type with all the qualifiers on both operands</td>
</tr>
</tbody>
</table>

In C, a conditional expression is not an lvalue, nor is its result.

Table 21. Types of operands and results in conditional C expressions
Table 21. Types of operands and results in conditional C expressions (continued)

<table>
<thead>
<tr>
<th>Type of one operand</th>
<th>Type of other operand</th>
<th>Type of result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>void</code></td>
<td><code>void</code></td>
<td><code>void</code></td>
</tr>
<tr>
<td>Pointer to compatible type</td>
<td>Pointer to compatible type</td>
<td>Pointer to type with all the qualifiers specified for the type</td>
</tr>
<tr>
<td>Pointer to type</td>
<td>NULL pointer (the constant 0)</td>
<td>Pointer to type</td>
</tr>
<tr>
<td>Pointer to object or incomplete type</td>
<td>Pointer to void</td>
<td>Pointer to <code>void</code> with all the qualifiers specified for the type</td>
</tr>
</tbody>
</table>

IBM extension

In GNU C, a conditional expression is a valid lvalue, provided that its type is not `void` and both of its branches are valid lvalues. The following conditional expression `(a ? b : c)` is legal under GNU C:

```
(a ? b : c) = 5
/* Under GNU C, equivalent to (a ? b = 5 : (c = 5)) */
```

This extension is available when compiling in one of the extended language levels.

End of IBM extension

End of C only

Types in conditional C++ expressions

C++ only

In C++, a conditional expression is a valid lvalue if its type is not `void`, and its result is an lvalue.

Table 22. Types of operands and results in C++ conditional expressions

<table>
<thead>
<tr>
<th>Type of one operand</th>
<th>Type of other operand</th>
<th>Type of result</th>
</tr>
</thead>
<tbody>
<tr>
<td>Reference to type</td>
<td>Reference to type</td>
<td>Reference after usual reference conversions</td>
</tr>
<tr>
<td>Class T</td>
<td>Class T</td>
<td>Class T</td>
</tr>
<tr>
<td>Class T</td>
<td>Class X</td>
<td>Class type for which a conversion exists. If more than one possible conversion exists, the result is ambiguous.</td>
</tr>
<tr>
<td>throw expression</td>
<td>Other (type, pointer, reference)</td>
<td>Type of the expression that is not a throw expression</td>
</tr>
</tbody>
</table>

End of C++ only

Examples of conditional expressions

The following expression determines which variable has the greater value, `y` or `z`, and assigns the greater value to the variable `x`:

```
x = (y > z) ? y : z;
```
The following is an equivalent statement:

```
if (y > z)
  x = y;
else
  x = z;
```

The following expression calls the function `printf`, which receives the value of the variable `c`, if `c` evaluates to a digit. Otherwise, `printf` receives the character constant 'x'.

```
printf("c = %c\n", isdigit(c) ? c : 'x');
```

If the last operand of a conditional expression contains an assignment operator, use parentheses to ensure the expression evaluates properly. For example, the `=` operator has higher precedence than the `?:` operator in the following expression:

```
int i,j,k;
(i == 7) ? j ++ : k = j;
```

The compiler will interpret this expression as if it were parenthesized this way:

```
int i,j,k;
((i == 7) ? j ++ : k) = j;
```

That is, `k` is treated as the third operand, not the entire assignment expression `k = j`.

To assign the value of `j` to `k` when `i == 7` is false, enclose the last operand in parentheses:

```
int i,j,k;
(i == 7) ? j + : (k = j);
```

---

**Assignment expressions**

An *assignment expression* stores a value in the object designated by the left operand. There are two types of assignment operators:

- **Simple assignment operator** =
- **Compound assignment operators**

The left operand in all assignment expressions must be a modifiable lvalue. The type of the expression is the type of the left operand. The value of the expression is the value of the left operand after the assignment has completed.

> C The result of an assignment expression is not an lvalue.  
> C++ The result of an assignment expression is an lvalue.

All assignment operators have the same precedence and have right-to-left associativity.

**Related information**

- “Lvalues and rvalues” on page 115
- “Pointers” on page 82
- “Type qualifiers” on page 68

**Simple assignment operator** =

The simple assignment operator has the following form:

```
lvalue = expr
```
The operator stores the value of the right operand \( expr \) in the object designated by the left operand \( lvalue \).

The left operand must be a modifiable \( lvalue \). The type of an assignment operation is the type of the left operand.

If the left operand is not a class type or a vector type, the right operand is implicitly converted to the type of the left operand. This converted type will not be qualified by \texttt{const} or \texttt{volatile}.

If the left operand is a class type, that type must be complete. The copy assignment operator of the left operand will be called.

If the left operand is an object of reference type, the compiler will assign the value of the right operand to the object denoted by the reference.

The assignment operator has been extended to permit operands of vector type. Both sides of an assignment expression must be of the same vector type.

**Compound assignment operators**

The compound assignment operators consist of a binary operator and the simple assignment operator. They perform the operation of the binary operator on both operands and store the result of that operation into the left operand, which must be a modifiable \( lvalue \).

The following table shows the operand types of compound assignment expressions:

<table>
<thead>
<tr>
<th>Operator</th>
<th>Left operand</th>
<th>Right operand</th>
</tr>
</thead>
<tbody>
<tr>
<td>+= or -=</td>
<td>Arithmetic</td>
<td>Arithmetic</td>
</tr>
<tr>
<td>+= or -=</td>
<td>Pointer</td>
<td>Integral type</td>
</tr>
<tr>
<td>*=, /=, and %=</td>
<td>Arithmetic</td>
<td>Arithmetic</td>
</tr>
<tr>
<td>&lt;&lt;=, &gt;&gt;=, &amp;=, ^=, and</td>
<td>=</td>
<td>Integral type</td>
</tr>
</tbody>
</table>

Note that the expression

\[
a \ *= b + c
\]

is equivalent to

\[
a = a \ * (b + c)
\]

and \textit{not}

\[
a = a \ * b + c
\]

The following table lists the compound assignment operators and shows an expression using each operator:

<table>
<thead>
<tr>
<th>Operator</th>
<th>Example</th>
<th>Equivalent expression</th>
</tr>
</thead>
<tbody>
<tr>
<td>+=</td>
<td>index += 2</td>
<td>index = index + 2</td>
</tr>
<tr>
<td>-=</td>
<td>*(pointer++) -= 1</td>
<td>*pointer = *(pointer++) - 1</td>
</tr>
<tr>
<td>*=</td>
<td>bonus *= increase</td>
<td>bonus = bonus * increase</td>
</tr>
<tr>
<td>Operator</td>
<td>Example</td>
<td>Equivalent expression</td>
</tr>
<tr>
<td>----------</td>
<td>-----------</td>
<td>--------------------------------</td>
</tr>
<tr>
<td>/=</td>
<td>time /= hours</td>
<td>time = time / hours</td>
</tr>
<tr>
<td>%=</td>
<td>allowance %= 1000</td>
<td>allowance = allowance % 1000</td>
</tr>
<tr>
<td>&lt;&lt;=</td>
<td>result &lt;&lt;= num</td>
<td>result = result &lt;&lt;= num</td>
</tr>
<tr>
<td>&gt;&gt;=</td>
<td>form &gt;&gt;= 1</td>
<td>form = form &gt;&gt;= 1</td>
</tr>
<tr>
<td>&amp;=</td>
<td>mask &amp;= 2</td>
<td>mask = mask &amp; 2</td>
</tr>
<tr>
<td>^=</td>
<td>test ^= pre_test</td>
<td>test = test ^= pre_test</td>
</tr>
<tr>
<td></td>
<td>=</td>
<td>flag</td>
</tr>
</tbody>
</table>

Although the equivalent expression column shows the left operands (from the example column) twice, it is in effect evaluated only once.

**C++**  
In addition to the table of operand types, an expression is implicitly converted to the cv-unqualified type of the left operand if it is not of class type. However, if the left operand is of class type, the class becomes complete, and assignment to objects of the class behaves as a copy assignment operation. Compound expressions and conditional expressions are lvalues in C++, which allows them to be a left operand in a compound assignment expression.

---

**IBM extension**

When GNU C language features have been enabled, compound expressions and conditional expressions are allowed as lvalues, provided that their operands are lvalues. The following compound assignment of the compound expression (a, b) is legal under GNU C, provided that expression b, or more generally, the last expression in the sequence, is an lvalue:

```
(a, b) += 5  /* Under GNU C, this is equivalent to
            a, (b += 5)     */
```

---

**Comma expressions**

A comma expression contains two operands of any type separated by a comma and has left-to-right associativity. The left operand is fully evaluated, possibly producing side effects, and its value, if there is one, is discarded. The right operand is then evaluated. The type and value of the result of a comma expression are those of its right operand, after the usual unary conversions.

**C only**

The result of a comma expression is not an lvalue.

---

**End of C only**

**C++ only**

In C++, the result is an lvalue if the right operand is an lvalue. The following statements are equivalent:
\[ r = (a, b, \ldots, c); \]
\[ a; b; r = c; \]

The difference is that the comma operator may be suitable for expression contexts, such as loop control expressions.

Similarly, the address of a compound expression can be taken if the right operand is an lvalue.
\[ \&(a, b) \]
\[ a, \&b \]

--- End of C++ only ---

Any number of expressions separated by commas can form a single expression because the comma operator is associative. The use of the comma operator guarantees that the subexpressions will be evaluated in left-to-right order, and the value of the last becomes the value of the entire expression. In the following example, if omega has the value 11, the expression increments delta and assigns the value 3 to alpha:
\[ \text{alpha} = (\text{delta}++, \text{omega} \% 4); \]

A sequence point occurs after the evaluation of the first operand. The value of delta is discarded. Similarly, in the following example, the value of the expression:
\[ \text{intensity}++, \text{shade} \times \text{increment}, \text{rotate}(\text{direction}); \]

is the value of the expression:
\[ \text{rotate}(\text{direction}) \]

In some contexts where the comma character is used, parentheses are required to avoid ambiguity. For example, the function
\[ f(a, (t = 3, t + 2), c); \]

has only three arguments: the value of a, the value 5, and the value of c. Other contexts in which parentheses are required are in field-length expressions in structure and union declarator lists, enumeration value expressions in enumeration declarator lists, and initialization expressions in declarations and initializers.

In the previous example, the comma is used to separate the argument expressions in a function invocation. In this context, its use does not guarantee the order of evaluation (left to right) of the function arguments.

The primary use of the comma operator is to produce side effects in the following situations:

- Calling a function
- Entering or repeating an iteration loop
- Testing a condition
- Other situations where a side effect is required but the result of the expression is not immediately needed
The following table gives some examples of the uses of the comma operator.

<table>
<thead>
<tr>
<th>Statement</th>
<th>Effects</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>for (i=0; i&lt;2; ++i, f() );</code></td>
<td>A for statement in which i is incremented and f() is called at each iteration.</td>
</tr>
<tr>
<td><code>if ( f(), ++i, i&gt;1 ) { /* ... */ }</code></td>
<td>An if statement in which function f() is called, variable i is incremented, and variable i is tested against a value. The first two expressions within this comma expression are evaluated before the expression i&gt;1. Regardless of the results of the first two expressions, the third is evaluated and its result determines whether the if statement is processed.</td>
</tr>
<tr>
<td><code>func( ( ++a, f(a) ) );</code></td>
<td>A function call to func() in which a is incremented, the resulting value is passed to a function f(), and the return value of f() is passed to func(). The function func() is passed only a single argument, because the comma expression is enclosed in parentheses within the function argument list.</td>
</tr>
</tbody>
</table>

**throw expressions (C++ only)**

A throw expression is used to throw exceptions to C++ exception handlers. A throw expression is of type void.

**Related information**
- [Chapter 16, “Exception handling (C++ only),” on page 363](#)
- [“The void type” on page 53](#)

**Operator precedence and associativity**

Two operator characteristics determine how operands group with operators: *precedence* and *associativity*. Precedence is the priority for grouping different types of operators with their operands. Associativity is the left-to-right or right-to-left order for grouping operands to operators that have the same precedence. An operator’s precedence is meaningful only if other operators with higher or lower precedence are present. Expressions with higher-precedence operators are evaluated first. The grouping of operands can be forced by using parentheses.

For example, in the following statements, the value of 5 is assigned to both a and b because of the right-to-left associativity of the * operator. The value of c is assigned to b first, and then the value of b is assigned to a.

```c
b = 9;
c = 5;
a = b = c;
```

Because the order of subexpression evaluation is not specified, you can explicitly force the grouping of operands with operators by using parentheses.

In the expression

```c
a + b * c / d
```

the * and / operations are performed before + because of precedence. b is multiplied by c before it is divided by d because of associativity.
The following tables list the C and C++ language operators in order of precedence and show the direction of associativity for each operator. Operators that have the same rank have the same precedence.

**Table 23. Precedence and associativity of postfix operators**

<table>
<thead>
<tr>
<th>Rank</th>
<th>Right associative?</th>
<th>Operator function</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>yes</td>
<td>C++ global scope resolution</td>
<td>:: name_or_qualified name</td>
</tr>
<tr>
<td>1</td>
<td>yes</td>
<td>C++ class or namespace scope resolution</td>
<td>class_or_namespace :: member</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>member selection</td>
<td>object . member</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>member selection</td>
<td>pointer -&gt; member</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>subscripting</td>
<td>pointer [ expr ]</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>function call</td>
<td>expr ( expr_list )</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>value construction</td>
<td>type ( expr_list )</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>postfix increment</td>
<td>lvalue ++</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>postfix decrement</td>
<td>lvalue --</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ type identification</td>
<td>typeid ( type )</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ type identification at run time</td>
<td>typeid ( expr )</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ conversion checked at compile time</td>
<td>static_cast &lt; type &gt; ( expr )</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ conversion checked at run time</td>
<td>dynamic_cast &lt; type &gt; ( expr )</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ unchecked conversion</td>
<td>reinterpret_cast &lt; type &gt; ( expr )</td>
</tr>
<tr>
<td>2</td>
<td>yes</td>
<td>C++ const conversion</td>
<td>const_cast &lt; type &gt; ( expr )</td>
</tr>
</tbody>
</table>

**Table 24. Precedence and associativity of unary operators**

<table>
<thead>
<tr>
<th>Rank</th>
<th>Right associative?</th>
<th>Operator function</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>yes</td>
<td>size of object in bytes</td>
<td>sizeof expr</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>size of type in bytes</td>
<td>sizeof ( type )</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>prefix increment</td>
<td>++ lvalue</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>prefix decrement</td>
<td>-- lvalue</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>bitwise negation</td>
<td>~ expr</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>not</td>
<td>! expr</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>unary minus</td>
<td>- expr</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>unary plus</td>
<td>+ expr</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>address of</td>
<td>&amp; lvalue</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>indirection or dereference</td>
<td>* expr</td>
</tr>
</tbody>
</table>
Table 24. Precedence and associativity of unary operators (continued)

<table>
<thead>
<tr>
<th>Rank</th>
<th>Right associative?</th>
<th>Operator function</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>yes</td>
<td>create (allocate memory)</td>
<td>new type</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>create (allocate and initialize memory)</td>
<td>new type (expr_list) type</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>create (placement)</td>
<td>new type (expr_list) type (expr_list)</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>destroy (deallocate memory)</td>
<td>delete pointer</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>destroy array</td>
<td>delete [ ] pointer</td>
</tr>
<tr>
<td>3</td>
<td>yes</td>
<td>type conversion (cast)</td>
<td>(type) expr</td>
</tr>
</tbody>
</table>

Table 25. Precedence and associativity of binary operators

<table>
<thead>
<tr>
<th>Rank</th>
<th>Right associative?</th>
<th>Operator function</th>
<th>Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td></td>
<td>object .* ptr_to_member</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td></td>
<td>object -&gt;* ptr_to_member</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>multiplication</td>
<td>expr * expr</td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>division</td>
<td>expr / expr</td>
</tr>
<tr>
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<td></td>
<td>modulo (remainder)</td>
<td>expr % expr</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>binary addition</td>
<td>expr + expr</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>binary subtraction</td>
<td>expr - expr</td>
</tr>
<tr>
<td>7</td>
<td></td>
<td>bitwise shift left</td>
<td>expr &lt;&lt; expr</td>
</tr>
<tr>
<td>7</td>
<td></td>
<td>bitwise shift right</td>
<td>expr &gt;&gt; expr</td>
</tr>
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<td>8</td>
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<td>less than or equal to</td>
<td>expr &lt;= expr</td>
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<td>greater than</td>
<td>expr &gt; expr</td>
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<td></td>
<td>greater than or equal to</td>
<td>expr &gt;= expr</td>
</tr>
<tr>
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<td></td>
<td>equal</td>
<td>expr == expr</td>
</tr>
<tr>
<td>9</td>
<td></td>
<td>not equal</td>
<td>expr != expr</td>
</tr>
<tr>
<td>10</td>
<td></td>
<td>bitwise AND</td>
<td>expr &amp; expr</td>
</tr>
<tr>
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<td></td>
<td>bitwise exclusive OR</td>
<td>expr ^ expr</td>
</tr>
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<td>12</td>
<td></td>
<td>bitwise inclusive OR</td>
<td>expr</td>
</tr>
<tr>
<td>13</td>
<td></td>
<td>logical AND</td>
<td>expr &amp;&amp; expr</td>
</tr>
<tr>
<td>14</td>
<td></td>
<td>logical inclusive OR</td>
<td>expr</td>
</tr>
<tr>
<td>15</td>
<td></td>
<td>conditional expression</td>
<td>expr ? expr : expr</td>
</tr>
<tr>
<td>16</td>
<td></td>
<td>simple assignment</td>
<td>lvalue = expr</td>
</tr>
<tr>
<td>16</td>
<td></td>
<td>multiply and assign</td>
<td>lvalue *= expr</td>
</tr>
<tr>
<td>16</td>
<td></td>
<td>divide and assign</td>
<td>lvalue /= expr</td>
</tr>
</tbody>
</table>
Examples of expressions and precedence

The parentheses in the following expressions explicitly show how the compiler groups operands and operators.

```
total = (4 + (5 * 3));
total = (((8 * 5) / 10) / 3);
total = (10 + (5/3));
```

If parentheses did not appear in these expressions, the operands and operators would be grouped in the same manner as indicated by the parentheses. For example, the following expressions produce the same output.

```
total = (4+(5*3));
total = 4+5*3;
```

Because the order of grouping operands with operators that are both associative and commutative is not specified, the compiler can group the operands and operators in the expression:

```
total = price + prov_tax + city_tax;
```

in the following ways (as indicated by parentheses):

```
total = (price + (prov_tax + city_tax));
total = ((price + prov_tax) + city_tax);
total = ((price + city_tax) + prov_tax);
```

The grouping of operands and operators does not affect the result unless one ordering causes an overflow and another does not. For example, if `price = 32767`, `prov_tax = -42`, and `city_tax = 32767`, and all three of these variables have been declared as integers, the third statement `total = ((price + city_tax) + prov_tax)` will cause an integer overflow and the rest will not.

Because intermediate values are rounded, different groupings of floating-point operators may give different results.
In certain expressions, the grouping of operands and operators can affect the result. For example, in the following expression, each function call might be modifying the same global variables.

```c
a = b() + c() + d();
```

This expression can give different results depending on the order in which the functions are called.

If the expression contains operators that are both associative and commutative and the order of grouping operands with operators can affect the result of the expression, separate the expression into several expressions. For example, the following expressions could replace the previous expression if the called functions do not produce any side effects that affect the variable `a`.

```c
a = b();
a += c();
a += d();
```

The order of evaluation for function call arguments or for the operands of binary operators is not specified. Therefore, the following expressions are ambiguous:

```c
z = (x * ++y) / func1(y);
func2(++i, x[i]);
```

If `y` has the value of 1 before the first statement, it is not known whether or not the value of 1 or 2 is passed to `func1()`. In the second statement, if `i` has the value of 1 before the expression is evaluated, it is not known whether `x[1]` or `x[2]` is passed as the second argument to `func2()`.
Chapter 7. Statements

A statement, the smallest independent computational unit, specifies an action to be performed. In most cases, statements are executed in sequence. The following is a summary of the statements available in C and C++:

- Labeled statements
- Expression statements
- Block statements
- Selection statements
- Iteration statements
- Jump statements
- Declaration statements
- C++ try blocks
- Null statement
- IBM Inline assembly statements

Related information
- Chapter 3, “Data objects and declarations,” on page 39
- “Function declarations” on page 192
- “try blocks” on page 363

Labeled statements

There are three kinds of labels: identifier, case, and default.

Labeled statement syntax

```
identifier: statement
```

The label consists of the identifier and the colon (:) character.

A label name must be unique within the function in which it appears.

In C++, an identifier label may only be used as the target of a goto statement. A goto statement can use a label before its definition. Identifier labels have their own namespace; you do not have to worry about identifier labels conflicting with other identifiers. However, you may not redeclare a label within a function.

Case and default label statements only appear in switch statements. These labels are accessible only within the closest enclosing switch statement.

case statement syntax

```
case constant_expression: statement
```

default statement syntax

The following are examples of labels:

```
comment_complete : ;  /* null statement label */
test_for_null : if (NULL == pointer)
```

Related information

- "The goto statement" on page 184
- "The switch statement" on page 173

Locally declared labels

A locally declared label, or local label, is an identifier label that is declared at the beginning of a statement expression and for which the scope is the statement expression in which it is declared and defined. This language feature is an extension of C and C++ to facilitate handling programs developed with GNU C.

A local label can be used as the target of a goto statement, jumping to it from within the same block in which it was declared. This language extension is particularly useful for writing macros that contain nested loops, capitalizing on the difference between its statement scope and the function scope of an ordinary label.

Locally declared label syntax

```
[label] identifier
```

In a statement expression, the declaration of a local label must appear immediately after the left parenthesis and left brace, and must precede any ordinary declarations and statements. The label is defined in the usual way, with a name and a colon, within the statements of the statement expression.

Related information

- "Statement expressions" on page 171

Labels as values

The address of a label defined in the current function or a containing function can be obtained and used as a value wherever a constant of type void* is valid. The address is the return value when the label is the operand of the unary operator &. The ability to use the address of label as a value is an extension to C99 and C++, implemented to facilitate porting programs developed with GNU C.
In the following example, the computed goto statements use the values of \texttt{label1} and \texttt{label2} to jump to those spots in the function.

```c
int main()
{
    void * ptr1, *ptr2;
    ...
    label1: ...
    ...
    label2: ...
    ...
    ptr1 = &&label1;
    ptr2 = &&label2;
    if (...) {
        goto *ptr1;
    }
    else {
        goto *ptr2;
    }
    ...
}
```

Related information
- “Label value operator &&” on page 140
- “Computed goto statement” on page 185

--- End of IBM extension ---

**Expression statements**

An expression statement contains an expression. The expression can be null.

**Expression statement syntax**

```
expression
```

An expression statement evaluates \texttt{expression}, then discards the value of the expression. An expression statement without an expression is a null statement.

The following are examples of statements:

```c
printf("Account Number: \n"); /* call to the printf */
marks = dollars * exch_rate; /* assignment to marks */
(difference < 0) ? ++losses : ++gain; /* conditional increment */
```

Related information
- Chapter 6, “Expressions and operators,” on page 115

**Resolution of ambiguous statements**

The C++ syntax does not disambiguate between expression statements and declaration statements. The ambiguity arises when an expression statement has a function-style cast as its left-most subexpression. (Note that, because C does not support function-style casts, this ambiguity does not occur in C programs.) If the statement can be interpreted both as a declaration and as an expression, the statement is interpreted as a declaration statement.
Note: The ambiguity is resolved only on a syntactic level. The disambiguation does not use the meaning of the names, except to assess whether or not they are type names.

The following expressions disambiguate into expression statements because the ambiguous subexpression is followed by an assignment or an operator. type_spec in the expressions can be any type specifier:

```
type_spec(i)++;  // expression statement
type_spec(i,3)<<d;  // expression statement
type_spec(i)->l=24;  // expression statement
```

In the following examples, the ambiguity cannot be resolved syntactically, and the statements are interpreted as declarations. type_spec is any type specifier:

```
type_spec(*i)(int);  // declaration
type_spec(j)[5];  // declaration
type_spec(m) = { 1, 2 };  // declaration
```

The last statement above causes a compile-time error because you cannot initialize a pointer with a float value.

Any ambiguous statement that is not resolved by the above rules is by default a declaration statement. All of the following are declaration statements:

```
type_spec(a);  // declaration
type_spec(b)();  // declaration
type_spec(c)=23;  // declaration
type_spec(d),e,f,g=0;  // declaration
type_spec(h)(e,3);  // declaration
```

Related information
- Chapter 3, “Data objects and declarations,” on page 39
- Chapter 6, “Expressions and operators,” on page 115
- “Function call operator ( )” on page 121

---

**Block statements**

A block statement, or compound statement, lets you group any number of data definitions, declarations, and statements into one statement. All definitions, declarations, and statements enclosed within a single set of braces are treated as a single statement. You can use a block wherever a single statement is allowed.

**Block statement syntax**

```
{  
  type_definition
  file_scope_data_declaration
  block_scope_data_declaration
  ...  
  statement
}
```

A block defines a local scope. If a data object is usable within a block and its identifier is not redefined, all nested blocks can use that data object.
Example of blocks

The following program shows how the values of data objects change in nested blocks:

```c
/**
 * This example shows how data objects change in nested blocks.
**/
#include <stdio.h>

int main(void)
{
    int x = 1; /* Initialize x to 1 */
    int y = 3;
    if (y > 0)
    {
        int x = 2; /* Initialize x to 2 */
        printf("second x = %4d\n", x);
    }
    printf("first x = %4d\n", x);
    return(0);
}
```

The program produces the following output:

second x = 2
first x = 1

Two variables named x are defined in main. The first definition of x retains storage while main is running. However, because the second definition of x occurs within a nested block, `printf("second x = %4d\n", x);` recognizes x as the variable defined on the previous line. Because `printf("first x = %4d\n", x);` is not part of the nested block, x is recognized as the first definition of x.

Statement expressions

IBM extension

A compound statement is a sequence of statements enclosed by braces. In GNU C, a compound statement inside parentheses may appear as an expression in what is called a statement expression.

Statement expression syntax

The value of a statement expression is the value of the last simple expression to appear in the entire construct. If the last statement is not an expression, then the construct is of type `void` and has no value.

The statement expression can be combined with the `typeof` operator to create complex function-like macros in which each operand is evaluated only once. For example:

```c
#define SWAP(a,b) ( { __typeof__(a) temp; temp=a; a=b; b=temp; } )
```

End of IBM extension
Selection statements

Selection statements consist of the following types of statements:

- The if statement
- The switch statement

The if statement

An if statement is a selection statement that allows more than one possible flow of control.

C++ An if statement lets you conditionally process a statement when the specified test expression, implicitly converted to bool, evaluates to true. If the implicit conversion to bool fails the program is ill-formed.

C In C, an if statement lets you conditionally process a statement when the specified test expression evaluates to a nonzero value. The test expression must be of arithmetic or pointer type.

You can optionally specify an else clause on the if statement. If the test expression evaluates to false (or in C, a zero value) and an else clause exists, the statement associated with the else clause runs. If the test expression evaluates to true, the statement following the expression runs and the else clause is ignored.

if statement syntax

```
if (expression) statement
```

When if statements are nested and else clauses are present, a given else is associated with the closest preceding if statement within the same block.

A single statement following any selection statements (if, switch) is treated as a compound statement containing the original statement. As a result any variables declared on that statement will be out of scope after the if statement. For example:

```c
if (x)
    int i;
```

is equivalent to:

```c
if (x)
    { int i; }
```

Variable i is visible only within the if statement. The same rule applies to the else part of the if statement.

Examples of if statements

The following example causes grade to receive the value A if the value of score is greater than or equal to 90.

```c
if (score >= 90)
    grade = 'A';
```

The following example displays Number is positive if the value of number is greater than or equal to 0. If the value of number is less than 0, it displays Number is negative.
if (number >= 0)
  printf("Number is positive\n");
else
  printf("Number is negative\n");

The following example shows a nested if statement:
if (paygrade == 7)
  if (level >= 0 && level <= 8)
    salary *= 1.05;
  else
    salary *= 1.04;
else
  salary *= 1.06;
cout << "salary is " << salary << endl;

The following example shows a nested if statement that does not have an else clause. Because an else clause always associates with the closest if statement, braces might be needed to force a particular else clause to associate with the correct if statement. In this example, omitting the braces would cause the else clause to associate with the nested if statement.
if (kegs > 0) {
  if (furlongs > kegs)
    fxph = furlongs/kegs;
} else
  fxph = 0;

The following example shows an if statement nested within an else clause. This example tests multiple conditions. The tests are made in order of their appearance. If one test evaluates to a nonzero value, a statement runs and the entire if statement ends.
if (value > 0)
  ++increase;
else if (value == 0)
  ++break_even;
else
  ++decrease;

Related information
• "Boolean types” on page 50

The switch statement

A switch statement is a selection statement that lets you transfer control to different statements within the switch body depending on the value of the switch expression. The switch expression must evaluate to an integral or enumeration value. The body of the switch statement contains case clauses that consist of
• A case label
• An optional default label
• A case expression
• A list of statements.

If the value of the switch expression equals the value of one of the case expressions, the statements following that case expression are processed. If not, the default label statements, if any, are processed.
**switch statement syntax**

```
switch(expression) switch_body
```

The *switch body* is enclosed in braces and can contain definitions, declarations, *case clauses*, and a *default clause*. Each case clause and default clause can contain statements.

```
{ 
  type_definition
  file_scope_data_declaration
  block_scope_data_declaration 
  case_clause 
}
```

**Note:** An initializer within a *type_definition, file_scope_data_declaration* or *block_scope_data_declaration* is ignored.

A *case clause* contains a *case label* followed by any number of statements. A case clause has the form:

**Case clause syntax**

```
case_label statement
```

A *case label* contains the word case followed by an integral constant expression and a colon. The value of each integral constant expression must represent a different value; you cannot have duplicate case labels. Anywhere you can put one case label, you can put multiple case labels. A case label has the form:

**case label syntax**

```
case integral_constant_expression : 
```

A *default clause* contains a default label followed by one or more statements. You can put a case label on either side of the default label. A *switch statement* can have only one default label. A *default_clause* has the form:

**Default clause statement**

```
default case_label : 
```
The switch statement passes control to the statement following one of the labels or to the statement following the switch body. The value of the expression that precedes the switch body determines which statement receives control. This expression is called the switch expression.

The value of the switch expression is compared with the value of the expression in each case label. If a matching value is found, control is passed to the statement following the case label that contains the matching value. If there is no matching value but there is a default label in the switch body, control passes to the default labelled statement. If no matching value is found, and there is no default label anywhere in the switch body, no part of the switch body is processed.

When control passes to a statement in the switch body, control only leaves the switch body when a break statement is encountered or the last statement in the switch body is processed.

If necessary, an integral promotion is performed on the controlling expression, and all expressions in the case statements are converted to the same type as the controlling expression. The switch expression can also be of class type if there is a single conversion to integral or enumeration type.

Compiling with option -qinfo=gen finds case labels that fall through when they should not.

Restrictions on switch statements
You can put data definitions at the beginning of the switch body, but the compiler does not initialize auto and register variables at the beginning of a switch body. You can have declarations in the body of the switch statement.

You cannot use a switch statement to jump over initializations.

When the scope of an identifier with a variably modified type includes a case or default label of a switch statement, the entire switch statement is considered to be within the scope of that identifier. That is, the declaration of the identifier must precede the switch statement.

In C++, you cannot transfer control over a declaration containing an explicit or implicit initializer unless the declaration is located in an inner block that is completely bypassed by the transfer of control. All declarations within the body of a switch statement that contain initializers must be contained in an inner block.

Examples of switch statements
The following switch statement contains several case clauses and one default clause. Each clause contains a function call and a break statement. The break statements prevent control from passing down through each statement in the switch body.

If the switch expression evaluated to '/', the switch statement would call the function divide. Control would then pass to the statement following the switch body.

```c
char key;

printf("Enter an arithmetic operator\n");
scanf("%c", &key);

switch (key)
{
```
case '+':
    add();
    break;

case '-':
    subtract();
    break;

case '*':
    multiply();
    break;

case '/':
    divide();
    break;

default:
    printf("invalid key\n");
    break;
}

If the switch expression matches a case expression, the statements following the case expression are processed until a break statement is encountered or the end of the switch body is reached. In the following example, break statements are not present. If the value of text[i] is equal to 'A', all three counters are incremented. If the value of text[i] is equal to 'a', lettera and total are increased. Only total is increased if text[i] is not equal to 'A' or 'a'.

char text[100];
int capa, lettera, total;

// ...
for (i=0; i<sizeof(text); i++) {
    switch (text[i])
    {
    case 'A':
        capa++;
    case 'a':
        lettera++;
    default:
        total++;
    }
}

The following switch statement performs the same statements for more than one case label:

/**
 ** This example contains a switch statement that performs
 ** the same statement for more than one case label.
 **/

#include <stdio.h>

int main(void)
{
    int month;

    /* Read in a month value */
    printf("Enter month: ");
    scanf("%d", &month);

    /* Tell what season it falls into */
    switch (month)
{  
case 12:
  case 1:
  case 2:
    printf("month %d is a winter month\n", month);
    break;
  case 3:
  case 4:
  case 5:
    printf("month %d is a spring month\n", month);
    break;
  case 6:
  case 7:
  case 8:
    printf("month %d is a summer month\n", month);
    break;
  case 9:
  case 10:
  case 11:
    printf("month %d is a fall month\n", month);
    break;
  case 66:
  case 99:
  default:
    printf("month %d is not a valid month\n", month);
}

return(0);
}

If the expression month has the value 3, control passes to the statement:
printf("month %d is a spring month\n", month);

The break statement passes control to the statement following the switch body.

Related information
  - __qinfo=gen__ in the XL C/C++ Compiler Reference
  - “Case and Default Labels” on page 167
  - “The break statement” on page 181

Iteration statements

Iteration statements consist of the following types of statements:
  - The while statement
  - The do statement
  - The for statement

Related information
  - “Boolean types” on page 50

The while statement

A while statement repeatedly runs the body of a loop until the controlling expression evaluates to false (or 0 in C).
while statement syntax

```
while (expression) statement
```

- **C** The expression must be of arithmetic or pointer type.
- **C++** The expression must be convertible to bool.

The expression is evaluated to determine whether or not to process the body of the loop. If the expression evaluates to false, the body of the loop never runs. If the expression does not evaluate to false, the loop body is processed. After the body has run, control passes back to the expression. Further processing depends on the value of the condition.

A break, return, or goto statement can cause a while statement to end, even when the condition does not evaluate to false.

- **C++** A throw expression also can cause a while statement to end prior to the condition being evaluated.

In the following example, item[index] triples and is printed out, as long as the value of the expression ++index is less than MAX_INDEX. When ++index evaluates to MAX_INDEX, the while statement ends.

```c
/**
 * This example illustrates the while statement.
 */

#define MAX_INDEX (sizeof(item) / sizeof(item[0]))
#include <stdio.h>

int main(void)
{
    static int item[] = { 12, 55, 62, 85, 102 };
    int index = 0;

    while (index < MAX_INDEX)
    {
        item[index] *= 3;
        printf("item[%d] = %d\n", index, item[index]);
        ++index;
    }

    return(0);
}
```

The do statement

A do statement repeatedly runs a statement until the test expression evaluates to false (or 0 in C). Because of the order of processing, the statement is run at least once.

do statement syntax

```
do statement while (expression);
```

- **C** The expression must be of arithmetic or pointer type.
- **C++** The controlling expression must be convertible to type bool.
The body of the loop is run before the controlling while clause is evaluated. Further processing of the do statement depends on the value of the while clause. If the while clause does not evaluate to false, the statement runs again. When the while clause evaluates to false, the statement ends.

A break, return, or goto statement can cause the processing of a do statement to end, even when the while clause does not evaluate to false.

A throw expression also can cause a while statement to end prior to the condition being evaluated.

The following example keeps incrementing i while i is less than 5:

```c
#include <stdio.h>

int main(void) {
    int i = 0;
    do {
        i++;
        printf("Value of i: %d\n", i);
    } while (i < 5);
    return 0;
}
```

The following is the output of the above example:

Value of i: 1
Value of i: 2
Value of i: 3
Value of i: 4
Value of i: 5

The for statement

A for statement lets you do the following:

- Evaluate an expression before the first iteration of the statement (initialization)
- Specify an expression to determine whether or not the statement should be processed (the condition)
- Evaluate an expression after each iteration of the statement (often used to increment for each iteration)
- Repeatedly process the statement if the controlling part does not evaluate to false (or 0 in C).

**for statement syntax**

```
for (expression1; expression2; expression3) statement
```

expression1 is the initialization expression. It is evaluated only before the statement is processed for the first time. You can use this expression to initialize a variable. You can also use this expression to declare a variable, provided that the variable is not declared as static (it must be automatic and may also be declared as register). If you declare a variable in this expression, or anywhere else in statement, that variable goes out of scope at the end of the for loop. If you do not want to evaluate an expression prior to the first iteration of the statement, you can omit this expression.
expression2 is the conditional expression. It is evaluated before each iteration of the statement. expression2 must be of arithmetic or pointer type. expression3 must be convertible to type bool.

If it evaluates to false (or 0 in C), the statement is not processed and control moves to the next statement following the for statement. If expression2 does not evaluate to false, the statement is processed. If you omit expression2, it is as if the expression had been replaced by true, and the for statement is not terminated by failure of this condition.

expression3 is evaluated after each iteration of the statement. This expression is often used for incrementing, decrementing, or assigning to a variable. This expression is optional.

A break, return, or goto statement can cause a for statement to end, even when the second expression does not evaluate to false. If you omit expression2, you must use a break, return, or goto statement to end the for statement.

You can set a compiler option that allows a variable declared in the scope of a for statement to have a scope that is not local to the for statement.

Related information
• -qlanglvl=noansifor in the XL C/C++ Compiler Reference

Examples of for statements
The following for statement prints the value of count 20 times. The for statement initially sets the value of count to 1. After each iteration of the statement, count is incremented.

```c
int count;
for (count = 1; count <= 20; count++)
    printf("count = %d\n", count);
```

The following sequence of statements accomplishes the same task. Note the use of the while statement instead of the for statement.

```c
int count = 1;
while (count <= 20)
{
    printf("count = %d\n", count);
    count++;
}
```

The following for statement does not contain an initialization expression:

```c
for (; index > 10; --index)
{
    list[index] = var1 + var2;
    printf("list[%d] = %d\n", index, list[index]);
}
```

The following for statement will continue running until scanf receives the letter e:

```c
for (;;)
{
    scanf("%c", &letter);
    if (letter == 'n')
        continue;
```
if (letter == 'e')
    break;
printf("You entered the letter %c
", letter);
}

The following for statement contains multiple initializations and increments. The comma operator makes this construction possible. The first comma in the for expression is a punctuator for a declaration. It declares and initializes two integers, i and j. The second comma, a comma operator, allows both i and j to be incremented at each step through the loop.
for (int i = 0, j = 50; i < 10; ++i, j += 50)
{
    cout << "i = " << i << " and j = " << j
    << endl;
}

The following example shows a nested for statement. It prints the values of an array having the dimensions [5][3].
for (row = 0; row < 5; row++)
    for (column = 0; column < 3; column++)
        printf("%d\n",
            table[row][column]);

The outer statement is processed as long as the value of row is less than 5. Each time the outer for statement is executed, the inner for statement sets the initial value of column to zero and the statement of the inner for statement is executed 3 times. The inner statement is executed as long as the value of column is less than 3.

---

**Jump statements**

Jump statements consist of the following types of statements:

- **The break statement**
- **The continue statement**
- **The return statement**
- **The goto statement**
- **Computed goto statement**

**The break statement**

A **break statement** lets you end an **iterative** (do, for, or while) statement or a **switch** statement and exit from it at any point other than the logical end. A break may only appear on one of these statements.

**break statement syntax**

```
break;
```

In an iterative statement, the break statement ends the loop and moves control to the next statement outside the loop. Within nested statements, the break statement ends only the smallest enclosing do, for, switch, or while statement.

In a **switch** statement, the break passes control out of the switch body to the next statement outside the switch statement.
The continue statement

A continue statement ends the current iteration of a loop. Program control is passed from the continue statement to the end of the loop body.

A continue statement has the form:

```
continue;
```

A continue statement can only appear within the body of an iterative statement, such as do, for, or while.

The continue statement ends the processing of the action part of an iterative statement and moves control to the loop continuation portion of the statement. For example, if the iterative statement is a for statement, control moves to the third expression in the condition part of the statement, then to the second expression (the test) in the condition part of the statement.

Within nested statements, the continue statement ends only the current iteration of the do, for, or while statement immediately enclosing it.

Examples of continue statements

The following example shows a continue statement in a for statement. The continue statement causes processing to skip over those elements of the array rates that have values less than or equal to 1.

```c
/**
 ** This example shows a continue statement in a for statement.
 **/

#include <stdio.h>
define SIZE 5

int main(void)
{
    int i;
    static float rates[SIZE] = { 1.45, 0.05, 1.88, 2.00, 0.75 };

    printf("Rates over 1.00\n");
    for (i = 0; i < SIZE; i++)
    {
        if (rates[i] <= 1.00) /* skip rates <= 1.00 */
            continue;
        printf("rate = %.2f\n", rates[i]);
    }

    return(0);
}
```

The program produces the following output:
Rates over 1.00
rate = 1.45
rate = 1.88
rate = 2.00

The following example shows a continue statement in a nested loop. When the inner loop encounters a number in the array strings, that iteration of the loop ends. Processing continues with the third expression of the inner loop. The inner loop ends when the ‘\0’ escape sequence is encountered.

The inner loop ends when the ‘\0’ escape sequence is encountered.
/**
 ** This program counts the characters in strings that are part
 ** of an array of pointers to characters. The count excludes
 ** the digits 0 through 9.
 ***/

#include <stdio.h>
#define SIZE 3

int main(void)
{
    static char *strings[SIZE] = { "ab", "c5d", "e5" };
    int i;
    int letter_count = 0;
    char *pointer;
    for (i = 0; i < SIZE; i++) /* for each string */
        /* for each character */
        for (pointer = strings[i]; *pointer != '\0';
            ++pointer)
            /* if a number */
            if (*pointer >= '0' && *pointer <= '9')
                continue;
            letter_count++;
    printf("letter count = %d\n", letter_count);
    return(0);
}

The program produces the following output:
letter count = 5

The return statement

A return statement ends the processing of the current function and returns control
to the caller of the function.

return statement syntax

\[
\text{return} \quad \left\langle \begin{array}{l}
\text{expression} \\
\end{array} \right\rangle ;
\]

A value-returning function should include a return statement, containing an
expression. If an expression is not given on a return statement in a
function declared with a non-void return type, the compiler issues a warning message. If an expression is not given on a return statement in a
function declared with a non-void return type, the compiler issues an error message.

If the data type of the expression is different from the function return type,
conversion of the return value takes place as if the value of the expression were
assigned to an object with the same function return type.

For a function of return type void, a return statement is not strictly necessary. If
the end of such a function is reached without encountering a return statement,
control is passed to the caller as if a return statement without an expression were
encountered. In other words, an implicit return takes place upon completion of the
final statement, and control automatically returns to the calling function. If a return statement is used, it must not contain an expression.
Examples of return statements

The following are examples of return statements:

```
return;            /* Returns no value */
return result;     /* Returns the value of result */
return 1;          /* Returns the value 1 */
return (x * x);    /* Returns the value of x * x */
```

The following function searches through an array of integers to determine if a match exists for the variable number. If a match exists, the function match returns the value of i. If a match does not exist, the function match returns the value -1 (negative one).

```
int match(int number, int array[ ], int n)
{
    int i;
    for (i = 0; i < n; i++)
        if (number == array[i])
            return (i);
    return(-1);
}
```

A function can contain multiple return statements. For example:

```
void copy( int *a, int *b, int c)
{
    /* Copy array a into b, assuming both arrays are the same size */
    if (!a || !b)       /* if either pointer is 0, return */
        return;
    if (a == b)         /* if both parameters refer */
        return;          /* to same array, return */
    if (c == 0)         /* nothing to copy */
        return;
    for (int i = 0; i < c; ++i;) /* do the copying */
        b[i] = a[i];     /* implicit return */
}
```

In this example, the return statement is used to cause a premature termination of the function, similar to a break statement.

An expression appearing in a return statement is converted to the return type of the function in which the statement appears. If no implicit conversion is possible, the return statement is invalid.

Related information

- "Function return type specifiers" on page 201
- "Function return values" on page 202

The goto statement

A goto statement causes your program to unconditionally transfer control to the statement associated with the label specified on the goto statement.

```
goto statement syntax
```

184  XL C/C++ Language Reference
Because the goto statement can interfere with the normal sequence of processing, it makes a program more difficult to read and maintain. Often, a break statement, a continue statement, or a function call can eliminate the need for a goto statement.

If an active block is exited using a goto statement, any local variables are destroyed when control is transferred from that block.

You cannot use a goto statement to jump over initializations.

A goto statement is allowed to jump within the scope of a variable length array, but not past any declarations of objects with variably modified types.

The following example shows a goto statement that is used to jump out of a nested loop. This function could be written without using a goto statement.

```c
/**
 * This example shows a goto statement that is used to jump out of a nested loop.
 */
#include <stdio.h>
void display(int matrix[3][3]);
int main(void)
{
    int matrix[3][3]={1,2,3,4,5,2,8,9,10};
    display(matrix);
    return(0);
}
void display(int matrix[3][3])
{
    int i, j;
    for (i = 0; i < 3; i++)
        for (j = 0; j < 3; j++)
            if ( (matrix[i][j] < 1) || (matrix[i][j] > 6) )
            goto out_of_bounds;
    printf("matrix[%d][%d] = %d\n", i, j, matrix[i][j]);
    return;
out_of_bounds: printf("number must be 1 through 6\n");
}
```

Related information

- "Labeled statements" on page 167

**Computed goto statement**

A computed goto is a goto statement for which the target is a label from the same function. The address of the label is a constant of type void*, and is obtained by applying the unary label value operator & to the label. The target of a computed goto is known at run time, and all computed goto statements from the same function will have the same targets. The language feature is an extension to C99 and C++, implemented to facilitate porting programs developed with GNU C.
Computed goto statement syntax

\[ \text{goto} \text{*expression*}; \]

The *expression* is an expression of type void*.

Related information

- “Labels as values” on page 168
- “Label value operator &&” on page 140

---

End of IBM extension

---

Null statement

The null statement performs no operation. It has the form:

\[ \text{;} \]

A null statement can hold the label of a labeled statement or complete the syntax of an iterative statement.

The following example initializes the elements of the array price. Because the initializations occur within the for expressions, a statement is only needed to finish the for syntax; no operations are required.

```c
for (i = 0; i < 3; price[i++] = 0) 
    ;
```

A null statement can be used when a label is needed before the end of a block statement. For example:

```c
void func(void) {
    if (error_detected)
        goto depart;
    /* further processing */
    depart: ; /* null statement required */
}
```

Inline assembly statements

Under extended language levels, the compiler provides partial support for embedded assembly code fragments among C and C++ source statements. This extension has been implemented for use in general system programming code, in the operating system kernel and device drivers, which were originally developed with GNU C.

The keyword `asm` stands for assembly code. When strict language levels are used in compilation, the compiler recognizes and ignores the keyword `asm` in a declaration.

The syntax is as follows:
asm statement syntax

\[
\text{asm} \quad \text{---} \quad \text{volatile} \quad \text{---} \quad \text{code_format_string} \quad \text{---} \quad \text{constraint} \quad \text{---} \quad \text{constraint} \quad \text{---} \quad \text{clobbers}
\]

input:

\[
\text{constraint} \quad \text{---} \quad \text{constraint}
\]

output:

\[
\text{constraint} \quad \text{---} \quad \text{constraint}
\]

The qualifier \texttt{volatile} instructs the compiler that the assembler instructions may update memory not listed in \texttt{output}, \texttt{input}, or \texttt{clobbers}.

The \texttt{code_format_string} is the source text of the \texttt{asm} instructions and is a string literal similar to a \texttt{printf} format specifier. The string consists of a comma-separated list of \% specifiers, each of which corresponds to an input or output operand. The \% specifiers can take either of the following forms:

- \%\texttt{integer}, where \texttt{integer} refers to the sequential number of the input or output operand.
- \%\texttt{[symbolic\_name]}, where the \texttt{symbolic\_name} can be referenced within the assembler code.

The \texttt{input} is a comma-separated list of input operands.

The \texttt{output} is a comma-separated list of output operands.

If pointer expressions are used in \texttt{input} or \texttt{output}, the assembly instructions should honor the ANSI aliasing rule (see "Type-based aliasing" on page 84 for more information). This means that indirect addressing using values in pointer expression operands should be consistent with the pointer types; otherwise, you must disable the \texttt{-qalias=ansi} option during compilation.

\texttt{clobbers} is a comma-separated list of register names enclosed in double quotes. These are registers that can be updated by the \texttt{asm} instruction.

The \texttt{constraint} is a string literal specifying the constraints for the operand, one character per constraint.

The \texttt{C\_expression} is a C or C++ expression whose value is used as the operand for the \texttt{asm} instruction. Output operands must be modifiable lvalues.

The following constraints are supported:
= Write-only operand.
+ Read and write operand.
& An operand may be modified before the instruction is finished using the input operands; a register that is used as input should not be reused here.
b Use a general register other than zero.
f Use a floating-point register.
g Use a general register, memory, or immediate operand.
i An immediate integer operand.
m A memory operand supported by the machine.
n Handle in the same way as i.
o Handle in the same way as m.
r Use a general register.
v Use a vector register.
0, 1, 2, ...66 A matching constraint. Allocate the same register in output as in the corresponding input.
I, J, K, M, N, O, P, G, S, T Constant values. Fold the expression in the operand and substitute the value into the % specifier.

Related information
• [qalias=ansi] in the XL C/C++ Compiler Reference

Examples of inline assembly statements
In the following example:
```c
int a;
int b = 100;
int c = 200;
asm("add \%0, \%1, \%2"
     : "=r"(a)
     : "r"(b), "r"(c)
);   
```
add is the op code of the instruction, understood by the assembler. \%0, \%1 and \%2 are the operands, which are to be substituted by the C expressions in the output/input operand fields. The output operand uses the = constraint to indicate that a modifiable operand is required; and the r constraint to indicate that a general purpose register is required. Likewise, the r in the input operands indicates that general purpose registers are required. Within these restrictions, the compiler is free to choose any registers to substitute for \%0, \%1, and \%2.

The following example illustrates the use of the symbolic names for input and output operands:
```c
int a;
int b = 1, c = 2, d = 3;
_asm("addc [%result],%[first],%[second]"
     :[result]="r"(a)
     :[first]="r"(b), [second]="r"(d) );
```
Restrictions on inline assembly statements

The assembler instructions must be self-contained within an `asm` statement. The `asm` statement can only be used to generate instructions. All connections to the rest of the program must be established through the `output` and `input` operand list. In particular:

- No more than 10 operands can be used in the assembly instructions within one inline assembly statement (operands that are listed but not used do not count).
- Branching to a label in another `asm` statement is not supported.
- Referencing an external symbol directly, without going through the operand list, is not supported.
- Pseudo-operators and directives, such as instructions with the suffix `.section`, `.text`, or `.data`, are not supported.
- The total number of instructions in one `asm` statement cannot exceed 63. The instruction count must also include the instructions generated by the compiler to handle the operands in the operand list.

Related information

- “Assembly labels” on page 17
- “Global variables in specified registers (C only)” on page 48
- `-qasm` and `-qasm_as` in the [XL C/C++ Compiler Reference](#)

End of IBM extension
Chapter 8. Functions

In the context of programming languages, the term function means an assemblage of statements used for computing an output value. The word is used less strictly than in mathematics, where it means a set relating input variables uniquely to output variables. Functions in C or C++ programs may not produce consistent outputs for all inputs, may not produce output at all, or may have side effects. Functions can be understood as user-defined operations, in which the parameters of the parameter list, if any, are the operands.

This section discusses the following topics:
- “Function declarations and definitions”
- “Function storage class specifiers” on page 196
- “Function specifiers” on page 196
- “Function return type specifiers” on page 201
- “Function declarators” on page 202
- “Function attributes” on page 206
- “The main() function” on page 212
- “Function calls” on page 213
- “Default arguments in C++ functions” on page 216
- “Pointers to functions” on page 219
- “Nested functions” on page 219

Function declarations and definitions

The distinction between a function declaration and function definition is similar to that of a data declaration and definition. The declaration establishes the names and characteristics of a function but does not allocate storage for it, while the definition allocates the body for a function, associates an identifier with the function, and allocates storage for it. Thus, the identifiers declared in this example:

```c
float square(float x);
```

do not allocate storage.

The function definition contains a function declaration and the body of a function. The body is a block of statements that perform the work of the function. The identifiers declared in this example allocate storage; they are both declarations and definitions.

```c
float square(float x)
{ return x*x; }
```

A function can be declared several times in a program, but all declarations for a given function must be compatible; that is, the return type is the same and the parameters have the same type. However, a function can only have one definition. Declarations are typically placed in header files, while definitions appear in source files.
Function declarations

A function identifier preceded by its return type and followed by its parameter list is called a function declaration or function prototype. The prototype informs the compiler of the format and existence of a function prior to its use. The compiler checks for mismatches between the parameters of a function call and those in the function declaration. The compiler also uses the declaration for argument type checking and argument conversions.

- C++ Implicit declaration of functions is not allowed: you must explicitly declare every function you can call it.

- C If a function declaration is not visible at the point at which a call to the function is made, the compiler assumes an implicit declaration of extern int func(); However, for conformance to C99, you should explicitly prototype every function before making a call to it.

The elements of a declaration for a function are as follows:

- **Function storage class specifiers** which specify linkage
- **Function return type specifiers** which specify the data type of a value to be returned
- **Function specifiers** which specify additional properties for functions
- **Function declarators** which include function identifiers as well as lists of parameters

All function declarations have the form:

**Function declaration syntax**

```
storage_classSpecifier functionSpecifier return_typeSpecifier
```

```
function_declarator ;
```

- IBM In addition, for compatibility with GNU C and C++, XL C/C++ allows you to use attributes to modify the properties of functions. They are described in “Function attributes” on page 206.

Function definitions

The elements of a function definition are as follows:

- **Function storage class specifiers** which specify linkage
- **Function return type specifiers** which specify the data type of a value to be returned
- **Function specifiers** which specify additional properties for functions
- **Function declarators** which include function identifiers as well as lists of parameters

- The **function body**, which is a braces-enclosed series of statements representing the actions that the function performs

- C++ Constructor-initializers, which are used only in constructor functions declared in classes; they are described in “Constructors” on page 305.
**C++** Try blocks, which are used in class functions; they are described in "try blocks" on page 363.

In addition, for compatibility with GNU C and C++, XL C/C++ allows you to use attributes to modify the properties of functions. They are described in "Function attributes" on page 206.

Function definitions take the following form:

```
C only

Function definition syntax

```

```
C++ only

Function definition syntax

```

Examples of function declarations
The following code fragments show several function declarations (or prototypes). The first declares a function f that takes two integer arguments and has a return type of `void`:

```c
void f(int, int);
```

This fragment declares a pointer p1 to a function that takes a pointer to a constant character and returns an integer:

```c
int (*p1) (const char*);
```

The following code fragment declares a function f1 that takes an integer argument, and returns a pointer to a function that takes an integer argument and returns an integer:

```c
int (*f1(int)) (int);
```

Alternatively, a typedef can be used for the complicated return type of function f1:

```c
typedef int f1_return_type(int);
int (*f1) (int);
```

```c
```
The following declaration is of an external function `f2` that takes a constant integer as its first argument, can have a variable number and variable types of other arguments, and returns type `int`.

```c
int extern f2(const int, ...); /* C version */
int extern f2(const int ...); // C++ version
```

Function `f6` is a `const` class member function of class `X`, takes no arguments, and has a return type of `int`:

```c
class X
{
public:
    int f6() const;
};
```

Function `f4` takes no arguments, has return type `void`, and can throw class objects of types `X` and `Y`.

```c
class X;
class Y;
// ...
void f4() throw(X,Y);
```

**Examples of function definitions**

The following example is a definition of the function `sum`:

```c
int sum(int x, int y)
{
    return(x + y);
}
```

The function `sum` has external linkage, returns an object that has type `int`, and has two parameters of type `int` declared as `x` and `y`. The function body contains a single statement that returns the sum of `x` and `y`.

The following function `set_date` declares a pointer to a structure of type `date` as a parameter. `date_ptr` has the storage class specifier `register`.

```c
void set_date(register struct date *date_ptr)
{
    date_ptr->mon = 12;
    date_ptr->day = 25;
    date_ptr->year = 87;
}
```

**Compatible functions**

<table>
<thead>
<tr>
<th>C only</th>
</tr>
</thead>
</table>

For two functions, the composite type must meet the following requirements:

- If one of the function types has a parameter type list, the composite type is a function prototype with the same parameter type list.
- If both types are function types with parameter lists, then each parameter in the parameter list of the composite is the composite type of the corresponding parameters.

If the function declarator is not part of the function declaration, the parameters may have incomplete type. The parameters may also specify variable length array
types by using the [*] notation in their sequences of declarator specifiers. The following are examples of compatible function prototype declarators:

double maximum(int n, int m, double a[n][m]);
double maximum(int n, int m, double a[*][*]);
double maximum(int n, int m, double a[][*]);
double maximum(int n, int m, double a[][m]);

Related information
- "Compatible types" on page 41

Multiple function declarations

C++ only

All function declarations for a particular function must have the same number and type of parameters, and must have the same return type.

These return and parameter types are part of the function type, although the default arguments and exception specifications are not.

If a previous declaration of an object or function is visible in an enclosing scope, the identifier has the same linkage as the first declaration. However, a variable or function that has no linkage and later declared with a linkage specifier will have the linkage you have specified.

For the purposes of argument matching, ellipsis and linkage keywords are considered a part of the function type. They must be used consistently in all declarations of a function. If the only difference between the parameter types in two declarations is in the use of typedef names or unspecified argument array bounds, the declarations are the same. A const or volatile type qualifier is also part of the function type, but can only be part of a declaration or definition of a nonstatic member function.

If two function declarations match in both return type and parameter lists, then the second declaration is treated as redeclaration of the first. The following example declares the same function:

```c++
int foo(const string &bar);
int foo(const string &);
```

Declaring two functions differing only in return type is not valid function overloading, and is flagged as a compile-time error. For example:

```c++
void f();
int f(); // error, two definitions differ only in
         // return type
int g()
{
    return f();
}
```

Related information
- "Overloading functions" on page 229
Function storage class specifiers

For a function, the storage class specifier determines the linkage of the function. By default, function definitions have external linkage, and can be called by functions defined in other files. An exception is inline functions, which are treated by default as having internal linkage; see “Linkage of inline functions” on page 197 for more information.

A storage class specifier may be used in both function declarations and definitions. The only storage class options for functions are:

- `static`
- `extern`

The static storage class specifier

A function declared with the static storage class specifier has internal linkage, which means that it may be called only within the translation unit in which it is defined.

The static storage class specifier can be used in a function declaration only if it is at file scope. You cannot declare functions within a block as static.

- C++ This use of static is deprecated in C++. Instead, place the function in the unnamed namespace.

Related information

- “Internal linkage” on page 8
- Chapter 9, “Namespaces (C++ only),” on page 221

The extern storage class specifier

A function that is declared with the extern storage class specifier has external linkage, which means that it can be called from other translation units. The keyword extern is optional; if you do not specify a storage class specifier, the function is assumed to have external linkage.

Related information

- “External linkage” on page 9
- “Language linkage (C++ only)” on page 10
- “Class scope (C++ only)” on page 5
- Chapter 9, “Namespaces (C++ only),” on page 221

Function specifiers

The available function specifiers for function definitions are:

- `inline` which instructs the compiler to expand a function definition at the point of a function call.
- C++ `explicit`, which can only be used for member functions of classes, and is described in “The explicit specifier” on page 323
- C++ `virtual`, which can only be used for member functions of classes, and is described in “Virtual functions” on page 295
The inline function specifier

An inline function is one for which the compiler copies the code from the function definition directly into the code of the calling function rather than creating a separate set of instructions in memory. Instead of transferring control to and from the function code segment, a modified copy of the function body may be substituted directly for the function call. In this way, the performance overhead of a function call is avoided.

> C  Any function, with the exception of main, can be declared or defined as inline with the inline function specifier. Static local variables are not allowed to be defined within the body of an inline function.

> C++  C++ functions implemented inside of a class declaration are automatically defined inline. Regular C++ functions and member functions declared outside of a class declaration, with the exception of main, can be declared or defined as inline with the inline function specifier. Static locals and string literals defined within the body of an inline function are treated as the same object across translation units; see “Linkage of inline functions” for details.

The following code fragment shows an inline function definition:
```c
inline int add(int i, int j) { return i + j; }
```

The use of the inline specifier does not change the meaning of the function. However, the inline expansion of a function may not preserve the order of evaluation of the actual arguments.

The most efficient way to code an inline function is to place the inline function definition in a header file, and then include the header in any file containing a call to the function which you would like to inline.

Note: The inline qualifier is represented by the following keywords:

- **C**  The inline keyword is only recognized under compilation with c99 or with the -qlanglvl=stdc99 or -qlanglvl=extc99 options (or equivalent pragmas) or -qkeyword=inline. Note that the latter option is enabled by default for xlc in the configuration file that is shipped with the compiler. The inline keyword is recognized at all language levels; however, see “Linkage of inline functions” below for the semantics of this keyword.

- **C++**  The inline and __inline__ keywords are recognized at all language levels.

Related information

- “The always_inline function attribute” on page 208
- “The noinline function attribute” on page 210
- -qlanglvl and -qkeyword in the XL C/C++ Compiler Reference

Linkage of inline functions

In C, inline functions are treated by default as having static linkage; that is, they are only visible within a single translation unit. Therefore, in the following example, even though function foo is defined in exactly the same way, foo in file
a.c and foo in file b.c are treated as separate functions: two function bodies are generated, and assigned two different addresses in memory:

```c
// a.c
#include <stdio.h>
inline int foo()
     return 3;
}

void g()
     printf("foo called from g: return value = %d, address = %p\n", foo(), &foo);
}

// b.c
#include <stdio.h>
inline int foo()
     return 3;
}

void g();

int main()
     printf("foo called from main: return value = %d, address = %p\n", foo(), &foo);
g();
```

The output from the compiled program is:

```
foo called from main: return value = 3, address = 0x10000580
foo called from g: return value = 3, address = 0x10000500
```

Since inline functions are treated as having internal linkage, an inline function definition can co-exist with a regular, external definition of a function with the same name in another translation unit. However, when you call the function from the file containing the inline definition, the compiler may choose either the inline version defined in the same file or the external version defined in another file for the call; your program should not rely on the inline version being called. In the following example, the call to foo from function g could return either 6 or 3:

```c
// a.c
#include <stdio.h>
inline int foo()
     return 6;
}

void g()
     printf("foo called from g: return value = %d\n", foo());
}

// b.c
#include <stdio.h>
int foo()
     return 3;
}

void g();
```
int main() {  
printf("foo called from main: return value = %d\n", foo());  
g();  
}

Similarly, if you define a function as extern inline, or redefine an inline function as extern, the function simply becomes a regular, external function and is not inlined.

If you specify the \_\_inline\_ keyword, with the trailing underscores, the compiler uses the GNU C semantics for inline functions. In contrast to the C99 semantics, a function defined as \_\_inline\_ provides an external definition only; a function defined as static \_\_inline\_ provides an inline definition with internal linkage (as in C99); and a function defined as extern \_\_inline\_, when compiled with optimization enabled, allows the co-existence of an inline and external definition of the same function. For more information on the GNU C implementation of inline functions, see the GCC documentation, available at http://gcc.gnu.org/onlinedocs/.

--- End of C only

---

**C++ only**

You must define an inline function in exactly the same way in each translation unit in which the function is used or called. Furthermore, if a function is defined as inline, but never used or called within the same translation unit, it is discarded by the compiler (unless you compile with the \texttt{-qkeepinlines} option).

Nevertheless, in C++, inline functions are treated by default as having external linkage, meaning that the program behaves as if there is only one copy of the function. The function will have the same address in all translation units and each translation unit will share any static locals and string literals. Therefore, compiling the previous example gives the following output:

```
foo called from main: return value = 3, address = 0x10000580
foo called from g: return value = 3, address = 0x10000580
```

Redefining an inline function with the same name but with a different function body is illegal; however, the compiler does not flag this as an error, but simply generates a function body for the version defined in the first file entered on the compilation command line, and discards the others. Therefore, the following example, in which inline function \texttt{foo} is defined differently in two different files, may not produce the expected results:

```
// a.C
#include <stdio.h>
inline int foo(){
  return 6;
}

void g() {
  printf("foo called from g: return value = %d, address = %p\n", foo(), &foo);
}

// b.C
#include <stdio.h>
```
inline int foo(){
    return 3;
}

void g();

int main()
{
    printf("foo called from main: return value = %d, address = %p\n", foo(), &foo);
    g();
}

When compiled with the command xlC a.C b.C, the output is:
foo called from main: return value = 6, address = 0x10001640
foo called from g: return value = 6, address = 0x10001640

The call to foo from main does not use the inline definition provided in b.C, but rather calls foo as a regular external function defined in a.C. It is your responsibility to ensure that inline function definitions with the same name match exactly across translation units, to avoid unexpected results.

Because inline functions are treated as having external linkage, any static local variables or string literals that are defined within the body of an inline function are treated as the same object across translation units. The following example demonstrates this:
// a.C
#include <stdio.h>

inline int foo(){
    static int x = 23;
    printf("address of x = %p\n", &x);
    x++;    
    return x;
}

void g(){
    printf("foo called from g: return value = %d\n", foo());
}

// b.C
#include <stdio.h>

inline int foo()
{
    static int x=23;
    printf("address of x = %p\n", &x);
    x++;    
    return x;
}

void g();

int main()
{
    printf("foo called from main: return value = %d\n", foo());
    g();
}

The output of this program shows that x in both definitions of foo is indeed the same object:
address of x = 0x10011d5c
foo called from main: return value = 24
address of x = 0x10011d5c
foo called from g: return value = 25

If you want to ensure that each instance of function defined as inline is treated as a separate function, you can use the static specifier in the function definition in each translation unit, or compile with the -qstaticinline option. Note, however, that static inline functions are removed from name lookup during template instantiation, and are not found.

Related information
• “The static storage class specifier” on page 196
• “The extern storage class specifier” on page 196
• -qstaticinline and -qkeepinlines in the XL C/C++ Compiler Reference

Function return type specifiers
The result of a function is called its return value and the data type of the return value is called the return type.

Every function declaration and definition must specify a return type, whether or not it actually returns a value.

If a function declaration does not specify a return type, the compiler assumes an implicit return type of int. However, for conformance to C99, you should specify a return type for every function declaration and definition, whether or not the function returns int.

A function may be defined to return any type of value, except an array type or a function type; these exclusions must be handled by returning a pointer to the array or function. When a function does not return a value, void is the type specifier in the function declaration and definition.

A function cannot be declared as returning a data object having a volatile or const type, but it can return a pointer to a volatile or const object.

A function can have a return type that is a user-defined type. For example:
enum count {one, two, three};
enum count counter();

The user-defined type may also be defined within the function declaration.

enum count {one, two, three} counter(); // legal in C
enum count {one, two, three} counter(); // error in C++

References can also be used as return types for functions. The reference returns the lvalue of the object to which it refers.

Related information
• “Type specifiers” on page 49
Function return values

C If a function is defined as having a return type of void, it should not return a value. C++ In C++, a function which is defined as having a return type of void, or is a constructor or destructor, must not return a value.

C If a function is defined as having a return type other than void, it should return a value. Under compilation for strict C99 conformance, a function defined with a return type must include an expression containing the value to be returned.

C++ A function defined with a return type must include an expression containing the value to be returned.

When a function returns a value, the value is returned via a return statement to the caller of the function, after being implicitly converted to the return type of the function in which it is defined. The following code fragment shows a function definition including the return statement:

```c
int add(int i, int j)
{
    return i + j; // return statement
}
```

The function add() can be called as shown in the following code fragment:

```c
int a = 10,
    b = 20;
int answer = add(a, b); // answer is 30
```

In this example, the return statement initializes a variable of the returned type. The variable answer is initialized with the int value 30. The type of the returned expression is checked against the returned type. All standard and user-defined conversions are performed as necessary.

Each time a function is called, new copies of its variables with automatic storage are created. Because the storage for these automatic variables may be reused after the function has terminated, a pointer or reference to an automatic variable should not be returned. C++ If a class object is returned, a temporary object may be created if the class has copy constructors or a destructor.

Related information
- “The return statement” on page 183
- “Temporary objects” on page 320
- “Overloading assignments” on page 237
- “Overloading subscripting” on page 239
- “The auto storage class specifier” on page 44

Function declarators

Function declarators consist of the following elements:

- An identifier, or name
- Parameter declarations, which specify the parameters that can be passed to the function in a function call
Exception declarations, which include throw expressions; exception specifications are described in Chapter 16, “Exception handling (C++ only),” on page 363.

The type qualifiers const and volatile, which are used only in class member functions; they are described in “Constant and volatile member functions” on page 258.

**C only**

Function declarator syntax

```
identifier (parameter_declaration)
```

End of C only

**C++ only**

Function declarator syntax

```
identifier (parameter_declaration) cv-qualifier
exception declaration
```

End of C++ only

Related information

- “Default arguments in C++ functions” on page 216

**Parameter declarations**

The function declarator includes the list of parameters that can be passed to the function when it is called by another function, or by itself.

In C++, the parameter list of a function is referred to as its signature. The name and signature of a function uniquely identify it. As the word itself suggests, the function signature is used by the compiler to distinguish among the different instances of overloaded functions.

**Function parameter declaration syntax**

```
(parameter ,...,)
```

```
register type_specifier declarator
```

**Chapter 8. Functions** 203
An empty argument list in a function declaration or definition indicates that a function that takes no arguments. To explicitly indicate that a function does not take any arguments, you can declare the function in two ways: with an empty parameter list, or with the keyword void:

```c
int f(void);
int f();
```

An empty argument list in a function declaration indicates that a function may take any number or type of arguments. Thus, `int f()` indicates that function f takes no arguments. However, `int f();` simply indicates that the number and type of parameters is not known. To explicitly indicate that a function does not take any arguments, you should define the function with the keyword void.

An ellipsis at the end of the parameter specifications is used to specify that a function has a variable number of parameters. The number of parameters is equal to, or greater than, the number of parameter specifications.

```c
int f(int, ...);
```

- The comma before the ellipsis is optional. In addition, a parameter declaration is not required before the ellipsis.
- At least one parameter declaration, as well as a comma before the ellipsis, are both required in C.

**Related information**
- “The void type” on page 53
- “Type specifiers” on page 49
- “Type qualifiers” on page 68
- “Exception specifications” on page 374

**Parameter types**

In a function declaration, or prototype, the type of each parameter must be specified. In the function definition, the type of each parameter must also be specified. In the function definition, if the type of a parameter is not specified, it is assumed to be `int`. 
A variable of a user-defined type may be declared in a parameter declaration, as in
the following example, in which x is declared for the first time:

```c
struct X { int i; }
void print(struct X x);
```

> C  The user-defined type can also be defined within the parameter
    declaration. > C++ The user-defined type can not be defined within the
    parameter declaration.

```c
void print(struct X { int i; } x); // legal in C
void print(struct X { int i; } x); // error in C++
```

**Parameter names**

In a function *definition*, each parameter must have an identifier. In a function
*declaration*, or prototype, specifying an identifier is optional. Thus, the following
example is legal in a function declaration:

```c
int func(int,long);
```

---

**C++ only**

The following constraints apply to the use of parameter names in function declarations:

- Two parameters cannot have the same name within a single declaration.
- If a parameter name is the same as a name outside the function, the name
  outside the function is hidden and cannot be used in the parameter declaration.

In the following example, the third parameter name `intersects` is meant to have
enumeration type `subway_line`, but this name is hidden by the name of the first
parameter. The declaration of the function `subway()` causes a compile-time error
because `subway_line` is not a valid type name because the first parameter name `subway_line`
hides the namespace scope `enum` type and cannot be used again in
the second parameter.

```c
enum subway_line {yonge, university, spadina, bloor};
int subway(char * subway_line, int stations,
           subway_line intersects);
```

---

**Static array indices in function parameter declarations (C only)**

Except in certain contexts, an unsubscripted array name (for example, `region`
instead of `region[4]`) represents a pointer whose value is the address of the first
element of the array, provided that the array has previously been declared. An
array type in the parameter list of a function is also converted to the corresponding
pointer type. Information about the size of the argument array is lost when the
array is accessed from within the function body.

To preserve this information, which is useful for optimization, you may declare the
index of the argument array using the `static` keyword. The constant expression
specifies the minimum pointer size that can be used as an assumption for
optimizations. This particular usage of the `static` keyword is highly prescribed.
The keyword may only appear in the outermost array type derivation and only in
function parameter declarations. If the caller of the function does not abide by
these restrictions, the behavior is undefined.

The following examples show how the feature can be used.
void foo(int arr[static 10]); /* arr points to the first of at least 10 ints */
void foo(int arr[const 10]); /* arr is a const pointer */
void foo(int arr[static const i]); /* arr points to at least i ints; i is computed at run time. */
void foo(int arr[const static i]); /* alternate syntax to previous example */
void foo(int arr[const]); /* const pointer to int */

Related information
- "The static storage class specifier" on page 44
- "Arrays" on page 85
- "Array subscipting operator [ ]" on page 122

## Function attributes

**IBM extension**

Function attributes are extensions implemented to enhance the portability of programs developed with GNU C. Specifiable attributes for functions provide explicit ways to help the compiler optimize function calls and to instruct it to check more aspects of the code. Others provide additional functionality.

IBM C and C++ implement a subset of the GNU C function attributes. If a particular function attribute is not implemented, its specification is accepted and the semantics are ignored. These language features are collectively available when compiling in any of the extended language levels.

A function attribute is specified with the keyword `__attribute__` followed by the attribute name and any additional arguments the attribute name requires. A function `__attribute__` specification is included in the declaration or definition of a function. The syntax takes the following forms:

**Function attribute syntax: function declaration**

```
>>>function declarator—__attribute__—
     ((—attribute_name—))—;
```

**C only**

**Function attribute syntax: function definition**

```
>>>__attribute__—((—attribute_name—))—function declarator
```
Function attribute syntax: function definition

```c
function declarator__attribute__

```__attribute_name__
```

End of C++ only

The function attribute in a function declaration is always placed after the declarator, including the parenthesized parameter declaration:

```c
/* Specify the attribute on a function prototype declaration */
void f(int i, int j) __attribute__((individual_attribute_name));
void f(int i, int j) {}
```

End of C++ only

In C++, the attribute specification must also follow any exception declaration that may be present for the function.

End of C++ only

C only

Due to ambiguities in parsing old-style parameter declarations, a function definition must have the attribute specification precede the declarator:

```c
int __attribute__((individual_attribute_name)) foo(int i) {}
```

End of C only

A function attribute specification using the form __attribute_name__ (that is, the attribute name with double underscore characters leading and trailing) reduces the likelihood of a name conflict with a macro of the same name.

Related information
- “Variable attributes” on page 101

The alias function attribute

The alias function attribute causes the function declaration to appear in the object file as an alias for another symbol. This language feature provides a technique for coping with duplicate or cumbersome names.
**alias function attribute syntax**

```
_attribute__((alias("original_function_name")))(__attribute__((alias)))
```

**C**  The aliased function can be defined after the specification of its alias with this function attribute. C also allows an alias specification in the absence of a definition of the aliased function in the same compilation unit.

**C++** The `original_function_name` must be the mangled name.

The following declares `bar` to be an alias for `__foo`:

```
void __foo(){ /* function body */ }
void bar() __attribute__((alias("__foo")));
```

The compiler does not check for consistency between the declaration of `bar` and definition of `__foo`. Such consistency remains the responsibility of the programmer.

**The always_inline function attribute**

The `always_inline` function attribute instructs the compiler to inline an `inline` function, regardless of whether optimization was specified at compile time. However, the attribute has no effect if the program is compiled at no-opt levels. Specifying this attribute for a function without an `inline` specification also has no effect. The attribute takes precedence over inlining compiler options.

**always_inline function attribute syntax**

```
_attribute__((always_inline))
```

**Related information**

- [“The inline function specifier” on page 197](#)
- [“The noinline function attribute” on page 210](#)

**The const function attribute**

The `const` function attribute allows you to tell the compiler that the function can safely be called fewer times than indicated in the source code. The language feature provides you with an explicit way to help the compiler optimize code by indicating that the function does not examine any values except its arguments and has no effects except for its return value.
The following kinds of functions should not be declared const:
• A function with pointer arguments which examines the data pointed to.
• A function that calls a non-const function.

The constructor and destructor function attributes

The constructor and destructor function attributes provide the ability to write a function that initializes data or releases storage that is used implicitly during program execution. A function to which the constructor function attribute has been applied is called automatically before execution enters main. Similarly, a function to which the destructor attribute has been applied is called automatically after calling exit or upon completion of main.

When the constructor or destructor function is called automatically, the return value of the function is ignored, and any parameters of the function are undefined.

A function declaration containing a constructor or destructor function attribute must match all of its other declarations.

The format function attribute

The format function attribute provides a way to identify user-defined functions that take format strings as arguments so that calls to these functions will be type-checked against a format string, similar to the way the compiler checks calls to the functions printf, scanf, strftime, and strftime for errors.

where

string_index
Is a constant integral expression that specifies which argument in the declaration of the user function is the format string argument.
C++, the minimum value of string_index for nonstatic member functions is 2 because the first argument is an implicit this argument. This behavior is consistent with that of GNU C++.

first_to_check
Is a constant integral expression that specifies the first argument to check against the format string. If there are no arguments to check against the format string (that is, diagnostics should only be performed on the format string syntax and semantics), first_to_check should have a value of 0. For strftime-style formats, first_to_check is required to be 0.

It is also possible to diagnose the same string for different format styles. All styles are diagnosed.

The format_arg function attribute
The format_arg function attribute provides a way to identify user-defined functions that modify format strings. Once the function is identified, calls to functions like printf, scanf, strftime, or strftime_mon, whose operands are a call to the user-defined function can be checked for errors.

format_arg function attribute syntax

|||attribute___(__(format_arg(__(string_index---))----string_index__))||

where string_index is a constant integral expression that specifies which argument is the format string argument, starting from 1. For non-static member functions in C++, string_index starts from 2 because the first parameter is an implicit this parameter.

It is possible to specify multiple format_arg attributes on the same function, in which case, all apply.

extern char* my_dgettext(const char* my_format, const char* my_format2)
 ATTRIBUTE__(format_arg__1)) ATTRIBUTE__(format_arg__2));

printf(my_dgettext("%","%"));
//printf-style format diagnostics are performed on both "%" strings

The noinline function attribute
The noinline function attribute prevents the function to which it is applied from being inlined, regardless of whether the function is declared inline or non-inline. The attribute takes precedence over inlining compiler options, the inline keyword, and the always_inline function attribute.

noinline function attribute syntax
Other than preventing inlining, the attribute does not remove the semantics of inline functions.

**The noreturn function attribute**

The `noreturn` function attribute allows you to indicate to the compiler that the function is not intended to return. The language feature provides the programmer with another explicit way to help the compiler optimize code and to reduce false warnings for uninitialized variables.

The return type of the function should be `void`.

**noreturn function attribute syntax**

```c
_attribute__((noreturn))
```

Registers saved by the calling function may not necessarily be restored before calling the nonreturning function.

**Related information**

- `#pragma leaves` in the [XL C/C++ Compiler Reference](#)

**The pure function attribute**

The `pure` function attribute allows you to declare a function that can be called fewer times than what is literally in the source code. Declaring a function with the attribute `pure` indicates that the function has no effect except a return value that depends only on the parameters, global variables, or both.

**pure function attribute syntax**

```c
_attribute__((pure))
```

**Related information**

- `#pragma isolated_call` in the [XL C/C++ Compiler Reference](#)

**The section function attribute**

The `section` function attribute specifies the section in the object file in which the compiler should place its generated code. The language feature provides the ability to control the section in which a function should appear.

**section function attribute syntax**

```c
_attribute__((section("section_name")))
```

where `section_name` is a string literal.
Each defined function can reside in only one section. The section indicated in a function definition should match that in any previous declaration. The section indicated in a function definition cannot be overwritten, whereas one in a function declaration can be overwritten by a later specification. Moreover, if a section attribute is applied to a function declaration, the function will be placed in the specified section only if it is defined in the same compilation unit.

Related information

- “The section variable attribute” on page 104

The weak function attribute

The weak function attribute causes the symbol resulting from the function declaration to appear in the object file as a weak symbol, rather than a global one. The language feature provides the programmer writing library functions with a way to allow function definitions in user code to override the library function declaration without causing duplicate name errors.

weak function attribute syntax

```c
__attribute__((weak))
```

Related information

- “The alias function attribute” on page 207
- `#pragma weak` in the XL C/C++ Compiler Reference
- “The weak variable attribute” on page 105

End of IBM extension

The main() function

When a program begins running, the system calls the function main, which marks the entry point of the program. By default, main has the storage class extern. Every program must have one function named main, and the following constraints apply:

- No other function in the program can be called main.
- main cannot be defined as inline or static.
- `C++` main cannot be called from within a program.
- `C++` The address of main cannot be taken.
- `C++` The main function cannot be overloaded.

The function main can be defined with or without parameters, using any of the following forms:

```c
int main (void)
int main ( )
int main(int argc, char *argv[])
int main (int argc, char ** argv)
```

Although any name can be given to these parameters, they are usually referred to as argc and argv. The first parameter, argc (argument count) is an integer that indicates how many arguments were entered on the command line when the program was started. The second parameter, argv (argument vector), is an array of
pointers to arrays of character objects. The array objects are null-terminated strings, representing the arguments that were entered on the command line when the program was started.

The first element of the array, argv[0], is a pointer to the character array that contains the program name or invocation name of the program that is being run from the command line. argv[1] indicates the first argument passed to the program, argv[2] the second argument, and so on. The following example program backward prints the arguments entered on a command line such that the last argument is printed first:

```c
#include <stdio.h>
int main(int argc, char *argv[]) {
    while (--argc > 0)
        printf("%s ", argv[argc]);
}
```

Invoking this program from a command line with the following:

```
backward string1 string2
```

gives the following output:

```
string2 string1
```

The arguments argc and argv would contain the following values:

<table>
<thead>
<tr>
<th>Object</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>argc</td>
<td>3</td>
</tr>
<tr>
<td>argv[0]</td>
<td>pointer to string &quot;backward&quot;</td>
</tr>
<tr>
<td>argv[1]</td>
<td>pointer to string &quot;string1&quot;</td>
</tr>
<tr>
<td>argv[2]</td>
<td>pointer to string &quot;string2&quot;</td>
</tr>
<tr>
<td>argv[3]</td>
<td>NULL</td>
</tr>
</tbody>
</table>

Related information

- “The extern storage class specifier” on page 46
- “The inline function specifier” on page 197
- “The static storage class specifier” on page 44
- “Function calls”

Function calls

Once a function has been declared and defined, it can be called from anywhere within the program: from within the main function, from another function, and even from itself. Calling the function involves specifying the function name, followed by the function call operator and any data values the function expects to receive. These values are the arguments for the parameters defined for the function, and the process just described is called passing arguments to the function.

> C++

A function may not be called if it has not already been declared.

Passing arguments can be done in two ways:

- **Pass by value**, which copies the value of an argument to the corresponding parameter in the called function
Pass by reference, which passes the address of an argument to the corresponding parameter in the called function

---

C++ only

If a class has a destructor or a copy constructor that does more than a bitwise copy, passing a class object by value results in the construction of a temporary object that is actually passed by reference.

It is an error when a function argument is a class object and all of the following properties hold:

- The class needs a copy constructor.
- The class does not have a user-defined copy constructor.
- A copy constructor cannot be generated for that class.

---

Related information

- "Function argument conversions" on page 112
- "Function call operator ()" on page 121
- "Constructors" on page 305

Pass by value

When you use pass-by-value, the compiler copies the value of an argument in a calling function to a corresponding non-pointer or non-reference parameter in the called function definition. The parameter in the called function is initialized with the value of the passed argument. As long as the parameter has not been declared as constant, the value of the parameter can be changed, but the changes are only performed within the scope of the called function only; they have no effect on the value of the argument in the calling function.

In the following example, main passes func two values: 5 and 7. The function func receives copies of these values and accesses them by the identifiers a and b. The function func changes the value of a. When control passes back to main, the actual values of x and y are not changed.

```c
/**
 ** This example illustrates calling a function by value
 **/

#include <stdio.h>

void func (int a, int b)
{
    a += b;
    printf("In func, a = %d b = %d\n", a, b);
}

int main(void)
{
    int x = 5, y = 7;
    func(x, y);
    printf("In main, x = %d y = %d\n", x, y);
    return 0;
}
```

The output of the program is:

21

XL C/C++ Language Reference
In func, a = 12 b = 7
In main, x = 5 y = 7

Pass by reference
Passing by \textit{by reference} refers to a method of passing the address of an argument in the calling function to a corresponding parameter in the called function.

\begin{itemize}
\item \textbf{C} In C, the corresponding parameter in the called function must be declared as a pointer type.
\item \textbf{C++} In C++, the corresponding parameter can be declared as any reference type, not just a pointer type.
\end{itemize}

In this way, the value of the argument in the calling function can be modified by the called function.

The following example shows how arguments are passed by reference. In C++, the reference parameters are initialized with the actual arguments when the function is called. In C, the pointer parameters are initialized with pointer values when the function is called.

\begin{verbatim}
C++ only
#include <stdio.h>
void swapnum(int &i, int &j) {
  int temp = i;
  i = j;
  j = temp;
}
int main(void) {
  int a = 10;
  int b = 20;
  swapnum(a, b);
  printf("A is %d and B is %d\n", a, b);
  return 0;
}
\end{verbatim}

\begin{verbatim}
C only
#include <stdio.h>
void swapnum(int *i, int *j) {
  int temp = *i;
  *i = *j;
  *j = temp;
}
int main(void) {
  int a = 10;
  int b = 20;
  swapnum(&a, &b);
  printf("A is %d and B is %d\n", a, b);
  return 0;
}
\end{verbatim}

When the function \texttt{swapnum()} is called, the actual values of the variables \texttt{a} and \texttt{b} are exchanged because they are passed by reference. The output is:
A is 20 and B is 10

In order to modify a reference that is const-qualified, you must cast away its constness with the const_cast operator. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

void f(const int& x) {
    int* y = const_cast<int*>(&x);
    (*y)++;
}

int main() {
    int a = 5;
    f(a);
    cout << a << endl;
}
```

This example outputs 6.

---

Related information

- [“References (C++ only)” on page 88](#)
- [“The const_cast operator (C++ only)” on page 128](#)

---

Default arguments in C++ functions

You can provide default values for function parameters. For example:

```cpp
#include <iostream>
using namespace std;

int a = 1;
int f(int a) { return a; }
int g(int x = f(a)) { return x; }

int h() {
    a = 2;
    {
        int a = 3;
        return g();
    }
}

int main() {
    cout << h() << endl;
}
```

This example prints 2 to standard output, because the a referred to in the declaration of g() is the one at file scope, which has the value 2 when g() is called.

The default argument must be implicitly convertible to the parameter type.
A pointer to a function must have the same type as the function. Attempts to take
the address of a function by reference without specifying the type of the function
will produce an error. The type of a function is not affected by arguments with
default values.

The following example shows that default arguments are not considered part of a
function's type. The default argument allows you to call a function without
specifying all of the arguments, it does not allow you to create a pointer to the
function that does not specify the types of all the arguments. Function f can be
called without an explicit argument, but the pointer badpointer cannot be defined
without specifying the type of the argument:

```c
int f(int = 0);
void g()
{
    int a = f(1);  // ok
    int b = f();   // ok, default argument used
}
int (*pointer)(int) = &f;  // ok, type of f() specified (int)
int (*badpointer)() = &f;  // error, badpointer and f have
different types. badpointer must
// be initialized with a pointer to
// a function taking no arguments.
```

In this example, function f3 has a return type int, and takes an int argument with
a default value that is the value returned from function f2:

```c
const int j = 5;
int f3( int x = f2(j) );
```

**Related information**

- "Pointers to functions" on page 219

### Restrictions on default arguments

Of the operators, only the function call operator and the operator new can have
default arguments when they are overloaded.

Parameters with default arguments must be the trailing parameters in the function
declaration parameter list. For example:

```c
void f(int a, int b = 2, int c = 3);  // trailing defaults
void g(int a = 1, int b = 2, int c);  // error, leading defaults
void h(int a, int b = 3, int c);     // error, default in middle
```

Once a default argument has been given in a declaration or definition, you cannot
redefine that argument, even to the same value. However, you can add default
arguments not given in previous declarations. For example, the last declaration
below attempts to redefine the default values for a and b:

```c
void f(int a, int b, int c=1);     // valid
void f(int a, int b=1, int c);     // valid, add another default
void f(int a=1, int b, int c);     // valid, add another default
void f(int a=1, int b=1, int c=1); // error, redefined defaults
```

You can supply any default argument values in the function declaration or in the
definition. Any parameters in the parameter list following a default argument
value must have a default argument value specified in this or a previous
declaration of the function.

You cannot use local variables in default argument expressions. For example, the
compiler generates errors for both function g() and function h() below:
void f(int a)
{
    int b=4;
    void g(int c=a); // Local variable "a" cannot be used here
    void h(int d=b); // Local variable "b" cannot be used here
}

Related information
• "Function call operator ( )" on page 121
• "The new operator (C++ only)" on page 140

Evaluation of default arguments

When a function defined with default arguments is called with trailing arguments missing, the default expressions are evaluated. For example:

```c
void f(int a, int b = 2, int c = 3); // declaration
// ...
int a = 1;
f(a);       // same as call f(a,2,3)
f(a,10);    // same as call f(a,10,3)
f(a,10,20); // no default arguments
```

Default arguments are checked against the function declaration and evaluated when the function is called. The order of evaluation of default arguments is undefined. Default argument expressions cannot use other parameters of the function. For example:

```c
int f(int q = 3, int r = q); // error
```

The argument `r` cannot be initialized with the value of the argument `q` because the value of `q` may not be known when it is assigned to `r`. If the above function declaration is rewritten:

```c
int q=5;
int f(int q = 3, int r = q); // error
```

The value of `r` in the function declaration still produces an error because the variable `q` defined outside of the function is hidden by the argument `q` declared for the function. Similarly:

```c
typedef double D;
int f(int D, int z = D(5.3)); // error
```

Here the type `D` is interpreted within the function declaration as the name of an integer. The type `D` is hidden by the argument `D`. The cast `D(5.3)` is therefore not interpreted as a cast because `D` is the name of the argument not a type.

In the following example, the nonstatic member `a` cannot be used as an initializer because `a` does not exist until an object of class `X` is constructed. You can use the static member `b` as an initializer because `b` is created independently of any objects of class `X`. You can declare the member `b` after its use as a default argument because the default values are not analyzed until after the final bracket `}` of the class declaration.

```c
class X
{
    int a;
    f(int z = a); // error
    g(int z = b); // valid
    static int b;
};
```
Pointers to functions

A pointer to a function points to the address of the executable code of the function. You can use pointers to call functions and to pass functions as arguments to other functions. You cannot perform pointer arithmetic on pointers to functions.

The type of a pointer to a function is based on both the return type and parameter types of the function.

A declaration of a pointer to a function must have the pointer name in parentheses. The function call operator () has a higher precedence than the dereference operator *. Without them, the compiler interprets the statement as a function that returns a pointer to a specified return type. For example:

```c
int *f(int a); /* function f returning an int*/
int (*g)(int a); /* pointer g to a function returning an int*/
char (*h)(int, int) /* h is a function that takes two integer parameters and returns char */
```

In the first declaration, f is interpreted as a function that takes an int as argument, and returns a pointer to an int. In the second declaration, g is interpreted as a pointer to a function that takes an int argument and that returns an int.

Related information

- “Language linkage (C++ only)” on page 10
- “Pointers” on page 82
- “Pointer conversions” on page 110
- “The extern storage class specifier” on page 196

Nested functions

A nested function is a function defined inside the definition of another function. It can be defined wherever a variable declaration is permitted, which allows nested functions within nested functions. Within the containing function, the nested function can be declared prior to being defined by using the auto keyword. Otherwise, a nested function has internal linkage. The language feature is an extension to C89 and C99, implemented to facilitate porting programs developed with GNU C.

A nested function can access all identifiers of the containing function that precede its definition.

A nested function must not be called after the containing function exits.

A nested function cannot use a goto statement to jump to a label in the containing function, or to a local label declared with the __label__ keyword inherited from the containing function.
Related information

- “Locally declared labels” on page 168

End of IBM extension
Chapter 9. Namespaces (C++ only)

A namespace is an optionally named scope. You declare names inside a namespace as you would for a class or an enumeration. You can access names declared inside a namespace the same way you access a nested class name by using the scope resolution (::) operator. However namespaces do not have the additional features of classes or enumerations. The primary purpose of the namespace is to add an additional identifier (the name of the namespace) to a name.

Related information
- “Scope resolution operator :: (C++ only)” on page 120

Defining namespaces

In order to uniquely identify a namespace, use the namespace keyword.

Namespace syntax

```
namespace identifier { namespace_body }
```

The identifier in an original namespace definition is the name of the namespace. The identifier may not be previously defined in the declarative region in which the original namespace definition appears, except in the case of extending namespace. If an identifier is not used, the namespace is an unnamed namespace.

Related information
- “Unnamed namespaces” on page 223

Declaring namespaces

The identifier used for a namespace name should be unique. It should not be used previously as a global identifier.

```
namespace Raymond {
  // namespace body here...
}
```

In this example, Raymond is the identifier of the namespace. If you intend to access a namespace’s elements, the namespace’s identifier must be known in all translation units.

Related information
- “File/global scope” on page 3

Creating a namespace alias

An alternate name can be used in order to refer to a specific namespace identifier.

```
namespace INTERNATIONAL_BUSINESS_MACHINES {
  void f();
}
```

```
namespace IBM = INTERNATIONAL_BUSINESS_MACHINES;
```
In this example, the IBM identifier is an alias for INTERNATIONAL_BUSINESS_MACHINES. This is useful for referring to long namespace identifiers.

If a namespace name or alias is declared as the name of any other entity in the same declarative region, a compiler error will result. Also, if a namespace name defined at global scope is declared as the name of any other entity in any global scope of the program, a compiler error will result.

**Related information**

- "File/global scope" on page 3

### Creating an alias for a nested namespace

Namespace definitions hold declarations. Since a namespace definition is a declaration itself, namespace definitions can be nested.

An alias can also be applied to a nested namespace.

```c
namespace INTERNATIONAL_BUSINESS_MACHINES {
  int j;
  namespace NESTED_IBM_PRODUCT {
    void a() { j++; }
    int j;
    void b() { j++; }
  }
}

namespace NIBM = INTERNATIONAL_BUSINESS_MACHINES::NESTED_IBM_PRODUCT
```

In this example, the NIBM identifier is an alias for the namespace NESTED_IBM_PRODUCT. This namespace is nested within the INTERNATIONAL_BUSINESS_MACHINES namespace.

**Related information**

- "Creating a namespace alias" on page 221

### Extending namespaces

Namespaces are extensible. You can add subsequent declarations to a previously defined namespace. Extensions may appear in files separate from or attached to the original namespace definition. For example:

```c
namespace X { // namespace definition
  int a;
  int b;
}

namespace X { // namespace extension
  int c;
  int d;
}

namespace Y { // equivalent to namespace X
  int a;
  int b;
  int c;
  int d;
}```
In this example, namespace X is defined with a and b and later extended with c and d. namespace X now contains all four members. You may also declare all of the required members within one namespace. This method is represented by namespace Y. This namespace contains a, b, c, and d.

Namespaces and overloading

You can overload functions across namespaces. For example:

```c
// Original X.h:
f(int);

// Original Y.h:
f(char);

// Original program.c:
#include "X.h"
#include "Y.h"

void z()
{
    f('a'); // calls f(char) from Y.h
}
```

Namespaces can be introduced to the previous example without drastically changing the source code.

```c
// New X.h:
namespace X {
    f(int);
}

// New Y.h:
namespace Y {
    f(char);
}

// New program.c:
#include "X.h"
#include "Y.h"

using namespace X;
using namespace Y;

void z()
{
    f('a'); // calls f() from Y.h
}
```

In program.c, function void z() calls function f(), which is a member of namespace Y. If you place the using directives in the header files, the source code for program.c remains unchanged.

Related information

- Chapter 10, “Overloading (C++ only),” on page 229

Unnamed namespaces

A namespace with no identifier before an opening brace produces an unnamed namespace. Each translation unit may contain its own unique unnamed namespace. The following example demonstrates how unnamed namespaces are useful.
```c++
#include <iostream>
using namespace std;

namespace {
    const int i = 4;
    int variable;
}

int main()
{
    cout << i << endl;
    variable = 100;
    return 0;
}

In the previous example, the unnamed namespace permits access to i and variable without using a scope resolution operator.

The following example illustrates an improper use of unnamed namespaces.
```
```c++
#include <iostream>
using namespace std;

namespace {
    const int i = 4;
}

int i = 2;

int main()
{
    cout << i << endl; // error
    return 0;
}

Inside main, i causes an error because the compiler cannot distinguish between the global name and the unnamed namespace member with the same name. In order for the previous example to work, the namespace must be uniquely identified with an identifier and i must specify the namespace it is using.

You can extend an unnamed namespace within the same translation unit. For example:
```
#include <iostream>
using namespace std;

namespace {
    int variable;
    void funct (int);
}

namespace {
    void funct (int i) { cout << i << endl; }
}

int main()
{
    funct(variable);
    return 0;
}
```
both the prototype and definition for funct are members of the same unnamed namespace.

Note: Items defined in an unnamed namespace have internal linkage. Rather than using the keyword static to define items with internal linkage, define them in an unnamed namespace instead.

Related information

- “Program linkage” on page 8
- “Internal linkage” on page 8

Namespace member definitions

A namespace can define its own members within itself or externally using explicit qualification. The following is an example of a namespace defining a member internally:

```cpp
namespace A {
    void b() { /* definition */ }
}
```

Within namespace A member void b() is defined internally.

A namespace can also define its members externally using explicit qualification on the name being defined. The entity being defined must already be declared in the namespace and the definition must appear after the point of declaration in a namespace that encloses the declaration’s namespace.

The following is an example of a namespace defining a member externally:

```cpp
namespace A {
    namespace B {
        void f();
    }
    void B::f() { /* defined outside of B */ }
}
```

In this example, function f() is declared within namespace B and defined (outside B) in A.

Namespaces and friends

Every name first declared in a namespace is a member of that namespace. If a friend declaration in a non-local class first declares a class or function, the friend class or function is a member of the innermost enclosing namespace.

The following is an example of this structure:

```cpp
// f has not yet been defined
void z(int);
namespace A {
    class X {
        friend void f(X); // A::f is a friend
    };
    // A::f is not visible here
    X x;
    void f(X) { /* definition */ } // f() is defined and known to be a friend
}
using A::x;
```
void z()
{
    A::f(x); // OK
    A::X::f(x); // error: f is not a member of A::X
}

In this example, function f() can only be called through namespace A using the
call A::f(s);. Attempting to call function f() through class X using the
A::X::f(x); call results in a compiler error. Since the friend declaration first occurs
in a non-local class, the friend function is a member of the innermost enclosing
namespace and may only be accessed through that namespace.

Related information
• “Friends” on page 272

The using directive

A using directive provides access to all namespace qualifiers and the scope
operator. This is accomplished by applying the using keyword to a namespace
identifier.

Using directive syntax

```
using namespace name;
```

The name must be a previously defined namespace. The using directive may be
applied at the global and local scope but not the class scope. Local scope takes
precedence over global scope by hiding similar declarations.

If a scope contains a using directive that nominates a second namespace and that
second namespace contains another using directive, the using directive from the
second namespace will act as if it resides within the first scope.

```c
namespace A {
    int i;
}
namespace B {
    int i;
    using namespace A;
}
void f() {
    using namespace B;
    i = 7; // error
}
```

In this example, attempting to initialize i within function f() causes a compiler
error, because function f() cannot know which i to call; i from namespace A, or i
from namespace B.

Related information
• “The using declaration and class members” on page 284

The using declaration and namespaces

A using declaration provides access to a specific namespace member. This is
accomplished by applying the using keyword to a namespace name with its
corresponding namespace member.
Using declaration syntax

In this syntax diagram, the qualifier name follows the using declaration and the member follows the qualifier name. For the declaration to work, the member must be declared inside the given namespace. For example:

```cpp
namespace A {
  int i;
  int k;
  void f;
  void g;
}
using A::k
```

In this example, the using declaration is followed by A, the name of namespace A, which is then followed by the scope operator (::), and k. This format allows k to be accessed outside of namespace A through a using declaration. After issuing a using declaration, any extension made to that specific namespace will not be known at the point at which the using declaration occurs.

Overloaded versions of a given function must be included in the namespace prior to that given function’s declaration. A using declaration may appear at namespace, block and class scope.

Related information
- “The using declaration and class members” on page 284

Explicit access

To explicitly qualify a member of a namespace, use the namespace identifier with a :: scope resolution operator.

Explicit access qualification syntax

```
namespace_name::member
```

For example:

```cpp
namespace VENDITTI {
  void j();
};
VENDITTI::j();
```

In this example, the scope resolution operator provides access to the function j held within namespace VENDITTI. The scope resolution operator :: is used to access identifiers in both global and local namespaces. Any identifier in an application can be accessed with sufficient qualification. Explicit access cannot be applied to an unnamed namespace.

Related information
- “Scope resolution operator :: (C++ only)” on page 120
Chapter 10. Overloading (C++ only)

If you specify more than one definition for a function name or an operator in the same scope, you have overloading that function name or operator. Overloaded functions and operators are described in “Overloading functions” and “Overloading operators” on page 231, respectively.

An overloaded declaration is a declaration that had been declared with the same name as a previously declared declaration in the same scope, except that both declarations have different types.

If you call an overloaded function name or operator, the compiler determines the most appropriate definition to use by comparing the argument types you used to call the function or operator with the parameter types specified in the definitions. The process of selecting the most appropriate overloaded function or operator is called overload resolution, as described in “Overload resolution” on page 241.

Overloading functions

You overload a function name f by declaring more than one function with the name f in the same scope. The declarations of f must differ from each other by the types and/or the number of arguments in the argument list. When you call an overloaded function named f, the correct function is selected by comparing the argument list of the function call with the parameter list of each of the overloaded candidate functions with the name f. A candidate function is a function that can be called based on the context of the call of the overloaded function name.

Consider a function print, which displays an int. As shown in the following example, you can overload the function print to display other types, for example, double and char*. You can have three functions with the same name, each performing a similar operation on a different data type:

```cpp
#include <iostream>
using namespace std;

void print(int i) { 
    cout << "Here is int " << i << endl; 
}
void print(double f) { 
    cout << "Here is float " << f << endl; 
}
void print(char* c) { 
    cout << "Here is char* " << c << endl; 
}

int main() { 
    print(10); 
    print(10.10); 
    print("ten"); 
}
```

The following is the output of the above example:

```
Here is int 10 
Here is float 10.1 
Here is char* ten 
```
Function overloading based on vector parameter types is supported.

**Related information**
- “Restrictions on overloaded functions”
- “Derivation” on page 279

**Restrictions on overloaded functions**

You cannot overload the following function declarations if they appear in the same scope. Note that this list applies only to explicitly declared functions and those that have been introduced through using declarations:

- Function declarations that differ only by return type. For example, you cannot declare the following declarations:
  ```
  int f();
  float f();
  ```

- Member function declarations that have the same name and the same parameter types, but one of these declarations is a static member function declaration. For example, you cannot declare the following two member function declarations of `f()`:
  ```
  struct A {
    static int f();
    int f();
  };
  ```

- Member function template declarations that have the same name, the same parameter types, and the same template parameter lists, but one of these declarations is a static template member function declaration.

- Function declarations that have equivalent parameter declarations. These declarations are not allowed because they would be declaring the same function.

- Function declarations with parameters that differ only by the use of `typedef` names that represent the same type. Note that a `typedef` is a synonym for another type, not a separate type. For example, the following two declarations of `f()` are declarations of the same function:
  ```
  typedef int I;
  void f(float, int);
  void f(float, I);
  ```

- Function declarations with parameters that differ only because one is a pointer and the other is an array. For example, the following are declarations of the same function:
  ```
  f(char*);
  f(char[10]);
  ```

  The first array dimension is insignificant when differentiating parameters; all other array dimensions are significant. For example, the following are declarations of the same function:
  ```
  g(char[10][20]);
  g(char[5][20]);
  ```

  The following two declarations are *not* equivalent:
  ```
  g(char[*][20]);
  g(char[*][40]);
  ```

- Function declarations with parameters that differ only because one is a function type and the other is a pointer to a function of the same type. For example, the following are declarations of the same function:
  ```
  void f(int(float));
  void f(int (*)(float));
  ```
• Function declarations with parameters that differ only because of cv-qualifiers 
const, volatile, and restrict. This restriction only applies if any of these 
qualifiers appears at the outermost level of a parameter type specification. For 
example, the following are declarations of the same function:

```cpp
int f(int);  
int f(const int);  
int f(volatile int);
```

Note that you can differentiate parameters with const, volatile and restrict 
qualifiers if you apply them within a parameter type specification. For example, 
the following declarations are not equivalent because const and volatile qualify 
int, rather than *, and thus are not at the outermost level of the parameter type 
specification.

```cpp
void g(int*);  
void g(const int*);  
void g(volatile int*);
```

The following declarations are also not equivalent:

```cpp
void g(float&);  
void g(const float&);  
void g(volatile float&);
```

• Function declarations with parameters that differ only because their default 
arguments differ. For example, the following are declarations of the same 
function:

```cpp
void f(int);  
void f(int i = 10);
```

• Multiple functions with extern "C" language-linkage and the same name, 
regardless of whether their parameter lists are different.

Related information

• “The using declaration and namespaces” on page 226
• “typedef definitions” on page 66
• “Type qualifiers” on page 68
• “Language linkage (C++ only)” on page 10

---

Overloading operators

You can redefine or overload the function of most built-in operators in C++. These 
operators can be overloaded globally or on a class-by-class basis. Overloaded 
operators are implemented as functions and can be member functions or global 
functions.

Operator overloading involving vector types is not supported.

An overloaded operator is called an operator function. You declare an operator 
function with the keyword operator preceding the operator. Overloaded operators 
are distinct from overloaded functions, but like overloaded functions, they are 
distinguished by the number and types of operands used with the operator.

Consider the standard + (plus) operator. When this operator is used with operands 
of different standard types, the operators have slightly different meanings. For 
example, the addition of two integers is not implemented in the same way as the 
addition of two floating-point numbers. C++ allows you to define your own 
meanings for the standard C++ operators when they are applied to class types. In 
the following example, a class called `complex` is defined to model complex numbers,
and the + (plus) operator is redefined in this class to add two complex numbers.
// This example illustrates overloading the plus (+) operator.

#include <iostream>
using namespace std;

class complx
{
    double real,
            imag;
public:
    complx( double real = 0., double imag = 0.); // constructor
    complx operator+(const complx&) const;       // operator+
};

// define constructor
complx::complx( double r, double i )
{
    real = r; imag = i;
}

// define overloaded + (plus) operator
complx complx::operator+(const complx& c) const
{
    complx result;
    result.real = (this->real + c.real);
    result.imag = (this->imag + c.imag);
    return result;
}

int main()
{
    complx x(4,4);
    complx y(6,6);
    complx z = x + y; // calls complx::operator+
}

You can overload any of the following operators:

<table>
<thead>
<tr>
<th>+</th>
<th>−</th>
<th>*</th>
<th>/</th>
<th>%</th>
<th>^</th>
<th>&amp;</th>
<th></th>
<th>~</th>
</tr>
</thead>
</table>

<table>
<thead>
<tr>
<th>!</th>
<th>=</th>
<th>&lt;</th>
<th>&gt;</th>
<th>+=</th>
<th>−=</th>
<th>*=</th>
<th>/=</th>
<th>%=</th>
</tr>
</thead>
</table>

| ^= | &= | |= | <<= | >>= | == | != | | != |
|---|---|---|---|---|---|---|---|---|

<table>
<thead>
<tr>
<th>&lt;=</th>
<th>&gt;=</th>
<th>&amp; &amp;</th>
<th></th>
<th>+</th>
<th>−</th>
<th></th>
<th>-&gt;</th>
<th>~&gt;</th>
</tr>
</thead>
</table>

| ( | ) | [ | ] | new | delete | new[] | delete[] |
|---|---|---|---|---|---|---|---|---|

where () is the function call operator and [] is the subscript operator.

You can overload both the unary and binary forms of the following operators:

<table>
<thead>
<tr>
<th>+</th>
<th>−</th>
<th>*</th>
<th>&amp;</th>
</tr>
</thead>
</table>

You cannot overload the following operators:

<table>
<thead>
<tr>
<th>.</th>
<th>.*</th>
<th>::</th>
<th>? :</th>
</tr>
</thead>
</table>

You cannot overload the preprocessor symbols # and ##.

An operator function can be either a nonstatic member function, or a nonmember function with at least one parameter that has class, reference to class, enumeration, or reference to enumeration type.

You cannot change the precedence, grouping, or the number of operands of an operator.
An overloaded operator (except for the function call operator) cannot have default arguments or an ellipsis in the argument list.

You must declare the overloaded =, [], (), and -> operators as nonstatic member functions to ensure that they receive lvalues as their first operands.

The operators new, delete, new[], and delete[] do not follow the general rules described in this section.

All operators except the = operator are inherited.

**Related information**
- “Free store” on page 315

### Overloading unary operators

You overload a unary operator with either a nonstatic member function that has no parameters, or a nonmember function that has one parameter. Suppose a unary operator @ is called with the statement @t, where t is an object of type T. A nonstatic member function that overloads this operator would have the following form:

```cpp
return_type operator@()
```

A nonmember function that overloads the same operator would have the following form:

```cpp
return_type operator@(T)
```

An overloaded unary operator may return any type.

The following example overloads the ! operator:

```cpp
#include <iostream>
using namespace std;

struct X { };
void operator!(X) {
    cout << "void operator!(X)" << endl;
}

struct Y {
    void operator() {
        cout << "void Y::operator()" << endl;
    }
};

struct Z { };

int main() {
    X ox; Y oy; Z oz;
    !ox;
    !oy;
    // !oz;
}
```

The following is the output of the above example:

```cpp
void operator!(X)
void Y::operator()
The operator function call !ox is interpreted as !operator!(X). The call !oy is interpreted as Y::operator!(). (The compiler would not allow !oz because the ! operator has not been defined for class Z.)

**Related information**

- "Unary expressions” on page 131

## Overloading increment and decrement operators

You overload the prefix increment operator ++ with either a nonmember function operator that has one argument of class type or a reference to class type, or with a member function operator that has no arguments.

In the following example, the increment operator is overloaded in both ways:

```cpp
class X {
public:
    // member prefix ++x
    void operator++() { }
};
class Y { }

// non-member prefix ++y
void operator++(Y&) { }

int main() {
    X x;
    Y y;

    // calls x.operator++()
    ++x;

    // explicit call, like ++x
    x.operator++();

    // calls operator++(y)
    ++y;

    // explicit call, like ++y
    operator++(y);
}
```

The postfix increment operator ++ can be overloaded for a class type by declaring a nonmember function operator operator++() with two arguments, the first having class type and the second having type int. Alternatively, you can declare a member function operator operator++() with one argument having type int. The compiler uses the int argument to distinguish between the prefix and postfix increment operators. For implicit calls, the default value is zero.

For example:

```cpp
class X {
public:
    // member postfix x++
    void operator++(int) { }
};
class Y { }

// nonmember postfix y++
void operator++(Y&, int) { }
```
int main() {
    X x;
    Y y;

    // calls x.operator++(0)
    // default argument of zero is supplied by compiler
    x++;
    // explicit call to member postfix x++
    x.operator++(0);

    // calls operator++(y, 0)
    y++;
    // explicit call to non-member postfix y++
    operator++(y, 0);
}

The prefix and postfix decrement operators follow the same rules as their increment counterparts.

Related information
- "Increment operator ++" on page 132
- "Decrement operator --" on page 133

Overloading binary operators

You overload a binary unary operator with either a nonstatic member function that has one parameter, or a nonmember function that has two parameters. Suppose a binary operator @ is called with the statement t @ u, where t is an object of type T, and u is an object of type U. A nonstatic member function that overloads this operator would have the following form:

return_type operator@(T)

A nonmember function that overloads the same operator would have the following form:

return_type operator@(T, U)

An overloaded binary operator may return any type.

The following example overloads the * operator:

struct X {
    // member binary operator
    void operator*(int) { }  // member binary operator
};

// non-member binary operator
void operator*(X, float) { }

int main() {
    X x;
    int y = 10;
    float z = 10;
    x * y;
    x * z;
}
The call \( x \times y \) is interpreted as \( x.\text{operator\ast}(y) \). The call \( x \times z \) is interpreted as \( \text{operator\ast}(x, z) \).

Related information
- "Binary expressions" on page 147

Overloading assignments
You overload the assignment operator, \( \text{operator\ast} \), with a nonstatic member function that has only one parameter. You cannot declare an overloaded assignment operator that is a nonmember function. The following example shows how you can overload the assignment operator for a particular class:

```cpp
struct X {
    int data;
    X& operator=(X& a) { return a; }
    X& operator=(int a) {
        data = a;
        return *this;
    }
};

int main() {
    X x1, x2;
    x1 = x2;  // call x1.operator=(x2)
    x1 = 5;  // call x1.operator=(5)
}
```

The assignment \( x1 = x2 \) calls the copy assignment operator \( X\& \text{X::operator\ast}(X\&) \). The assignment \( x1 = 5 \) calls the copy assignment operator \( X\& \text{X::operator\ast}(\text{int}) \). The compiler implicitly declares a copy assignment operator for a class if you do not define one yourself. Consequently, the copy assignment operator (\( \text{operator\ast} \)) of a derived class hides the copy assignment operator of its base class.

However, you can declare any copy assignment operator as virtual. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    A& operator=(char) {
        cout << "A\& A::operator=(char)" << endl;
        return *this;
    }
    virtual A& operator=(const A&) {
        cout << "A\& A::operator=(const A\&)" << endl;
        return *this;
    }
};

struct B : A {
    B& operator=(char) {
        cout << "B\& B::operator=(char)" << endl;
        return *this;
    }
    virtual B& operator=(const A&) {
        cout << "B\& B::operator=(const A\&)" << endl;
        return *this;
    }
};

struct C : B { }

int main() {
```
The following is the output of the above example:

```cpp
A & A::operator=(char)
B & B::operator=(const A&)
```

The assignment `*ap1 = 'z'` calls `A & A::operator=(char)`. Because this operator has not been declared virtual, the compiler chooses the function based on the type of the pointer `ap1`. The assignment `*ap2 = b2` calls `B & B::operator=(const &A)`. Because this operator has been declared virtual, the compiler chooses the function based on the type of the object that the pointer `ap1` points to. The compiler would not allow the assignment `c1 = 'z'` because the implicitly declared copy assignment operator declared in class `C` hides `B & B::operator=(char)`. Related information
- "Copy assignment operators" on page 326
- "Assignment expressions" on page 158

### Overloading function calls

The function call operator, when overloaded, does not modify how functions are called. Rather, it modifies how the operator is to be interpreted when applied to objects of a given type.

You overload the function call operator, `operator()`, with a nonstatic member function that has any number of parameters. If you overload a function call operator for a class its declaration will have the following form:

```
return_type operator()(parameter_list)
```

Unlike all other overloaded operators, you can provide default arguments and ellipses in the argument list for the function call operator.

The following example demonstrates how the compiler interprets function call operators:

```cpp
struct A {
    void operator()(int a, char b, ...) {}
    void operator()(char c, int d = 20) {}
};

int main() {
    A a;
    a(5, 'z', 'a', 0);
    a('z');
    // a();
}
```

The function call `a(5, 'z', 'a', 0)` is interpreted as `a.operator()(5, 'z', 'a', 0)`. This calls `A::operator()(int a, char b, ...)`. The function call `a('z')` is interpreted as `a.operator()('z')`. This calls `void A::operator()(char c, int d =`
The compiler would not allow the function call `a()` because its argument list does not match any function call parameter list defined in class `A`.

The following example demonstrates an overloaded function call operator:

```cpp
class Point {
    private:
        int x, y;
    public:
        Point() : x(0), y(0) {}
        Point& operator()(int dx, int dy) {
            x += dx;
            y += dy;
            return *this;
        }
};

int main() {
    Point pt;
    // Offset this coordinate x with 3 points
    // and coordinate y with 2 points.
    pt(3, 2);
}
```

The above example reinterprets the function call operator for objects of class `Point`. If you treat an object of `Point` like a function and pass it two integer arguments, the function call operator will add the values of the arguments you passed to `Point::x` and `Point::y` respectively.

**Related information**

- "Function call operator ()" on page 121

**Overloading subscripting**

You overload `operator[]` with a nonstatic member function that has only one parameter. The following example is a simple array class that has an overloaded subscripting operator. The overloaded subscripting operator throws an exception if you try to access the array outside of its specified bounds:

```cpp
#include <iostream>
using namespace std;

template <class T> class MyArray {
    private:
        T* storage;
        int size;
    public:
        MyArray(int arg = 10) {
            storage = new T[arg];
            size = arg;
        }

        ~MyArray() {
            delete[] storage;
            storage = 0;
        }

        T& operator[](const int location) throw (const char*);
};

template <class T> T& MyArray<T>::operator[](const int location) throw (const char*){
    if (location < 0 || location >= size) throw "Invalid array access";
    else return storage[location];
}
```
int main() {
    try {
        MyArray<int> x(13);
        x[0] = 45;
        x[1] = 2435;
        cout << x[0] << endl;
        cout << x[1] << endl;
        x[13] = 84;
    } catch (const char* e) {
        cout << e << endl;
    }
}

The following is the output of the above example:
45
2435
Invalid array access

The expression x[1] is interpreted as x.operator[](1) and calls int& MyArray<int>::operator[](const int).

Related information
- "Array subscripting operator [ ]" on page 122

Overloading class member access

You overload operator-> with a nonstatic member function that has no parameters. The following example demonstrates how the compiler interprets overloaded class member access operators:

```c++
struct Y {
    void f() {}
};

struct X {
    Y* ptr;
    Y* operator->() {
        return ptr;
    }
};

int main() {
    X x;
    x->f();
}
```

The statement x->f() is interpreted as (x.operator->())->f().

The operator-> is used (often in conjunction with the pointer-dereference operator) to implement "smart pointers." These pointers are objects that behave like normal pointers except they perform other tasks when you access an object through them, such as automatic object deletion (either when the pointer is destroyed, or the pointer is used to point to another object), or reference counting (counting the number of smart pointers that point to the same object, then automatically deleting the object when that count reaches zero).

One example of a smart pointer is included in the C++ Standard Library called auto_ptr. You can find it in the <memory> header. The auto_ptr class implements automatic object deletion.
Related information

- "Arrow operator –>" on page 123

Overload resolution

The process of selecting the most appropriate overloaded function or operator is called overload resolution.

Suppose that f is an overloaded function name. When you call the overloaded function f(), the compiler creates a set of candidate functions. This set of functions includes all of the functions named f that can be accessed from the point where you called f(). The compiler may include as a candidate function an alternative representation of one of those accessible functions named f to facilitate overload resolution.

After creating a set of candidate functions, the compiler creates a set of viable functions. This set of functions is a subset of the candidate functions. The number of parameters of each viable function agrees with the number of arguments you used to call f().

The compiler chooses the best viable function, the function declaration that the C++ run-time environment will use when you call f(), from the set of viable functions. The compiler does this by implicit conversion sequences. An implicit conversion sequence is the sequence of conversions required to convert an argument in a function call to the type of the corresponding parameter in a function declaration. The implicit conversion sequences are ranked; some implicit conversion sequences are better than others. The best viable function is the one whose parameters all have either better or equal-ranked implicit conversion sequences than all of the other viable functions. The compiler will not allow a program in which the compiler was able to find more than one best viable function. Implicit conversion sequences are described in more detail in "Implicit conversion sequences" on page 242.

When a variable length array is a function parameter, the leftmost array dimension does not distinguish functions among candidate functions. In the following, the second definition of f is not allowed because void f(int []) has already been defined.

```c++
void f(int a[*]) {}
void f(int a[5]) {} // illegal
```

However, array dimensions other than the leftmost in a variable length array do differentiate candidate functions when the variable length array is a function parameter. For example, the overload set for function f might comprise the following:

```c++
void f(int a[][5]) {}
void f(int a[][4]) {}
void f(int a[][g]) {} // assume g is a global int
```

but cannot include

```c++
void f(int a[][g2]) {} // illegal, assuming g2 is a global int
```

because having candidate functions with second-level array dimensions g and g2 creates ambiguity about which function f should be called: neither g nor g2 is known at compile time.
You can override an exact match by using an explicit cast. In the following example, the second call to f() matches with f(void*):

```c
void f(int) {};
void f(void*) {};

int main() {
    f(0xaabb);  // matches f(int);
    f((void*) 0xaabb);  // matches f(void*)
}
```

**Implicit conversion sequences**

An *implicit conversion sequence* is the sequence of conversions required to convert an argument in a function call to the type of the corresponding parameter in a function declaration.

The compiler will try to determine an implicit conversion sequence for each argument. It will then categorize each implicit conversion sequence in one of three categories and rank them depending on the category. The compiler will not allow any program in which it cannot find an implicit conversion sequence for an argument.

The following are the three categories of conversion sequences in order from best to worst:

- **Standard conversion sequences**
- **User-defined conversion sequences**
- **Ellipsis conversion sequences**

**Note:** Two standard conversion sequences or two user-defined conversion sequences may have different ranks.

**Standard conversion sequences**

Standard conversion sequences are categorized in one of three ranks. The ranks are listed in order from best to worst:

- **Exact match:** This rank includes the following conversions:
  - Identity conversions
  - Lvalue-to-rvalue conversions
  - Array-to-pointer conversions
  - Qualification conversions
- **Promotion:** This rank includes integral and floating point promotions.
- **Conversion:** This rank includes the following conversions:
  - Integral and floating-point conversions
  - Floating-integral conversions
  - Pointer conversions
  - Pointer-to-member conversions
  - Boolean conversions

The compiler ranks a standard conversion sequence by its worst-ranked standard conversion. For example, if a standard conversion sequence has a floating-point conversion, then that sequence has conversion rank.

**Related information**

- “Lvalue-to-rvalue conversions” on page 110
- “Pointer conversions” on page 110
User-defined conversion sequences

A user-defined conversion sequence consists of the following:

- A standard conversion sequence
- A user-defined conversion
- A second standard conversion sequence

A user-defined conversion sequence A is better than a user-defined conversion sequence B if the both have the same user-defined conversion function or constructor, and the second standard conversion sequence of A is better than the second standard conversion sequence of B.

Ellipsis conversion sequences

An ellipsis conversion sequence occurs when the compiler matches an argument in a function call with a corresponding ellipsis parameter.

Resolving addresses of overloaded functions

If you use an overloaded function name f without any arguments, that name can refer to a function, a pointer to a function, a pointer to member function, or a specialization of a function template. Because you did not provide any arguments, the compiler cannot perform overload resolution the same way it would for a function call or for the use of an operator. Instead, the compiler will try to choose the best viable function that matches the type of one of the following expressions, depending on where you have used f:

- An object or reference you are initializing
- The left side of an assignment
- A parameter of a function or a user-defined operator
- The return value of a function, operator, or conversion
- An explicit type conversion

If the compiler chose a declaration of a nonmember function or a static member function when you used f, the compiler matched the declaration with an expression of type pointer-to-function or reference-to-function. If the compiler chose a declaration of a nonstatic member function, the compiler matched that declaration with an expression of type pointer-to-member function. The following example demonstrates this:

```cpp
struct X {
    int f(int) { return 0; }
    static int f(char) { return 0; }
};

int main() {
    int (X::*a)(int) = &X::f;
    // int (*b)(int) = &X::f;
}
```

The compiler will not allow the initialization of the function pointer b. No nonmember function or static function of type int(int) has been declared.

If f is a template function, the compiler will perform template argument deduction to determine which template function to use. If successful, it will add that function
to the list of viable functions. If there is more than one function in this set, including a non-template function, the compiler will eliminate all template functions from the set and choose the non-template function. If there are only template functions in this set, the compiler will choose the most specialized template function. The following example demonstrates this:

```c
// Example code

int main() {
    int (*a)(int) = f;
    a(1);
}
```

The function call `a(1)` calls `int f(int)`.

**Related information**

- “Pointers to functions” on page 219
- “Pointers to members” on page 260
- “Function templates” on page 340
- “Explicit specialization” on page 351
Chapter 11. Classes (C++ only)

A class is a mechanism for creating user-defined data types. It is similar to the C language structure data type. In C, a structure is composed of a set of data members. In C++, a class type is like a C structure, except that a class is composed of a set of data members and a set of operations that can be performed on the class.

In C++, a class type can be declared with the keywords union, struct, or class. A union object can hold any one of a set of named members. Structure and class objects hold a complete set of members. Each class type represents a unique set of class members including data members, member functions, and other type names. The default access for members depends on the class key:

- The members of a class declared with the keyword class are private by default.
  A class is inherited privately by default.
- The members of a class declared with the keyword struct are public by default.
  A structure is inherited publicly by default.
- The members of a union (declared with the keyword union) are public by default.
  A union cannot be used as a base class in derivation.

Once you create a class type, you can declare one or more objects of that class type. For example:

```cpp
class X
{
    /* define class members here */
};
int main()
{
    X xobject1; // create an object of class type X
    X xobject2; // create another object of class type X
}
```

You may have polymorphic classes in C++. Polymorphism is the ability to use a function name that appears in different classes (related by inheritance), without knowing exactly the class the function belongs to at compile time.

C++ allows you to redefine standard operators and functions through the concept of overloading. Operator overloading facilitates data abstraction by allowing you to use classes as easily as built-in types.

Related information

- “Structures and unions” on page 55
- Chapter 12, “Class members and friends (C++ only),” on page 255
- Chapter 13, “Inheritance (C++ only),” on page 277
- Chapter 10, “Overloading (C++ only),” on page 229
- “Virtual functions” on page 295
Declaring class types

A class declaration creates a unique type class name.

A class specifier is a type specifier used to declare a class. Once a class specifier has been seen and its members declared, a class is considered to be defined even if the member functions of that class are not yet defined.

Class specifier syntax

```
class class_name [ : base_clause ] { member_list }
```

The `class_name` is a unique identifier that becomes a reserved word within its scope. Once a class name is declared, it hides other declarations of the same name within the enclosing scope.

The `member_list` specifies the class members, both data and functions, of the class `class_name`. If the `member_list` of a class is empty, objects of that class have a nonzero size. You can use a `class_name` within the `member_list` of the class specifier itself as long as the size of the class is not required.

The `base_clause` specifies the base class or classes from which the class `class_name` inherits members. If the `base_clause` is not empty, the class `class_name` is called a derived class.

A structure is a class declared with the `class_key` `struct`. The members and base classes of a structure are public by default. A union is a class declared with the `class_key` `union`. The members of a union are public by default; a union holds only one data member at a time.

An aggregate class is a class that has no user-defined constructors, no private or protected non-static data members, no base classes, and no virtual functions.

Related information

- “Class member lists” on page 255
- “Derivation” on page 279

Using class objects

You can use a class type to create instances or objects of that class type. For example, you can declare a class, structure, and union with class names `X`, `Y`, and `Z` respectively:

```
class X {
    // members of class X
};

struct Y {
    // members of struct Y
};

union Z {
    // members of union Z
};
```
You can then declare objects of each of these class types. Remember that classes, structures, and unions are all types of C++ classes.

```cpp
int main()
{
    X xobj;       // declare a class object of class type X
    Y yobj;       // declare a struct object of class type Y
    Z zobj;       // declare a union object of class type Z
}
```

In C++, unlike C, you do not need to precede declarations of class objects with the keywords `union`, `struct`, and `class` unless the name of the class is hidden. For example:

```cpp
struct Y { /* ... */ };  
class X { /* ... */ };  
int main()
{
    int X;               // hides the class name X
    Y yobj;             // valid
    X xobj;             // error, class name X is hidden
    class X xobj;       // valid
}
```

When you declare more than one class object in a declaration, the declarators are treated as if declared individually. For example, if you declare two objects of class `S` in a single declaration:

```cpp
class S { /* ... */ };  
int main()
{
    S S,T; // declare two objects of class type S
}
```

this declaration is equivalent to:

```cpp
class S { /* ... */ };  
int main()
{
    S S;  
    class S T;       // keyword class is required
                      // since variable S hides class type S
}
```

but is not equivalent to:

```cpp
class S { /* ... */ };  
int main()
{
    S S;  
    S T;       // error, S class type is hidden
}
```

You can also declare references to classes, pointers to classes, and arrays of classes. For example:

```cpp
class X { /* ... */ };  
struct Y { /* ... */ };  
union Z { /* ... */ };  
int main()
{
    X xobj;
    X &xref = xobj;       // reference to class object of type X
    Y *ypt;              // pointer to struct object of type Y
    Z zarray[10];        // array of 10 union objects of type Z
}```
You can initialize classes in external, static, and automatic definitions. The initializer contains an = (equal sign) followed by a brace-enclosed, comma-separated list of values. You do not need to initialize all members of a class.

Objects of class types that are not copy restricted can be assigned, passed as arguments to functions, and returned by functions.

**Related information**
- “Structures and unions” on page 55
- “References (C++ only)” on page 88
- “Scope of class names” on page 249

## Classes and structures

The C++ class is an extension of the C language structure. Because the only difference between a structure and a class is that structure members have public access by default and class members have private access by default, you can use the keywords class or struct to define equivalent classes.

For example, in the following code fragment, the class X is equivalent to the structure Y:

```cpp
class X {
    // private by default
    int a;

public:
    // public member function
    int f() { return a = 5; };
};

struct Y {
    // public by default
    int f() { return a = 5; };

private:
    // private data member
    int a;
};
```

If you define a structure and then declare an object of that structure using the keyword class, the members of the object are still public by default. In the following example, main() has access to the members of obj_X even though obj_X has been declared using an elaborated type specifier that uses the class key class:

```cpp
#include <iostream>
using namespace std;

struct X {
    int a;
    int b;
};

class X obj_X;

int main() {
```
The following is the output of the above example:
Here are a and b: 0 1

Related information
• “Structures and unions” on page 55

Scope of class names

A class declaration introduces the class name into the scope where it is declared. Any class, object, function or other declaration of that name in an enclosing scope is hidden.

If a class name is declared in the same scope as a function, enumerator, or object with the same name, you must refer to that class using an elaborated type specifier:

Elaborated type specifier syntax

```
class
struct
union
typename

::

namespace name

::

template

::
	nested_name_specifier

identifier

template

template_name

::

nested_name_specifier

identifier

::

template_name

::

nested_name_specifier
```

Nested name specifier:

The following example must use an elaborated type specifier to refer to class A because this class is hidden by the definition of the function A():
```cpp
class A {}
void A (class A*) {}

int main()
{
    class A* x;
    A(x);
}
```

The declaration class A* x is an elaborated type specifier. Declaring a class with the same name of another function, enumerator, or object as demonstrated above is not recommended.

An elaborated type specifier can also be used in the incomplete declaration of a class type to reserve the name for a class type within the current scope.

Related information
• “Class scope (C++ only)” on page 5
Incomplete class declarations

An incomplete class declaration is a class declaration that does not define any class members. You cannot declare any objects of the class type or refer to the members of a class until the declaration is complete. However, an incomplete declaration allows you to make specific references to a class prior to its definition as long as the size of the class is not required.

For example, you can define a pointer to the structure first in the definition of the structure second. Structure first is declared in an incomplete class declaration prior to the definition of second, and the definition of oneptr in structure second does not require the size of first:

```c
struct first; // incomplete declaration of struct first

struct second // complete declaration of struct second
{
    first* oneptr; // pointer to struct first refers to // struct first prior to its complete // declaration

    first one; // error, you cannot declare an object of // an incompletely declared class type
    int x, y;
};

struct first // complete declaration of struct first
{
    second two; // define an object of class type second
    int z;
};
```

However, if you declare a class with an empty member list, it is a complete class declaration. For example:

```c
class X; // incomplete class declaration
class Z {}; // empty member list
class Y
{
    public:
        X yobj; // error, cannot create an object of an // incomplete class type
        Z zobj; // valid
};
```

Related information

- “Class member lists” on page 255

Nested classes

A nested class is declared within the scope of another class. The name of a nested class is local to its enclosing class. Unless you use explicit pointers, references, or object names, declarations in a nested class can only use visible constructs, including type names, static members, and enumerators from the enclosing class and global variables.

Member functions of a nested class follow regular access rules and have no special access privileges to members of their enclosing classes. Member functions of the enclosing class have no special access to members of a nested class. The following example demonstrates this:
class A {
    int x;
}

class B {}

class C {
    // The compiler cannot allow the following
    // declaration because A::B is private:
    //   B b;
    int y;
    void f(A* p, int i) {
        // The compiler cannot allow the following
        // statement because A::x is private:
        //   p->x = i;
    }
};

void g(C* p) {
    // The compiler cannot allow the following
    // statement because C::y is private:
    //   int z = p->y;
}
};

int main() {
    // The compiler would not allow the declaration of object b because class A::B is private. The compiler would not allow the statement p->x = i because A::x is private. The compiler would not allow the statement int z = p->y because C::y is private.

    You can define member functions and static data members of a nested class in namespace scope. For example, in the following code fragment, you can access the static members x and y and member functions f() and g() of the nested class nested by using a qualified type name. Qualified type names allow you to define a typedef to represent a qualified class name. You can then use the typedef with the :: (scope resolution) operator to refer to a nested class or class member, as shown in the following example:

class outside {
public:
    class nested {
        public:
            static int x;
            static int y;
            int f();
            int g();
    }
};
int outside::nested::x = 5;
int outside::nested::f() { return 0; }

typedef outside::nested outnest; // define a typedef
int outnest::y = 10; // use typedef with ::
int outnest::g() { return 0; }

    However, using a typedef to represent a nested class name hides information and may make the code harder to understand.
You cannot use a typedef name in an elaborated type specifier. To illustrate, you cannot use the following declaration in the above example:

```cpp
class outnest obj;
```

A nested class may inherit from private members of its enclosing class. The following example demonstrates this:

```cpp
class A {
    private:
        class B {
        };
    B *z;

class C : private B {
    private:
        B y;
        // A::B y2;
        C *x;
        // A::C *x2;
    };
};
```

The nested class A::C inherits from A::B. The compiler does not allow the declarations A::B y2 and A::C *x2 because both A::B and A::C are private.

Related information

- "Class scope (C++ only)" on page 5
- "Scope of class names" on page 249
- "Member access" on page 269
- "Static members" on page 265

Local classes

A local class is declared within a function definition. Declarations in a local class can only use type names, enumerations, static variables from the enclosing scope, as well as external variables and functions.

For example:

```cpp
int x;                     // global variable
void f()                   // function definition
{
    static int y;         // static variable y can be used by
                            // local class
    int x;                // auto variable x cannot be used by
                            // local class
    extern int g();      // external function g can be used by
                            // local class

class local               // local class
    {
        int g() { return x; }    // error, local variable x
                                   // cannot be used by g
        int h() { return y; }    // valid, static variable y
        int k() { return ::x; }  // valid, global x
        int l() { return g(); }  // valid, extern function g
    };
}

int main()
{
    local* z;                // error: the class local is not visible
    // ...
}
```
Member functions of a local class have to be defined within their class definition, if they are defined at all. As a result, member functions of a local class are inline functions. Like all member functions, those defined within the scope of a local class do not need the keyword inline.

A local class cannot have static data members. In the following example, an attempt to define a static member of a local class causes an error:

```cpp
void f()
{
    class local
    {
        int f(); // error, local class has noninline
        // member function
        int g() {return 0;} // valid, inline member function
        static int a; // error, static is not allowed for
        // local class
        int b; // valid, nonstatic variable
    };
    // . . .
}
```

An enclosing function has no special access to members of the local class.

**Related information**
- ["Member functions" on page 257](#)
- ["The inline function specifier" on page 197](#)

**Local type names**

Local type names follow the same scope rules as other names. Type names defined within a class declaration have class scope and cannot be used outside their class without qualification.

If you use a class name, typedef name, or a constant name that is used in a type name, in a class declaration, you cannot redefine that name after it is used in the class declaration.

For example:

```cpp
int main ()
{
    typedef double db;
    struct st
    {
        db x;
        typedef int db; // error
        db y;
    };
}
```

The following declarations are valid:

```cpp
typedef float T;
class s {
    typedef int T;
    void f(const T);
};
```

Here, function f() takes an argument of type s::T. However, the following declarations, where the order of the members of s has been reversed, cause an error:
typedef float T;
class s {
    void f(const T);
    typedef int T;
};

In a class declaration, you cannot redefine a name that is not a class name, or a typedef name to a class name or typedef name once you have used that name in the class declaration.

Related information
- “Scope” on page 2
- “typedef definitions” on page 66
Chapter 12. Class members and friends (C++ only)

This section discusses the declaration of class members with respect to the information hiding mechanism and how a class can grant functions and classes access to its nonpublic members by the use of the friend mechanism. C++ expands the concept of information hiding to include the notion of having a public class interface but a private implementation. It is the mechanism for limiting direct access to the internal representation of a class type by functions in a program.

Related information
- “Member access” on page 269
- “Inherited member access” on page 282

Class member lists

An optional member list declares subobjects called class members. Class members can be data, functions, nested types, and enumerators.

Class member list syntax

```
member_declaration
: = 0 constant_expression

access_specifier
```

The member list follows the class name and is placed between braces. The following applies to member lists, and members of member lists:

- A member declaration or a member definition may be a declaration or definition of a data member, member function, nested type, or enumeration. (The enumerators of an enumeration defined in a class member list are also members of the class.)
- A member list is the only place where you can declare class members.
- Friend declarations are not class members but must appear in member lists.
- The member list in a class definition declares all the members of a class; you cannot add members elsewhere.
- You cannot declare a member twice in a member list.
- You may declare a data member or member function as static but not auto, extern, or register.
- You may declare a nested class, a member class template, or a member function, and define it outside the class.
- You must define static data members outside the class.
- Nonstatic members that are class objects must be objects of previously defined classes; a class A cannot contain an object of class A, but it can contain a pointer or reference to an object of class A.
- You must specify all dimensions of a nonstatic array member.
A constant initializer (= constant_expression) may only appear in a class member of integral or enumeration type that has been declared static.

A pure specifier (= 0) indicates that a function has no definition. It is only used with member functions declared as virtual and replaces the function definition of a member function in the member list.

An access specifier is one of public, private, or protected.

A member declaration declares a class member for the class containing the declaration.

The order of allocation of nonstatic class members separated by an access specifier is implementation-dependent. The compiler allocates class members in the order in which they are declared.

Suppose A is a name of a class. The following class members of A must have a name different from A:

- All data members
- All type members
- All enumerators of enumerated type members
- All members of all anonymous union members

Related information

- "Declaring class types" on page 246
- "Member access" on page 269
- "Inherited member access" on page 282
- "Static members" on page 265

Data members

Data members include members that are declared with any of the fundamental types, as well as other types, including pointer, reference, array types, bit fields, and user-defined types. You can declare a data member the same way as a variable, except that explicit initializers are not allowed inside the class definition. However, a const static data member of integral or enumeration type may have an explicit initializer.

If an array is declared as a nonstatic class member, you must specify all of the dimensions of the array.

A class can have members that are of a class type or are pointers or references to a class type. Members that are of a class type must be of a class type that has been previously declared. An incomplete class type can be used in a member declaration as long as the size of the class is not needed. For example, a member can be declared that is a pointer to an incomplete class type.

A class X cannot have a member that is of type X, but it can contain pointers to X, references to X, and static objects of X. Member functions of X can take arguments of type X and have a return type of X. For example:

```c
class X
{
    X();
    X *xptr;
};
```
```cpp
X &xref;
static X xcount;
X xfunc(X);
);

Related information
• “Member access” on page 269
• “Inherited member access” on page 282
• “Static members” on page 265

Member functions

Member functions are operators and functions that are declared as members of a class. Member functions do not include operators and functions declared with the friend specifier. These are called friends of a class. You can declare a member function as static; this is called a static member function. A member function that is not declared as static is called a nonstatic member function.

The definition of a member function is within the scope of its enclosing class. The body of a member function is analyzed after the class declaration so that members of that class can be used in the member function body, even if the member function definition appears before the declaration of that member in the class member list. When the function add() is called in the following example, the data variables a, b, and c can be used in the body of add().

class x
{
public:
   int add() // inline member function add
   {return a+b+c;};
private:
   int a,b,c;
};

Inline member functions

You may either define a member function inside its class definition, or you may define it outside if you have already declared (but not defined) the member function in the class definition.

A member function that is defined inside its class member list is called an inline member function. Member functions containing a few lines of code are usually declared inline. In the above example, add() is an inline member function. If you define a member function outside of its class definition, it must appear in a namespace scope enclosing the class definition. You must also qualify the member function name using the scope resolution (::) operator.

An equivalent way to declare an inline member function is to either declare it in the class with the inline keyword (and define the function outside of its class) or to define it outside of the class declaration using the inline keyword.

In the following example, member function Y::f() is an inline member function:

```
The following example is equivalent to the previous example; `Y::f()` is an inline member function:

```cpp
struct Y {
    private:
        char* a;
    public:
        char* f();
};

inline char* Y::f() { return a; }
```

The `inline` specifier does not affect the linkage of a member or nonmember function: linkage is external by default.

Member functions of a local class must be defined within their class definition. As a result, member functions of a local class are implicitly inline functions. These inline member functions have no linkage.

**Related information**
- “Friends” on page 272
- “Static member functions” on page 268
- “The inline function specifier” on page 197
- “Local classes” on page 252

**Constant and volatile member functions**

A member function declared with the `const` qualifier can be called for constant and nonconstant objects. A nonconstant member function can only be called for a nonconstant object. Similarly, a member function declared with the `volatile` qualifier can be called for volatile and nonvolatile objects. A nonvolatile member function can only be called for a nonvolatile object.

**Related information**
- “Type qualifiers” on page 68
- “The this pointer” on page 261

**Virtual member functions**

Virtual member functions are declared with the keyword `virtual`. They allow dynamic binding of member functions. Because all virtual functions must be member functions, virtual member functions are simply called `virtual functions`.

If the definition of a virtual function is replaced by a pure specifier in the declaration of the function, the function is said to be declared pure. A class that has at least one pure virtual function is called an `abstract class`.

**Related information**
- “Virtual functions” on page 295
- “Abstract classes” on page 300

**Special member functions**

`Special member functions` are used to create, destroy, initialize, convert, and copy class objects. These include the following:

- Constructors
- Destructors
• Conversion constructors
• Conversion functions
• Copy constructors

For full descriptions of these functions, see Chapter 14, “Special member functions (C++ only),” on page 303.

**Member scope**

Member functions and static members can be defined outside their class declaration if they have already been declared, but not defined, in the class member list. Nonstatic data members are defined when an object of their class is created. The declaration of a static data member is not a definition. The declaration of a member function is a definition if the body of the function is also given.

Whenever the definition of a class member appears outside of the class declaration, the member name must be qualified by the class name using the :: (scope resolution) operator.

The following example defines a member function outside of its class declaration.

```cpp
#include <iostream>
using namespace std;

struct X {
    int a, b;
    // member function declaration only
    int add();
};

// global variable
int a = 10;

// define member function outside its class declaration
int X::add() { return a + b; }

int main() {
    int answer;
    X xobject;
    xobject.a = 1;
    xobject.b = 2;
    answer = xobject.add();
    cout << xobject.a << " + " << xobject.b << " = " << answer << endl;
}
```

The output for this example is: 1 + 2 = 3

All member functions are in class scope even if they are defined outside their class declaration. In the above example, the member function add() returns the data member a, not the global variable a.

The name of a class member is local to its class. Unless you use one of the class access operators, . (dot), or -> (arrow), or :: (scope resolution) operator, you can only use a class member in a member function of its class and in nested classes. You can only use types, enumerations and static members in a nested class without qualification with the :: operator.

The order of search for a name in a member function body is:
1. Within the member function body itself
2. Within all the enclosing classes, including inherited members of those classes
3. Within the lexical scope of the body declaration

The search of the enclosing classes, including inherited members, is demonstrated in the following example:

class A { /* ... */ };  
class B { /* ... */ };  
class C { /* ... */ };  
class Z : A {   
    class Y : B {   
        class X : C { int f(); /* ... */ };   
    }; 
    int Z::*Y::X f() 
    { 
        char j; 
        return 0; 
    }
}

In this example, the search for the name j in the definition of the function f follows this order:
1. In the body of the function f
2. In X and in its base class C
3. In Y and in its base class B
4. In Z and in its base class A
5. In the lexical scope of the body of f. In this case, this is global scope.

Note that when the containing classes are being searched, only the definitions of the containing classes and their base classes are searched. The scope containing the base class definitions (global scope, in this example) is not searched.

Related information
• “Class scope (C++ only)” on page 5

Pointers to members

Pointers to members allow you to refer to nonstatic members of class objects. You cannot use a pointer to member to point to a static class member because the address of a static member is not associated with any particular object. To point to a static class member, you must use a normal pointer.

You can use pointers to member functions in the same manner as pointers to functions. You can compare pointers to member functions, assign values to them, and use them to call member functions. Note that a member function does not have the same type as a nonmember function that has the same number and type of arguments and the same return type.

Pointers to members can be declared and used as shown in the following example:

```c++
#include <iostream>
using namespace std;

class X {
public:
    int a;
    void f(int b) { 
        cout << "The value of b is " << b << endl; 
    }
};
```
int main() {

    // declare pointer to data member
    int X::*ptiptr = &X::a;

    // declare a pointer to member function
    void (X::*ptfptr) (int) = &X::f;

    // create an object of class type X
    X xobject;

    // initialize data member
    xobject.*ptiptr = 10;

    cout << "The value of a is " << xobject.*ptiptr << endl;

    // call member function
    (xobject.*ptfptr) (20);
}

The output for this example is:
The value of a is 10
The value of b is 20

To reduce complex syntax, you can declare a typedef to be a pointer to a member. A pointer to a member can be declared and used as shown in the following code fragment:

typedef int X::*my_pointer_to_member;
typedef void (X::*my_pointer_to_function) (int);

int main() {
    my_pointer_to_member ptiptr = &X::a;
    my_pointer_to_function ptfptr = &X::f;
    X xobject;
    xobject.*ptiptr = 10;
    cout << "The value of a is " << xobject.*ptiptr << endl;
    (xobject.*ptfptr) (20);
}

The pointer to member operators .* and ->* are used to bind a pointer to a member of a specific class object. Because the precedence of () (function call operator) is higher than .* and ->*, you must use parentheses to call the function pointed to by ptf.

Pointer-to-member conversion can occur when pointers to members are initialized, assigned, or compared. Note that pointer to a member is not the same as a pointer to an object or a pointer to a function.

Related information
- ["Pointer to member operators . * ->*(C++ only)"
  on page 155]

The this pointer

The keyword this identifies a special type of pointer. Suppose that you create an object named x of class A, and class A has a nonstatic member function f(). If you call the function x.f(), the keyword this in the body of f() stores the address of x. You cannot declare the this pointer or make assignments to it.

A static member function does not have a this pointer.
The type of the this pointer for a member function of a class type X, is X* const. If the member function is declared with the const qualifier, the type of the this pointer for that member function for class X, is const X* const.

A const this pointer can by used only with const member functions. Data members of the class will be constant within that function. The function is still able to change the value, but requires a const_cast to do so:

```cpp
void foo::p() const{
    member = 1; // illegal
    const_cast<int&> (member) = 1; // a bad practice but legal
}
```

A better technique would be to declare member mutable.

If the member function is declared with the volatile qualifier, the type of the this pointer for that member function for class X is volatile X* const. For example, the compiler will not allow the following:

```cpp
struct A {
    int a;
    int f() const { return a++; }
};
```

The compiler will not allow the statement a++ in the body of function f(). In the function f(), the this pointer is of type A* const. The function f() is trying to modify part of the object to which this points.

The this pointer is passed as a hidden argument to all nonstatic member function calls and is available as a local variable within the body of all nonstatic functions.

For example, you can refer to the particular class object that a member function is called for by using the this pointer in the body of the member function. The following code example produces the output a = 5:

```cpp
#include <iostream>
using namespace std;

struct X {
    private:
        int a;
    public:
        void Set_a(int a) {
            // The 'this' pointer is used to retrieve 'xobj.a'
            // hidden by the automatic variable 'a'
            this->a = a;
        }
        void Print_a() { cout << "a = " << a << endl; }
};

int main() {
    X xobj;
    int a = 5;
    xobj.Set_a(a);
    xobj.Print_a();
}
```

In the member function Set_a(), the statement this->a = a uses the this pointer to retrieve xobj.a hidden by the automatic variable a.
Unless a class member name is hidden, using the class member name is equivalent to using the class member name with the this pointer and the class member access operator (->).

The example in the first column of the following table shows code that uses class members without the this pointer. The code in the second column uses the variable THIS to simulate the first column’s hidden use of the this pointer:
<table>
<thead>
<tr>
<th>Code without using this pointer</th>
<th>Equivalent code, the THIS variable simulating the hidden use of the this pointer</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>#include &lt;string&gt;</code></td>
<td><code>#include &lt;string&gt;</code></td>
</tr>
<tr>
<td><code>#include &lt;iostream&gt;</code></td>
<td><code>#include &lt;iostream&gt;</code></td>
</tr>
<tr>
<td>using namespace std;</td>
<td>using namespace std;</td>
</tr>
<tr>
<td>struct X {</td>
<td>struct X {</td>
</tr>
<tr>
<td>private:</td>
<td>private:</td>
</tr>
<tr>
<td>int len;</td>
<td>int len;</td>
</tr>
<tr>
<td>char *ptr;</td>
<td>char *ptr;</td>
</tr>
<tr>
<td>public:</td>
<td>public:</td>
</tr>
<tr>
<td>int GetLen() {</td>
<td>int GetLen (X* const THIS) {</td>
</tr>
<tr>
<td>return len;</td>
<td>return THIS-&gt;len;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>char * GetPtr() {</td>
<td>char * GetPtr (X* const THIS) {</td>
</tr>
<tr>
<td>return ptr;</td>
<td>return THIS-&gt;ptr;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>X&amp; Set(char *);</td>
<td>X&amp; Set(X* const, char *);</td>
</tr>
<tr>
<td>X&amp; Cat(char *);</td>
<td>X&amp; Cat(X* const, char *);</td>
</tr>
<tr>
<td>X&amp; Copy(X&amp;);</td>
<td>X&amp; Copy(X* const, X&amp;);</td>
</tr>
<tr>
<td>void Print();</td>
<td>void Print(X* const);</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>X&amp; X::Set(char *pc) {</td>
<td>X&amp; X::Set(X* const THIS, char *pc) {</td>
</tr>
<tr>
<td>len = strlen(pc);</td>
<td>THIS-&gt;len = strlen(pc);</td>
</tr>
<tr>
<td>ptr = new char[len];</td>
<td>THIS-&gt;ptr = new char[THIS-&gt;len];</td>
</tr>
<tr>
<td>strcpy(ptr, pc);</td>
<td>strcat(THIS-&gt;ptr, pc);</td>
</tr>
<tr>
<td>return *this;</td>
<td>return *THIS;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>X&amp; X::Cat(char *pc) {</td>
<td>X&amp; X::Cat(X* const THIS, char *pc) {</td>
</tr>
<tr>
<td>len += strlen(pc);</td>
<td>THIS-&gt;len += strlen(pc);</td>
</tr>
<tr>
<td>strcat(ptr, pc);</td>
<td>strcat(THIS-&gt;ptr, pc);</td>
</tr>
<tr>
<td>return *this;</td>
<td>return *THIS;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>X&amp; X::Copy(X&amp; x) {</td>
<td>X&amp; X::Copy(X* const THIS, X&amp; x) {</td>
</tr>
<tr>
<td>Set(x.GetPtr());</td>
<td>THIS-&gt;Set(THIS, x.GetPtr(&amp;x));</td>
</tr>
<tr>
<td>return *this;</td>
<td>return *THIS;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>void X::Print() {</td>
<td>void X::Print(X* const THIS) {</td>
</tr>
<tr>
<td>cout &lt;&lt; ptr &lt;&lt; endl;</td>
<td>cout &lt;&lt; THIS-&gt;ptr &lt;&lt; endl;</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
<tr>
<td>int main() {</td>
<td>int main() {</td>
</tr>
<tr>
<td>X xobj1;</td>
<td>X xobj1;</td>
</tr>
<tr>
<td>xobj1.Set(&quot;abcd&quot;)</td>
<td>xobj1.Set(&amp;xobj1, &quot;abcd&quot;)</td>
</tr>
<tr>
<td>.Cat(&quot;efgh&quot;);</td>
<td>.Cat(&amp;xobj1, &quot;efgh&quot;);</td>
</tr>
<tr>
<td>xobj1.Print();</td>
<td>xobj1.Print(&amp;xobj1);</td>
</tr>
<tr>
<td>X xobj2;</td>
<td>X xobj2;</td>
</tr>
<tr>
<td>xobj2.Copy(xobj1)</td>
<td>xobj2.Copy(&amp;xobj2, xobj1)</td>
</tr>
<tr>
<td>.Cat(&quot;ijkl&quot;);</td>
<td>.Cat(&amp;xobj2, &quot;ijkl&quot;);</td>
</tr>
<tr>
<td>xobj2.Print();</td>
<td>xobj2.Print(&amp;xobj2);</td>
</tr>
<tr>
<td>}</td>
<td>}</td>
</tr>
</tbody>
</table>

Both examples produce the following output:

```
abcdefgh
abcdefghijkl
```
Static members

Class members can be declared using the storage class specifier `static` in the class member list. Only one copy of the static member is shared by all objects of a class in a program. When you declare an object of a class having a static member, the static member is not part of the class object.

A typical use of static members is for recording data common to all objects of a class. For example, you can use a static data member as a counter to store the number of objects of a particular class type that are created. Each time a new object is created, this static data member can be incremented to keep track of the total number of objects.

You access a static member by qualifying the class name using the `::` (scope resolution) operator. In the following example, you can refer to the static member `f()` of class type `X` as `X::f()` even if no object of type `X` is ever declared:

```cpp
struct X {
    static int f();
};

int main() {
    X::f();
}
```

Using the class access operators with static members

You do not have to use the class member access syntax to refer to a static member; to access a static member `s` of class `X`, you could use the expression `X::s`. The following example demonstrates accessing a static member:

```cpp
#include <iostream>
using namespace std;

struct A {
    static void f() { cout << "In static function A::f()" << endl; }
};

int main() {
    // no object required for static member
    A::f();

    A a;
    A* ap = &a;
    a.f();
    ap->f();
}
```

The three statements `A::f()`, `a.f()`, and `ap->f()` all call the same static member function `A::f()`.
You can directly refer to a static member in the same scope of its class, or in the scope of a class derived from the static member’s class. The following example demonstrates the latter case (directly referring to a static member in the scope of a class derived from the static member’s class):

```cpp
#include <iostream>
using namespace std;

int g() {
    cout << "In function g()" << endl;
    return 0;
}

class X {
    public:
        static int g() {
            cout << "In static member function X::g()" << endl;
            return 1;
        }
    };

class Y: public X {
    public:
        static int i;
    };

int Y::i = g();

int main() {
}

The following is the output of the above code:

```
In static member function X::g()
```

The initialization `int Y::i = g()` calls `X::g()`, not the function `g()` declared in the global namespace.

Related information

- “The static storage class specifier” on page 44
- “Scope resolution operator :: (C++ only)” on page 120
- “Dot operator .” on page 123
- “Arrow operator ->” on page 123

**Static data members**

The declaration of a static data member in the member list of a class is not a definition. You must define the static member outside of the class declaration, in namespace scope. For example:

```cpp
class X {
    public:
        static int i;
    };

int X::i = 0; // definition outside class declaration
```

Once you define a static data member, it exists even though no objects of the static data member’s class exist. In the above example, no objects of class X exist even though the static data member X::i has been defined.

Static data members of a class in namespace scope have external linkage. The initializer for a static data member is in the scope of the class declaring the member.
A static data member can be of any type except for void or void qualified with const or volatile. You cannot declare a static data member as mutable.

You can only have one definition of a static member in a program. Unnamed classes, classes contained within unnamed classes, and local classes cannot have static data members.

Static data members and their initializers can access other static private and protected members of their class. The following example shows how you can initialize static members using other static members, even though these members are private:

```cpp
class C {
    static int i;
    static int j;
    static int k;
    static int l;
    static int m;
    static int n;
    static int p;
    static int q;
    static int s;
    static int f() { return 0; }
    int a;
public:
    C() { a = 0; }
};
C c;
int C::i = C::f();  // initialize with static member function
int C::j = C::i;   // initialize with another static data member
int C::k = c.f();  // initialize with member function from an object
int C::l = c.j;   // initialize with data member from an object
int C::s = c.a;   // initialize with nonstatic data member
int C::r = 1;    // initialize with a constant value

class Y : private C {} y;
int C::m = Y::f();
int C::n = Y::r;
int C::p = y.r;   // error
int C::q = y.f();  // error
```

The initializations of C::p and C::q cause errors because y is an object of a class that is derived privately from C, and its members are not accessible to members of C.

If a static data member is of const integral or const enumeration type, you may specify a constant initializer in the static data member’s declaration. This constant initializer must be an integral constant expression. Note that the constant initializer is not a definition. You still need to define the static member in an enclosing namespace. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct X {
    static const int a = 76;
};

const int X::a;
```
int main() {
    cout << X::a << endl;
}

The tokens = 76 at the end of the declaration of static data member a is a constant initializer.

Related information
• “External linkage” on page 9
• “Member access” on page 269
• “Local classes” on page 252

Static member functions
You cannot have static and nonstatic member functions with the same names and the same number and type of arguments.

Like static data members, you may access a static member function f() of a class A without using an object of class A.

A static member function does not have a this pointer. The following example demonstrates this:
#include <iostream>
using namespace std;

struct X {
private:
    int i;
    static int si;
public:
    void set_i(int arg) { i = arg; }
    static void set_si(int arg) { si = arg; }
    void print_i() {
        cout << "Value of i = " << i << endl;
        cout << "Again, value of i = " << this->i << endl;
    }
    static void print_si() {
        cout << "Value of si = " << si << endl;
        // cout << "Again, value of si = " << this->si << endl;
    }
};

int X::si = 77;       // Initialize static data member

int main() {
    X xobj;
    xobj.set_i(11);
    xobj.print_i();

    // static data members and functions belong to the class and
    // can be accessed without using an object of class X
    X::print_si();
    X::set_si(22);
    X::print_si();
}

The following is the output of the above example:
Value of \( i = 11 \)
Again, value of \( i = 11 \)
Value of \( s_i = 77 \)
Value of \( s_i = 22 \)

The compiler does not allow the member access operation \( \text{this-} > s_i \) in function \( A::\text{print}_s() \) because this member function has been declared as static, and therefore does not have a this pointer.

You can call a static member function using the this pointer of a nonstatic member function. In the following example, the nonstatic member function \( \text{printall}() \) calls the static member function \( f() \) using the this pointer:

```cpp
#include <iostream>
using namespace std;

class C {
    static void f() {
        cout << "Here is i: " << i << endl;
    }
    static int i;
    int j;
public:
    C(int firstj): j(firstj) { }
    void printall();
};

void C::printall() {
    cout << "Here is j: " << this->j << endl;
    this->f();
}

int C::i = 3;

int main() {
    C obj_C(0);
    obj_C.printall();
}
```

The following is the output of the above example:
Here is \( j: 0 \)
Here is \( i: 3 \)

A static member function cannot be declared with the keywords `virtual`, `const`, `volatile`, or `const volatile`.

A static member function can access only the names of static members, enumerators, and nested types of the class in which it is declared. Suppose a static member function \( f() \) is a member of class \( X \). The static member function \( f() \) cannot access the nonstatic members \( X \) or the nonstatic members of a base class of \( X \).

**Related information**
- ["The this pointer" on page 261](#)

**Member access**

*Member access* determines if a class member is accessible in an expression or declaration. Suppose \( x \) is a member of class \( A \). Class member \( x \) can be declared to have one of the following levels of accessibility:
public: x can be used anywhere without the access restrictions defined by private or protected.

private: x can be used only by the members and friends of class A.

protected: x can be used only by the members and friends of class A, and the members and friends of classes derived from class A.

Members of classes declared with the keyword class are private by default. Members of classes declared with the keyword struct or union are public by default.

To control the access of a class member, you use one of the access specifiers public, private, or protected as a label in a class member list. The following example demonstrates these access specifiers:

```cpp
struct A {
    friend class C;
private:
    int a;
public:
    int b;
protected:
    int c;
};

struct B : A {
    void f() {
        // a = 1;
        b = 2;
        c = 3;
    }
};

struct C {
    void f(A x) {
        x.a = 4;
        x.b = 5;
        x.c = 6;
    }
};

int main() {
    A y;
    // y.a = 7;
    y.b = 8;
    // y.c = 9;

    B z;
    // z.a = 10;
    z.b = 11;
    // z.c = 12;
}
```

The following table lists the access of data members A::a, A::b, and A::c in various scopes of the above example.

<table>
<thead>
<tr>
<th>Scope</th>
<th>A::a</th>
<th>A::b</th>
<th>A::c</th>
</tr>
</thead>
<tbody>
<tr>
<td>function B::f()</td>
<td>No access. Member A::a is private.</td>
<td>Access. Member A::b is public.</td>
<td>Access. Class B inherits from A.</td>
</tr>
<tr>
<td>function C::f()</td>
<td>Access. Class C is a friend of A.</td>
<td>Access. Member A::b is public.</td>
<td>Access. Class C is a friend of A.</td>
</tr>
<tr>
<td>Scope</td>
<td>A::a</td>
<td>A::b</td>
<td>A::c</td>
</tr>
<tr>
<td>-------</td>
<td>------</td>
<td>------</td>
<td>------</td>
</tr>
<tr>
<td>object y in main()</td>
<td>No access. Member y.a is private.</td>
<td>Access. Member y.a is public.</td>
<td>No access. Member y.c is protected.</td>
</tr>
<tr>
<td>object z in main()</td>
<td>No access. Member z.a is private.</td>
<td>Access. Member z.a is public.</td>
<td>No access. Member z.c is protected.</td>
</tr>
</tbody>
</table>

An access specifier specifies the accessibility of members that follow it until the next access specifier or until the end of the class definition. You can use any number of access specifiers in any order. If you later define a class member within its class definition, its access specification must be the same as its declaration. The following example demonstrates this:

```cpp
class A {
    class B;
    public:
        class B {}
};
```

The compiler will not allow the definition of class B because this class has already been declared as private.

A class member has the same access control regardless whether it has been defined within its class or outside its class.

Access control applies to names. In particular, if you add access control to a typedef name, it affects only the typedef name. The following example demonstrates this:

```cpp
class A {
    class B {}
    public:
        typedef B C;
};

int main() {
    A::C x;
    // A::B y;
}
```

The compiler will allow the declaration A::C x because the typedef name A::C is public. The compiler would not allow the declaration A::B y because A::B is private.

Note that accessibility and visibility are independent. Visibility is based on the scoping rules of C++. A class member can be visible and inaccessible at the same time.

**Related information**

- “Scope” on page 2
- “Class member lists” on page 255
- “Inherited member access” on page 282
Friends

A friend of a class X is a function or class that is not a member of X, but is granted the same access to X as the members of X. Functions declared with the friend specifier in a class member list are called friend functions of that class. Classes declared with the friend specifier in the member list of another class are called friend classes of that class.

A class Y must be defined before any member of Y can be declared a friend of another class.

In the following example, the friend function print is a member of class Y and accesses the private data members a and b of class X.

```cpp
#include <iostream>
using namespace std;

class X;

class Y {
public:
  void print(X& x);
};

class X {
  int a, b;
  friend void Y::print(X& x);
public:
  X() : a(1), b(2) { }
};

void Y::print(X& x) {
  cout << "a is " << x.a << endl;
  cout << "b is " << x.b << endl;
}

int main() {
  X xobj;
  Y yobj;
  yobj.print(xobj);
}
```

The following is the output of the above example:

```
a is 1
b is 2
```

You can declare an entire class as a friend. Suppose class F is a friend of class A. This means that every member function and static data member definition of class F has access to class A.

In the following example, the friend class F has a member function print that accesses the private data members a and b of class X and performs the same task as the friend function print in the above example. Any other members declared in class F also have access to all members of class X.

```cpp
#include <iostream>
using namespace std;

class X {
  int a, b;
  friend class F;
public:
  X() : a(1), b(2) { }
};
```
class F {
public:
    void print(X & x) {
        cout << "a is " << x.a << endl;
        cout << "b is " << x.b << endl;
    }
};

int main() {
    X xobj;
    F fobj;
    fobj.print(xobj);
}

The following is the output of the above example:

a is 1
b is 2

You must use an elaborated type specifier when you declare a class as a friend. The following example demonstrates this:

class F;
class G;
class X {
    friend class F;
    friend G;
};

The compiler will warn you that the friend declaration of G must be an elaborated class name.

You cannot define a class in a friend declaration. For example, the compiler will not allow the following:

class F;
class X {
    friend class F { }
};

However, you can define a function in a friend declaration. The class must be a non-local class, function, the function name must be unqualified, and the function has namespace scope. The following example demonstrates this:

class A {
    void g();
};

void z() {
    class B {
        // friend void f() { }
    };
}

class C {
    // friend void A::g() { }
    friend void h() { }
};

The compiler would not allow the function definition of f() or g(). The compiler will allow the definition of h().

You cannot declare a friend with a storage class specifier.
Friend scope

The name of a friend function or class first introduced in a friend declaration is not in the scope of the class granting friendship (also called the enclosing class) and is not a member of the class granting friendship.

The name of a function first introduced in a friend declaration is in the scope of the first nonclass scope that contains the enclosing class. The body of a function provided in a friend declaration is handled in the same way as a member function defined within a class. Processing of the definition does not start until the end of the outermost enclosing class. In addition, unqualified names in the body of the function definition are searched for starting from the class containing the function definition.

If the name of a friend class has been introduced before the friend declaration, the compiler searches for a class name that matches the name of the friend class beginning at the scope of the friend declaration. If the declaration of a nested class is followed by the declaration of a friend class with the same name, the nested class is a friend of the enclosing class.

The scope of a friend class name is the first nonclass enclosing scope. For example:

```cpp
class A {
    class B { // arbitrary nested class definitions
        friend class C;
    };
};
```

is equivalent to:

```cpp
class C;
class A {
    class B { // arbitrary nested class definitions
        friend class C;
    };
};
```

If the friend function is a member of another class, you need to use the scope resolution operator (::). For example:

```cpp
class A {
    public:
    int f() {} }
};
class B {
    friend int A::f();
};
```

Friends of a base class are not inherited by any classes derived from that base class. The following example demonstrates this:

```cpp
class A {
    friend class B;
    int a;
};
class B { }
```
class C : public B {
    void f(A* p) {
        // p->a = 2;
    }
};

The compiler would not allow the statement p->a = 2 because class C is not a friend of class A, although C inherits from a friend of A.

Friendship is not transitive. The following example demonstrates this:

```cpp
class A {
    friend class B;
    int a;
};
class B {
    friend class C;
};
class C {
    void f(A* p) {
        // p->a = 2;
    }
};
```

The compiler would not allow the statement p->a = 2 because class C is not a friend of class A, although C is a friend of a friend of A.

If you declare a friend in a local class, and the friend’s name is unqualified, the compiler will look for the name only within the innermost enclosing nonclass scope. You must declare a function before declaring it as a friend of a local scope. You do not have to do so with classes. However, a declaration of a friend class will hide a class in an enclosing scope with the same name. The following example demonstrates this:

```cpp
class X { }
void a();

void f() {
    class Y { }
    void b();
    class A {
        friend class X;
        friend class Y;
        friend class Z;
        // friend void a();
        friend void b();
        // friend void c();
    };
    ::X moocow;
    // X moocow2;
}
```

In the above example, the compiler will allow the following statements:

- `friend class X`: This statement does not declare `::X` as a friend of `A`, but the local class `X` as a friend, even though this class is not otherwise declared.
- `friend class Y`: Local class `Y` has been declared in the scope of `f()`.
- `friend class Z`: This statement declares the local class `Z` as a friend of `A` even though `Z` is not otherwise declared.
- `friend void b()`: Function `b()` has been declared in the scope of `f()`.
- `::X moocow`: This declaration creates an object of the nonlocal class `::X`. 

The compiler would not allow the following statements:

- `friend void a();`: This statement does not consider function `a()` declared in namespace scope. Since function `a()` has not been declared in the scope of `f()`, the compiler would not allow this statement.
- `friend void c();`: Since function `c()` has not been declared in the scope of `f()`, the compiler would not allow this statement.
- `X moo2`: This declaration tries to create an object of the local class `X`, not the nonlocal class `::X`. Since local class `X` has not been defined, the compiler would not allow this statement.

**Related information**

- [“Scope of class names” on page 249](#)
- [“Nested classes” on page 250](#)
- [“Local classes” on page 252](#)

**Friend access**

A friend of a class can access the private and protected members of that class. Normally, you can only access the private members of a class through member functions of that class, and you can only access the protected members of a class through member functions of a class or classes derived from that class.

Friend declarations are not affected by access specifiers.

**Related information**

- [“Member access” on page 269](#)
Chapter 13. Inheritance (C++ only)

Inheritance is a mechanism of reusing and extending existing classes without modifying them, thus producing hierarchical relationships between them.

Inheritance is almost like embedding an object into a class. Suppose that you declare an object x of class A in the class definition of B. As a result, class B will have access to all the public data members and member functions of class A. However, in class B, you have to access the data members and member functions of class A through object x. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

class A {
  int data;
public:
  void f(int arg) { data = arg; }
  int g() { return data; }
};
class B {
public:
  A x;
};

int main() {
  B obj;
  obj.x.f(20);
  cout << obj.x.g() << endl;
  // cout << obj.g() << endl;
}
```

In the main function, object obj accesses function A::f() through its data member B::x with the statement obj.x.f(20). Object obj accesses A::g() in a similar manner with the statement obj.x.g(). The compiler would not allow the statement obj.g() because g() is a member function of class A, not class B.

The inheritance mechanism lets you use a statement like obj.g() in the above example. In order for that statement to be legal, g() must be a member function of class B.

Inheritance lets you include the names and definitions of another class’s members as part of a new class. The class whose members you want to include in your new class is called a base class. Your new class is derived from the base class. The new class contains a subobject of the type of the base class. The following example is the same as the previous example except it uses the inheritance mechanism to give class B access to the members of class A:

```cpp
#include <iostream>
using namespace std;

class A {
  int data;
public:
  void f(int arg) { data = arg; }
  int g() { return data; }
};
```
class B : public A { }

int main() {
    B obj;
    obj.f(20);
    cout << obj.g() << endl;
}

Class A is a base class of class B. The names and definitions of the members of class A are included in the definition of class B; class B inherits the members of class A. Class B is derived from class A. Class B contains a subobject of type A.

You can also add new data members and member functions to the derived class. You can modify the implementation of existing member functions or data by overriding base class member functions or data in the newly derived class.

You may derive classes from other derived classes, thereby creating another level of inheritance. The following example demonstrates this:

struct A { }
struct B : A { }
struct C : B { }

Class B is a derived class of A, but is also a base class of C. The number of levels of inheritance is only limited by resources.

*Multiple inheritance* allows you to create a derived class that inherits properties from more than one base class. Because a derived class inherits members from all its base classes, ambiguities can result. For example, if two base classes have a member with the same name, the derived class cannot implicitly differentiate between the two members. Note that, when you are using multiple inheritance, the access to names of base classes may be ambiguous. See "Multiple inheritance" on page 288 for more detailed information.

A direct base class is a base class that appears directly as a base specifier in the declaration of its derived class.

An indirect base class is a base class that does not appear directly in the declaration of the derived class but is available to the derived class through one of its base classes. For a given class, all base classes that are not direct base classes are indirect base classes. The following example demonstrates direct and indirect base classes:

class A {  
    public:
        int x;
};
class B : public A {  
    public:
        int y;
};
class C : public B { }

Class B is a direct base class of C. Class A is a direct base class of B. Class A is an indirect base class of C. (Class C has x and y as its data members.)

Polymorphic functions are functions that can be applied to objects of more than one type. In C++, polymorphic functions are implemented in two ways:
- Overloaded functions are statically bound at compile time.
C++ provides virtual functions. A *virtual function* is a function that can be called for a number of different user-defined types that are related through derivation. Virtual functions are bound dynamically at run time. They are described in more detail in “Virtual functions” on page 295.

**Derivation**

Inheritance is implemented in C++ through the mechanism of derivation. Derivation allows you to derive a class, called a *derived class*, from another class, called a *base class*.

**Derived class syntax**

```
derived_class::
```

In the declaration of a derived class, you list the base classes of the derived class. The derived class inherits its members from these base classes.

The *qualified_class_specifier* must be a class that has been previously declared in a class declaration.

An *access specifier* is one of public, private, or protected.

The virtual keyword can be used to declare virtual base classes.

The following example shows the declaration of the derived class D and the base classes V, B1, and B2. The class B1 is both a base class and a derived class because it is derived from class V and is a base class for D:

```c++
class V { /* ... */ };
class B1 : virtual public V { /* ... */ };
class B2 { /* ... */ };
class D : public B1, private B2 { /* ... */ };
```

Classes that are declared but not defined are not allowed in base lists.

For example:

```c++
class X;
// error
class Y: public X { };
```

The compiler will not allow the declaration of class Y because X has not been defined.
When you derive a class, the derived class inherits class members of the base class. You can refer to inherited members (base class members) as if they were members of the derived class. For example:

class Base {
   public:
      int a, b;
};

class Derived : public Base {
   public:
      int c;
};

int main() {
   Derived d;
   d.a = 1;  // Base::a
   d.b = 2;  // Base::b
   d.c = 3;  // Derived::c
}

The derived class can also add new class members and redefine existing base class members. In the above example, the two inherited members, a and b, of the derived class d, in addition to the derived class member c, are assigned values. If you redefine base class members in the derived class, you can still refer to the base class members by using the :: (scope resolution) operator. For example:

#include <iostream>
using namespace std;

class Base {
   public:
      char* name;
      void display() { 
         cout << name << endl;
      }
};

class Derived: public Base {
   public:
      char* name;
      void display() { 
         cout << name << ", " << Base::name << endl;
      }
};

int main() {
   Derived d;
   d.name = "Derived Class";
   d.Base::name = "Base Class";

   // call Derived::display()
   d.display();

   // call Base::display()
   d.Base::display();
}

The following is the output of the above example:
Derived Class, Base Class
Base Class

You can manipulate a derived class object as if it were a base class object. You can use a pointer or a reference to a derived class object in place of a pointer or reference to its base class. For example, you can pass a pointer or reference to a
derived class object \( D \) to a function expecting a pointer or reference to the base class of \( D \). You do not need to use an explicit cast to achieve this; a standard conversion is performed. You can implicitly convert a pointer to a derived class to point to an accessible unambiguous base class. You can also implicitely convert a reference to a derived class to a reference to a base class.

The following example demonstrates a standard conversion from a pointer to a derived class to a pointer to a base class:

```cpp
#include <iostream>
using namespace std;

class Base {
public:
    char* name;
    void display() {
        cout << name << endl;
    }
};

class Derived: public Base {
public:
    char* name;
    void display() {
        cout << name << ", " << Base::name << endl;
    }
};

int main() {
    Derived d;
    d.name = "Derived Class";
    d.Base::name = "Base Class";
    Derived* dptr = &d;

    // standard conversion from Derived* to Base*
    Base* bptr = dptr;

    // call Base::display()
    bptr->display();
}
```

The following is the output of the above example:

```
Base Class
```

The statement `Base* bptr = dptr` converts a pointer of type `Derived` to a pointer of type `Base`.

The reverse case is not allowed. You cannot implicitly convert a pointer or a reference to a base class object to a pointer or reference to a derived class. For example, the compiler will not allow the following code if the classes `Base` and `Class` are defined as in the above example:

```cpp
int main() {
    Base b;
    b.name = "Base class";

    Derived* dptr = &b;
}
```

The compiler will not allow the statement `Derived* dptr = &b` because the statement is trying to implicitly convert a pointer of type `Base` to a pointer of type `Derived`.

Chapter 13. Inheritance (C++ only) 281
If a member of a derived class and a member of a base class have the same name, the base class member is hidden in the derived class. If a member of a derived class has the same name as a base class, the base class name is hidden in the derived class.

**Related information**
- “Virtual base classes” on page 289
- “Incomplete class declarations” on page 250
- “Scope resolution operator :: (C++ only)” on page 120

**Inherited member access**

The following sections discuss the access rules affecting a protected nonstatic base class member and how to declare a derived class using an access specifier:

- “Protected members”
- “Access control of base class members” on page 283

**Related information**
- “Member access” on page 269

**Protected members**

A protected nonstatic base class member can be accessed by members and friends of any classes derived from that base class by using one of the following:

- A pointer to a directly or indirectly derived class
- A reference to a directly or indirectly derived class
- An object of a directly or indirectly derived class

If a class is derived privately from a base class, all protected base class members become private members of the derived class.

If you reference a protected nonstatic member x of a base class A in a friend or a member function of a derived class B, you must access x through a pointer to, reference to, or object of a class derived from A. However, if you are accessing x to create a pointer to member, you must qualify x with a nested name specifier that names the derived class B. The following example demonstrates this:

```c++
class A {
    public:
    protected:
        int i;
};

class B : public A {
    friend void f(A*, B*);
    void g(A*);
};

void f(A* pa, B* pb) {
    // pa->i = 1;
    pb->i = 2;
    // int A::* point_i = &A::i;
    int A::* point_i2 = &B::i;
}

void B::g(A* pa) {
    // pa->i = 1;
```
i = 2;

// int A::* point_i = &A::i;
int A::* point_i2 = &B::i;
}

void h(A* pa, B* pb) {
// pa->i = 1;
// pb->i = 2;
}

int main() { }

Class A contains one protected data member, an integer i. Because B derives from A, the members of B have access to the protected member of A. Function f() is a friend of class B:
• The compiler would not allow pa->i = 1 because pa is not a pointer to the derived class B.
• The compiler would not allow int A::* point_i = &A::i because i has not been qualified with the name of the derived class B.

Function g() is a member function of class B. The previous list of remarks about which statements the compiler would and would not allow apply for g() except for the following:
• The compiler allows i = 2 because it is equivalent to this->i = 2.

Function h() cannot access any of the protected members of A because h() is neither a friend or a member of a derived class of A.

Access control of base class members

When you declare a derived class, an access specifier can precede each base class in the base list of the derived class. This does not alter the access attributes of the individual members of a base class as seen by the base class, but allows the derived class to restrict the access control of the members of a base class.

You can derive classes using any of the three access specifiers:
• In a public base class, public and protected members of the base class remain public and protected members of the derived class.
• In a protected base class, public and protected members of the base class are protected members of the derived class.
• In a private base class, public and protected members of the base class become private members of the derived class.

In all cases, private members of the base class remain private. Private members of the base class cannot be used by the derived class unless friend declarations within the base class explicitly grant access to them.

In the following example, class d is derived publicly from class b. Class b is declared a public base class by this declaration.

class b { }
class d : public b // public derivation
{ }

You can use both a structure and a class as base classes in the base list of a derived class declaration:
• If the derived class is declared with the keyword class, the default access specifier in its base list specifiers is private.
• If the derived class is declared with the keyword `struct`, the default access specifier in its base list specifiers is `public`.

In the following example, private derivation is used by default because no access specifier is used in the base list and the derived class is declared with the keyword `class`:

```cpp
struct B
{
};
class D : B // private derivation
{
};
```

Members and friends of a class can implicitly convert a pointer to an object of that class to a pointer to either:
• A direct private base class
• A protected base class (either direct or indirect)

**Related information**

- “Member access” on page 269
- “Member scope” on page 259

## The using declaration and class members

A using declaration in a definition of a class `A` allows you to introduce a *name* of a data member or member function from a base class of `A` into the scope of `A`.

You would need a using declaration in a class definition if you want to create a set of overload a member functions from base and derived classes, or you want to change the access of a class member.

**using declaration syntax**

```cpp
using [typename] nested_name_specifier::unqualified_id;
```

A using declaration in a class `A` may name one of the following:
• A member of a base class of `A`
• A member of an anonymous union that is a member of a base class of `A`
• An enumerator for an enumeration type that is a member of a base class of `A`

The following example demonstrates this:

```cpp
struct Z {
    int g();
};

struct A {
    void f();
    enum E { e };
    union { int u; };
};

struct B : A {
    using A::f;
    using A::e;
    using A::u;
    // using Z::g;
};
```
The compiler would not allow the using declaration using Z::g because Z is not a base class of A.

A using declaration cannot name a template. For example, the compiler will not allow the following:

```cpp
struct A {
    template<class T> void f(T);
};

struct B : A {
    using A::f<int>;
};
```

Every instance of the name mentioned in a using declaration must be accessible. The following example demonstrates this:

```cpp
struct A {
    private:
        void f(int);
    public:
        int f();
    protected:
        void g();
};

struct B : A {
    // using A::f;
    using A::g;
};
```

The compiler would not allow the using declaration using A::f because void A::f(int) is not accessible from B even though int A::f() is accessible.

Related information
- "Scope of class names" on page 249
- "The using declaration and namespaces" on page 226

Overloading member functions from base and derived classes

A member function named f in a class A will hide all other members named f in the base classes of A, regardless of return types or arguments. The following example demonstrates this:

```cpp
struct A {
    void f() { }  
};

struct B : A {
    void f(int) { }  
};

int main() {
    B obj_B;
    obj_B.f(3);
    // obj_B.f();
}
```

The compiler would not allow the function call obj_B.f() because the declaration of void B::f(int) has hidden A::f().
To overload, rather than hide, a function of a base class A in a derived class B, you introduce the name of the function into the scope of B with a using declaration. The following example is the same as the previous example except for the using declaration using A::f:

```cpp
struct A {
    void f() { }
};

struct B : A {
    using A::f;
    void f(int) { }
};

int main() {
    B obj_B;
    obj_B.f(3);
    obj_B.f();
}
```

Because of the using declaration in class B, the name f is overloaded with two functions. The compiler will now allow the function call obj_B.f().

You can overload virtual functions in the same way. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "void A::f()" << endl; }
    virtual void f(int) { cout << "void A::f(int)" << endl; }
};

struct B : A {
    using A::f;
    void f(int) { cout << "void B::f(int)" << endl; }
};

int main() {
    B obj_B;
    B* pb = &obj_B;
    pb->f(3);
    pb->f();
}
```

The following is the output of the above example:

```
void B::f(int)
void A::f()
```

Suppose that you introduce a function f from a base class A a derived class B with a using declaration, and there exists a function named B::f that has the same parameter types as A::f. Function B::f will hide, rather than conflict with, function A::f. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    void f() { }
    void f(int) { cout << "void A::f(int)" << endl; }
};

struct B : A {
    using A::f;
};
```
```cpp
void f(int) { cout << "void B::f(int)" << endl; }

int main() {
    B obj_B;
    obj_B.f(3);
}
```

The following is the output of the above example:

```cpp
void B::f(int)
```

**Related information**

- Chapter 10, “Overloading (C++ only)” on page 229
- “Name hiding (C++ only)” on page 7
- “The using declaration and class members” on page 284

## Changing the access of a class member

Suppose class B is a direct base class of class A. To restrict access of class B to the members of class A, derive B from A using either the access specifiers protected or private.

To increase the access of a member x of class A inherited from class B, use a using declaration. You cannot restrict the access to x with a using declaration. You may increase the access of the following members:

- A member inherited as private. (You cannot increase the access of a member declared as private because a using declaration must have access to the member’s name.)

The following example demonstrates this:

```cpp
struct A {
    protected:
    int y;
    public:
    int z;
};

struct B : private A { };

struct C : private A {
    public:
    using A::y;
    using A::z;
};

struct D : private A {
    protected:
    using A::y;
    using A::z;
};

struct E : D {
    void f() {
        y = 1;
        z = 2;
    }
};

struct F : A {
    public:
    using A::y;
```
private:
  using A::z;
};

int main()
{
  B obj_B;
  // obj_B.y = 3;
  // obj_B.z = 4;
  C obj_C;
  obj_C.y = 5;
  obj_C.z = 6;
  D obj_D;
  // obj_D.y = 7;
  // obj_D.z = 8;
  F obj_F;
  obj_F.y = 9;
  obj_F.z = 10;
}

The compiler would not allow the following assignments from the above example:
- obj_B.y = 3 and obj_B.z = 4: Members y and z have been inherited as private.
- obj_D.y = 7 and obj_D.z = 8: Members y and z have been inherited as private, but their access have been changed to protected.

The compiler allows the following statements from the above example:
- y = 1 and z = 2 in D::f(): Members y and z have been inherited as private, but their access have been changed to protected.
- obj_C.y = 5 and obj_C.z = 6: Members y and z have been inherited as private, but their access have been changed to public.
- obj_F.y = 9: The access of member y has been changed from protected to public.
- obj_F.z = 10: The access of member z is still public. The private using declaration using A::z has no effect on the access of z.

Related information
- “Member access” on page 269
- “Inherited member access” on page 282

Multiple inheritance

You can derive a class from any number of base classes. Deriving a class from more than one direct base class is called multiple inheritance.

In the following example, classes A, B, and C are direct base classes for the derived class X:

```c
class A { /* ... */ }
class B { /* ... */ }
class C { /* ... */ }
class X : public A, private B, public C { /* ... */ }
```

The following inheritance graph describes the inheritance relationships of the above example. An arrow points to the direct base class of the class at the tail of the arrow:
The order of derivation is relevant only to determine the order of default initialization by constructors and cleanup by destructors.

A direct base class cannot appear in the base list of a derived class more than once:

```cpp
class B1 { /* ... */ }; // direct base class
class D : public B1, private B1 { /* ... */ }; // error
```

However, a derived class can inherit an indirect base class more than once, as shown in the following example:

```cpp
class L { /* ... */ }; // indirect base class
class B2 : public L { /* ... */ };
class B3 : public L { /* ... */ };
class D : public B2, public B3 { /* ... */ }; // valid
```

In the above example, class D inherits the indirect base class L once through class B2 and once through class B3. However, this may lead to ambiguities because two subobjects of class L exist, and both are accessible through class D. You can avoid this ambiguity by referring to class L using a qualified class name. For example:

```
D::L
```

or

```
B3::L.
```

You can also avoid this ambiguity by using the base specifier `virtual` to declare a base class, as described in "Derivation" on page 279.

**Virtual base classes**

Suppose you have two derived classes B and C that have a common base class A, and you also have another class D that inherits from B and C. You can declare the base class A as `virtual` to ensure that B and C share the same subobject of A.

In the following example, an object of class D has two distinct subobjects of class L, one through class B1 and another through class B2. You can use the keyword `virtual` in front of the base class specifiers in the `base lists` of classes B1 and B2 to indicate that only one subobject of type L, shared by class B1 and class B2, exists.
For example:

```cpp
class L { /* ... */ }; // indirect base class
class B1 : virtual public L { /* ... */ };
class B2 : virtual public L { /* ... */ };
class D : public B1, public B2 { /* ... */ }; // valid
```

Using the keyword `virtual` in this example ensures that an object of class D inherits only one subobject of class L.

A derived class can have both virtual and nonvirtual base classes. For example:

```cpp
class V { /* ... */ };
class B1 : virtual public V { /* ... */ };
class B2 : virtual public V { /* ... */ };
class B3 : public V { /* ... */ };
class X : public B1, public B2, public B3 { /* ... */ };
```

In the above example, class X has two subobjects of class V, one that is shared by classes B1 and B2 and one through class B3.

**Related information**

- "Derivation" on page 279

**Multiple access**

In an inheritance graph containing virtual base classes, a name that can be reached through more than one path is accessed through the path that gives the most access.

For example:

```cpp
class L {
public:
    void f();
};
```
class B1 : private virtual L { }

class B2 : public virtual L { }

class D : public B1, public B2 {
public:
  void f() {
    // L::f() is accessed through B2
    // and is public
    L::f();
  }
};

In the above example, the function f() is accessed through class B2. Because class B2 is inherited publicly and class B1 is inherited privately, class B2 offers more access.

Related information
- “Member access” on page 269
- “Protected members” on page 282
- “Access control of base class members” on page 283

Ambiguous base classes

When you derive classes, ambiguities can result if base and derived classes have members with the same names. Access to a base class member is ambiguous if you use a name or qualified name that does not refer to a unique function or object. The declaration of a member with an ambiguous name in a derived class is not an error. The ambiguity is only flagged as an error if you use the ambiguous member name.

For example, suppose that two classes named A and B both have a member named x, and a class named C inherits from both A and B. An attempt to access x from class C would be ambiguous. You can resolve ambiguity by qualifying a member with its class name using the scope resolution (::) operator.

```cpp
class B1 {
public:
  int i;
  int j;
  void g(int) { }
};

class B2 {
public:
  int j;
  void g() { }
};

class D : public B1, public B2 {
public:
  int i;
};

int main() {
  D dobj;
  D *dptr = &dobj;
  dptr->i = 5;
  // dptr->j = 10;
  dptr->B1::j = 10;
  // dobj.g();
  dobj.B2::g();
}
The statement `dptr->j = 10` is ambiguous because the name `j` appears both in `B1` and `B2`. The statement `dobj.g()` is ambiguous because the name `g` appears both in `B1` and `B2`, even though `B1::g(int)` and `B2::g()` have different parameters.

The compiler checks for ambiguities at compile time. Because ambiguity checking occurs before access control or type checking, ambiguities may result even if only one of several members with the same name is accessible from the derived class.

**Name hiding**
Suppose two subobjects named `A` and `B` both have a member name `x`. The member name `x` of subobject `B` hides the member name `x` of subobject `A` if `A` is a base class of `B`. The following example demonstrates this:

```c
struct A { int x; };
struct B: A { int x; };
struct C: A, B { void f() { x = 0; } };
int main() {
  C i;
  i.f();
}
```

The assignment `x = 0` in function `C::f()` is not ambiguous because the declaration `B::x` has hidden `A::x`. However, the compiler will warn you that deriving `C` from `A` is redundant because you already have access to the subobject `A` through `B`.

A base class declaration can be hidden along one path in the inheritance graph and not hidden along another path. The following example demonstrates this:

```c
struct A { int x; };
struct B { int y; };
struct C: A, virtual B { };
struct D: A, virtual B {
  int x;
  int y;
};
struct E: C, D { };
int main() {
  E e;
  // e.x = 1;
  e.y = 2;
}
```

The assignment `e.x = 1` is ambiguous. The declaration `D::x` hides `A::x` along the path `D::A::x`, but it does not hide `A::x` along the path `C::A::x`. Therefore the variable `x` could refer to either `D::x` or `A::x`. The assignment `e.y = 2` is not ambiguous. The declaration `D::y` hides `B::y` along both paths `D::B::y` and `C::B::y` because `B` is a virtual base class.

**Ambiguity and using declarations**
Suppose you have a class named `C` that inherits from a class named `A`, and `x` is a member name of `A`. If you use a using declaration to declare `A::x` in `C`, then `x` is
also a member of C; C::x does not hide A::x. Therefore using declarations cannot resolve ambiguities due to inherited members. The following example demonstrates this:

```cpp
struct A {
    int x;
};
struct B: A {};
struct C: A {
    using A::x;
};
struct D: B, C {
    void f() { x = 0; }
};
int main() {
    D i;
    i.f();
}
```

The compiler will not allow the assignment `x = 0` in function `D::f()` because it is ambiguous. The compiler can find `x` in two ways: as `B::x` or as `C::x`.

**Unambiguous class members**

The compiler can unambiguously find static members, nested types, and enumerators defined in a base class `A` regardless of the number of subobjects of type `A` an object has. The following example demonstrates this:

```cpp
struct A {
    int x;
    static int s;
    typedef A* Pointer_A;
    enum { e };
};
int A::s;
struct B: A {};
struct C: A {};
struct D: B, C {
    void f() {
        s = 1;
        Pointer_A pa;
        int i = e;
        // x = 1;
    }
};
int main() {
    D i;
    i.f();
}
```

The compiler allows the assignment `s = 1`, the declaration `Pointer_A pa`, and the statement `int i = e`. There is only one static variable `s`, only one typedef `Pointer_A`, and only one enumerator `e`. The compiler would not allow the assignment `x = 1` because `x` can be reached either from class `B` or class `C`. 
**Pointer conversions**

Conversions (either implicit or explicit) from a derived class pointer or reference to a base class pointer or reference must refer unambiguously to the same accessible base class object. (An accessible base class is a publicly derived base class that is neither hidden nor ambiguous in the inheritance hierarchy.) For example:

```c++
class W { /* ... */ };  
class X : public W { /* ... */ };  
class Y : public W { /* ... */ };  
class Z : public X, public Y { /* ... */ };  
int main ()
{
    Z z;
    X* xptr = &z;    // valid
    Y* yptr = &z;    // valid
    W* wptr = &z;    // error, ambiguous reference to class W
                     // X's W or Y's W?
}
```

You can use virtual base classes to avoid ambiguous reference. For example:

```c++
class W { /* ... */ };  
class X : public virtual W { /* ... */ };  
class Y : public virtual W { /* ... */ };  
class Z : public X, public Y { /* ... */ };  
int main ()
{
    Z z;
    X* xptr = &z;    // valid
    Y* yptr = &z;    // valid
    W* wptr = &z;    // valid, W is virtual therefore only one
                     // W subobject exists
}
```

A pointer to a member of a base class can be converted to a pointer to a member of a derived class if the following conditions are true:

- The conversion is not ambiguous. The conversion is ambiguous if multiple instances of the base class are in the derived class.
- A pointer to the derived class can be converted to a pointer to the base class. If this is the case, the base class is said to be accessible.
- Member types must match. For example suppose class A is a base class of class B. You cannot convert a pointer to member of A of type int to a pointer to member of type B of type float.
- The base class cannot be virtual.

**Overload resolution**

Overload resolution takes place after the compiler unambiguously finds a given function name. The following example demonstrates this:

```c++
struct A {
    int f() { return 1; }
};

struct B {
    int f(int arg) { return arg; }
};

struct C: A, B {
    int g() { return f(); }
};
```
The compiler will not allow the function call to f() in C::g() because the name f has been declared both in A and B. The compiler detects the ambiguity error before overload resolution can select the base match A::f().

Related information
- “Scope resolution operator :: (C++ only)” on page 120
- “Virtual base classes” on page 289

Virtual functions

By default, C++ matches a function call with the correct function definition at compile time. This is called static binding. You can specify that the compiler match a function call with the correct function definition at run time; this is called dynamic binding. You declare a function with the keyword virtual if you want the compiler to use dynamic binding for that specific function.

The following examples demonstrate the differences between static and dynamic binding. The first example demonstrates static binding:

```cpp
#include <iostream>
using namespace std;

struct A {
    void f() { cout << "Class A" << endl; }
};

struct B: A {
    void f() { cout << "Class B" << endl; }
};

void g(A& arg) {
    arg.f();
}

int main() {
    B x;
    g(x);
}
```

The following is the output of the above example:

Class A

When function g() is called, function A::f() is called, although the argument refers to an object of type B. At compile time, the compiler knows only that the argument of function g() will be a reference to an object derived from A; it cannot determine whether the argument will be a reference to an object of type A or type B. However, this can be determined at run time. The following example is the same as the previous example, except that A::f() is declared with the virtual keyword:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "Class A" << endl; }
};

struct B: A {
    void f() { cout << "Class B" << endl; }
};

void g(A& arg) {
    arg.f();
}
```

Chapter 13. Inheritance (C++ only) 295
```cpp
int main() {
    B x;
    y(x);
}
```

The following is the output of the above example:

Class B

The virtual keyword indicates to the compiler that it should choose the appropriate definition of f() not by the type of reference, but by the type of object that the reference refers to.

Therefore, a virtual function is a member function you may redefine for other derived classes, and can ensure that the compiler will call the redefined virtual function for an object of the corresponding derived class, even if you call that function with a pointer or reference to a base class of the object.

A class that declares or inherits a virtual function is called a polymorphic class.

You redefine a virtual member function, like any member function, in any derived class. Suppose you declare a virtual function named f in a class A, and you derive directly or indirectly from A a class named B. If you declare a function named f in class B with the same name and same parameter list as A::f, then B::f is also virtual (regardless whether or not you declare B::f with the virtual keyword) and it overrides A::f. However, if the parameter lists of A::f and B::f are different, A::f and B::f are considered different, B::f does not override A::f, and B::f is not virtual (unless you have declared it with the virtual keyword). Instead B::f hides A::f. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "Class A" << endl; }
};

struct B : A {
    void f(int) { cout << "Class B" << endl; }
};

struct C : B {
    void f() { cout << "Class C" << endl; }
};

int main() {
    B b; C c;
    A* pa1 = &b;
    A* pa2 = &c;
    // b.f();
    pa1->f();
    pa2->f();
}
```

The following is the output of the above example:

Class A
Class C
The function B::f is not virtual. It hides A::f. Thus the compiler will not allow the function call b.f(). The function C::f is virtual; it overrides A::f even though A::f is not visible in C.

If you declare a base class destructor as virtual, a derived class destructor will override that base class destructor, even though destructors are not inherited.

The return type of an overriding virtual function may differ from the return type of the overridden virtual function. This overriding function would then be called a covariant virtual function. Suppose that B::f overrides the virtual function A::f. The return types of A::f and B::f may differ if all the following conditions are met:

- The function B::f returns a reference or pointer to a class of type T, and A::f returns a pointer or a reference to an unambiguous direct or indirect base class of T.
- The const or volatile qualification of the pointer or reference returned by B::f has the same or less const or volatile qualification of the pointer or reference returned by A::f.
- The return type of B::f must be complete at the point of declaration of B::f, or it can be of type B.

The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {}

class B : private A {
    friend class D;
    friend class F;
};

A global_A;
B global_B;

struct C {
    virtual A* f() {
        cout << "A* C::f()" << endl;
        return &global_A;
    }
};

struct D : C {
    B* f() {
        cout << "B* D::f()" << endl;
        return &global_B;
    }
};

struct E;

struct F : C {
    // Error:
    // E is incomplete
    // E* f();
};

struct G : C {
    // Error:
    // A is an inaccessible base class of B
    // B* f();
};
```
int main() {
    D d;
    C* cp = &d;
    D* dp = &d;
    A* ap = cp->f();
    B* bp = dp->f();
}

The following is the output of the above example:
B* D::f()
B* D::f()

The statement A* ap = cp->f() calls D::f() and converts the pointer returned to type A*. The statement B* bp = dp->f() calls D::f() as well but does not convert the pointer returned; the type returned is B*. The compiler would not allow the declaration of the virtual function F::f() because E is not a complete class. The compiler would not allow the declaration of the virtual function G::f() because class A is not an accessible base class of B (unlike friend classes D and F, the definition of B does not give access to its members for class G).

A virtual function cannot be global or static because, by definition, a virtual function is a member function of a base class and relies on a specific object to determine which implementation of the function is called. You can declare a virtual function to be a friend of another class.

If a function is declared virtual in its base class, you can still access it directly using the scope resolution (::) operator. In this case, the virtual function call mechanism is suppressed and the function implementation defined in the base class is used. In addition, if you do not override a virtual member function in a derived class, a call to that function uses the function implementation defined in the base class.

A virtual function must be one of the following:
• Defined
• Declared pure
• Defined and declared pure

A base class containing one or more pure virtual member functions is called an abstract class.

Related information
• “Abstract classes” on page 300

Ambiguous virtual function calls
You cannot override one virtual function with two or more ambiguous virtual functions. This can happen in a derived class that inherits from two nonvirtual bases that are derived from a virtual base class.

For example:
class V {
    public:
        virtual void f() { }
};
class A : virtual public V {
    void f() { }
};

class B : virtual public V {
    void f() { }
};

// Error:
// Both A::f() and B::f() try to override V::f()
class D : public A, public B {};

int main() {
    D d;
    V* vptr = &d;

    // which f(), A::f() or B::f()?
    vptr->f();
}

The compiler will not allow the definition of class D. In class A, only A::f() will override V::f(). Similarly, in class B, only B::f() will override V::f(). However, in class D, both A::f() and B::f() will try to override V::f(). This attempt is not allowed because it is not possible to decide which function to call if a D object is referenced with a pointer to class V, as shown in the above example. Only one function can override a virtual function.

A special case occurs when the ambiguous overriding virtual functions come from separate instances of the same class type. In the following example, class D has two separate subobjects of class A:
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "A::f()" << endl; }
};

struct B : A {
    void f() { cout << "B::f()" << endl; }
};

struct C : A {
    void f() { cout << "C::f()" << endl; }
};

struct D : B, C {};

int main() {
    D d;
    B* bp = &d;
    A* ap = bp;
    D* dp = &d;

    ap->f();
    // dp->f();
}

Class D has two occurrences of class A, one inherited from B, and another inherited from C. Therefore there are also two occurrences of the virtual function A::f. The statement ap->f() calls D::B::f. However the compiler would not allow the statement dp->f() because it could either call D::B::f or D::C::f.
Virtual function access

The access for a virtual function is specified when it is declared. The access rules for a virtual function are not affected by the access rules for the function that later overrides the virtual function. In general, the access of the overriding member function is not known.

If a virtual function is called with a pointer or reference to a class object, the type of the class object is not used to determine the access of the virtual function. Instead, the type of the pointer or reference to the class object is used.

In the following example, when the function \( f() \) is called using a pointer having type \( B* \), \( bptr \) is used to determine the access to the function \( f() \). Although the definition of \( f() \) defined in class \( D \) is executed, the access of the member function \( f() \) in class \( B \) is used. When the function \( f() \) is called using a pointer having type \( D* \), \( dptr \) is used to determine the access to the function \( f() \). This call produces an error because \( f() \) is declared private in class \( D \).

```cpp
class B {
    public:
    virtual void f();
};

class D : public B {
    private:
    void f();
};

int main() {
    D dobj;
    B* bptr = &dobj;
    D* dptr = &dobj;

    // valid, virtual B::f() is public,
    // D::f() is called
    bptr->f();

    // error, D::f() is private
    dptr->f();
}
```

Abstract classes

An abstract class is a class that is designed to be specifically used as a base class. An abstract class contains at least one pure virtual function. You declare a pure virtual function by using a pure specifier (= 0) in the declaration of a virtual member function in the class declaration.

The following is an example of an abstract class:

```cpp
class AB {
    public:
    virtual void f() = 0;
};
```

Function \( AB::f \) is a pure virtual function. A function declaration cannot have both a pure specifier and a definition. For example, the compiler will not allow the following:

```cpp
struct A {
    virtual void g() { } = 0;
};
```
You cannot use an abstract class as a parameter type, a function return type, or the
type of an explicit conversion, nor can you declare an object of an abstract class.
You can, however, declare pointers and references to an abstract class. The
following example demonstrates this:

```cpp
struct A {
    virtual void f() = 0;
};

struct B : A {
    virtual void f() {}  
};

// Error:
// Class A is an abstract class
// A g();

// Error:
// Class A is an abstract class
// void h(A);
A& i(A&);

int main() {
    // Error:
    // Class A is an abstract class
    // A a;

    A* pa;
    B b;

    // Error:
    // Class A is an abstract class
    // static_cast<A>(b);
}
```

Class A is an abstract class. The compiler would not allow the function declarations
A g() or void h(A), declaration of object a, nor the static cast of b to type A.

Virtual member functions are inherited. A class derived from an abstract base class
will also be abstract unless you override each pure virtual function in the derived
class.

For example:

```cpp
class AB {
public:
    virtual void f() = 0;
};

class D2 : public AB {
    void g();
};

int main() {
    D2 d;
}
```

The compiler will not allow the declaration of object d because D2 is an abstract
class; it inherited the pure virtual function f() from AB. The compiler will allow the
declaration of object d if you define function D2::g().

Note that you can derive an abstract class from a nonabstract class, and you can
override a non-pure virtual function with a pure virtual function.
You can call member functions from a constructor or destructor of an abstract class. However, the results of calling (directly or indirectly) a pure virtual function from its constructor are undefined. The following example demonstrates this:

```c
struct A {
    A() {
        direct();
        indirect();
    }
    virtual void direct() = 0;
    virtual void indirect() { direct(); }
};
```

The default constructor of A calls the pure virtual function `direct()` both directly and indirectly (through `indirect()`).

The compiler issues a warning for the direct call to the pure virtual function, but not for the indirect call.

**Related information**

- "Virtual functions" on page 295
- "Virtual function access" on page 300
Chapter 14. Special member functions (C++ only)

The default constructor, destructor, copy constructor, and copy assignment operator are *special member functions*. These functions create, destroy, convert, initialize, and copy class objects, and are discussed in the following sections:

- “Overview of constructors and destructors” on page 305
- “Constructors” on page 316
- “Destructors” on page 312
- “Conversion by constructor” on page 322
- “Conversion functions” on page 324
- “Copy constructors” on page 325

Overview of constructors and destructors

Because classes have complicated internal structures, including data and functions, object initialization and cleanup for classes is much more complicated than it is for simple data structures. Constructors and destructors are special member functions of classes that are used to construct and destroy class objects. Construction may involve memory allocation and initialization for objects. Destruction may involve cleanup and deallocation of memory for objects.

Like other member functions, constructors and destructors are declared within a class declaration. They can be defined inline or external to the class declaration. Constructors can have default arguments. Unlike other member functions, constructors can have member initialization lists. The following restrictions apply to constructors and destructors:

- Constructors and destructors do not have return types nor can they return values.
- References and pointers cannot be used on constructors and destructors because their addresses cannot be taken.
- Constructors cannot be declared with the keyword *virtual*.
- Constructors and destructors cannot be declared *static*, *const*, or *volatile*.
- Unions cannot contain class objects that have constructors or destructors.

Constructors and destructors obey the same access rules as member functions. For example, if you declare a constructor with protected access, only derived classes and friends can use it to create class objects.

The compiler automatically calls constructors when defining class objects and calls destructors when class objects go out of scope. A constructor does not allocate memory for the class object its *this* pointer refers to, but may allocate storage for more objects than its class object refers to. If memory allocation is required for objects, constructors can explicitly call the `new` operator. During cleanup, a destructor may release objects allocated by the corresponding constructor. To release objects, use the `delete` operator.

Derived classes do not inherit or overload constructors or destructors from their base classes, but they do call the constructor and destructor of base classes. Destructors can be declared with the keyword *virtual*. 

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303
Constructors are also called when local or temporary class objects are created, and destructors are called when local or temporary objects go out of scope.

You can call member functions from constructors or destructors. You can call a virtual function, either directly or indirectly, from a constructor or destructor of a class A. In this case, the function called is the one defined in A or a base class of A, but not a function overridden in any class derived from A. This avoids the possibility of accessing an unconstructed object from a constructor or destructor. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "void A::f()" << endl; }
    virtual void g() { cout << "void A::g()" << endl; }
    virtual void h() { cout << "void A::h()" << endl; }
};

struct B : A {
    virtual void f() { cout << "void B::f()" << endl; }
    B() { f(); g(); h(); }
};

struct C : B {
    virtual void f() { cout << "void C::f()" << endl; }
    virtual void g() { cout << "void C::g()" << endl; }
    virtual void h() { cout << "void C::h()" << endl; }
};

int main() {
    C obj;
}
```

The following is the output of the above example:

```
void B::f()
void A::g()
void A::h()
```

The constructor of B does not call any of the functions overridden in C because C has been derived from B, although the example creates an object of type C named obj.

You can use the typeid or the `dynamic_cast` operator in constructors or destructors, as well as member initializers of constructors.

**Related information**

- "The constructor and destructor function attributes" on page 209
- "The new operator (C++ only)" on page 140
- "The delete operator (C++ only)" on page 142
- "Free store" on page 315
Constructors

A constructor is a member function with the same name as its class. For example:

```cpp
class X {
public:
    X(); // constructor for class X
};
```

Constructors are used to create, and can initialize, objects of their class type.

You cannot declare a constructor as virtual or static, nor can you declare a constructor as const, volatile, or const volatile.

You do not specify a return type for a constructor. A return statement in the body of a constructor cannot have a return value.

Related information

- “The constructor and destructor function attributes” on page 209
- “Free store” on page 315

Default constructors

A default constructor is a constructor that either has no parameters, or if it has parameters, all the parameters have default values.

If no user-defined constructor exists for a class A and one is needed, the compiler implicitly declares a default parameterless constructor A::A(). This constructor is an inline public member of its class. The compiler will implicitly define A::A() when the compiler uses this constructor to create an object of type A. The constructor will have no constructor initializer and a null body.

The compiler first implicitly defines the implicitly declared constructors of the base classes and nonstatic data members of a class A before defining the implicitly declared constructor of A. No default constructor is created for a class that has any constant or reference type members.

A constructor of a class A is trivial if all the following are true:

- It is implicitly defined
- A has no virtual functions and no virtual base classes
- All the direct base classes of A have trivial constructors
- The classes of all the nonstatic data members of A have trivial constructors

If any of the above are false, then the constructor is nontrivial.

A union member cannot be of a class type that has a nontrivial constructor.

Like all functions, a constructor can have default arguments. They are used to initialize member objects. If default values are supplied, the trailing arguments can be omitted in the expression list of the constructor. Note that if a constructor has any arguments that do not have default values, it is not a default constructor.

A copy constructor for a class A is a constructor whose first parameter is of type A&, const A&, volatile A&, or const volatile A&. Copy constructors are used to make a copy of one class object from another class object of the same class type. You cannot use a copy constructor with an argument of the same type as its class; you
must use a reference. You can provide copy constructors with additional parameters as long as they all have default arguments. If a user-defined copy constructor does not exist for a class and one is needed, the compiler implicitly creates a copy constructor, with public access, for that class. A copy constructor is not created for a class if any of its members or base classes have an inaccessible copy constructor.

The following code fragment shows two classes with constructors, default constructors, and copy constructors:

```c++
class X {
public:
    // default constructor, no arguments
    X();
    // constructor
    X(int, int , int = 0);
    // copy constructor
    X(const X&);
    // error, incorrect argument type
    X(X);
};

class Y {
public:
    // default constructor with one
    // default argument
    Y( int = 0);
    // default argument
    // copy constructor
    Y(const Y&, int = 0);
};
```

Related information
- “The constructor and destructor function attributes” on page 209
- “Copy constructors” on page 325

Explicit initialization with constructors

A class object with a constructor must be explicitly initialized or have a default constructor. Except for aggregate initialization, explicit initialization using a constructor is the only way to initialize nonstatic constant and reference class members.

A class object that has no constructors, no virtual functions, no private or protected members, and no base classes is called an aggregate. Examples of aggregates are C-style structures and unions.

You explicitly initialize a class object when you create that object. There are two ways to initialize a class object:
- Using a parenthesized expression list. The compiler calls the constructor of the class using this list as the constructor’s argument list.
- Using a single initialization value and the = operator. Because this type of expression is an initialization, not an assignment, the assignment operator function, if one exists, is not called. The type of the single argument must match
the type of the first argument to the constructor. If the constructor has remaining arguments, these arguments must have default values.

**Initializer syntax**

```
expression (expression, {expression})
```

The following example shows the declaration and use of several constructors that explicitly initialize class objects:

```cpp
// This example illustrates explicit initialization
// by constructor.
#include <iostream>
using namespace std;

class complx {
   double re, im;
public:
   // default constructor
   complx() : re(0), im(0) {}
   // copy constructor
   complx(const complx& c) { re = c.re; im = c.im; }
   // constructor with default trailing argument
   complx(double r, double i = 0.0) { re = r; im = i; }
   void display() {
      cout << "re = " << re << " im = " << im << endl;
   }
};

int main() {
   // initialize with complx(double, double)
   complx one(1);
   // initialize with a copy of one
   // using complx::complx(const complx&
   complx two = one;
   // construct complx(3,4)
   // directly into three
   complx three = complx(3,4);
   // initialize with default constructor
   complx four;
   // complx(double, double) and construct
   // directly into five
   complx five = 5;

   one.display();
two.display();
three.display();
four.display();
five.display();
}
```
The above example produces the following output:

```c
re = 1  im = 0
re = 1  im = 0
re = 3  im = 4
re = 0  im = 0
re = 5  im = 0
```

Related information

- “The constructor and destructor function attributes” on page 209
- “Initializers” on page 89

### Initializing base classes and members

Constructors can initialize their members in two different ways. A constructor can use the arguments passed to it to initialize member variables in the constructor definition:

```c
complx(double r, double i = 0.0) { re = r; im = i; }
```

Or a constructor can have an *initializer list* within the definition but prior to the constructor body:

```c
complx(double r, double i = 0) : re(r), im(i) { /* ... */ }
```

Both methods assign the argument values to the appropriate data members of the class.

### Initializer list syntax

Include the initialization list as part of the constructor definition, not as part of the constructor declaration. For example:

```c
#include <iostream>
using namespace std;

class B1 {
    int b;
public:
    B1() { cout << "B1::B1()" << endl; };
    // inline constructor
    B1(int i) : b(i) { cout << "B1::B1(int)" << endl; }
};
class B2 {
    int b;
protected:
    B2() { cout << "B2::B2()" << endl; }
    // noninline constructor
    B2(int i);
};

// B2 constructor definition including initialization list
B2::B2(int i) : b(i) { cout << "B2::B2(int)" << endl; }

class D : public B1, public B2 {
```
```cpp
int d1, d2;
public:
  D(int i, int j) : B1(i+1), B2(), d1(i) {
    cout << "D1::D1(int, int)" << endl;
    d2 = j;
  }

int main() {
  D obj(1, 2);
}

The following is the output of the above example:
B1::B1(int)
B1::B1()
D1::D1(int, int)
```

If you do not explicitly initialize a base class or member that has constructors by calling a constructor, the compiler automatically initializes the base class or member with a default constructor. In the above example, if you leave out the call B2() in the constructor of class D (as shown below), a constructor initializer with an empty expression list is automatically created to initialize B2. The constructors for class D, shown above and below, result in the same construction of an object of class D:

```cpp
class D : public B1, public B2 {
  int d1, d2;

  // call B2() generated by compiler
  D(int i, int j) : B1(i+1), d1(i) {
    cout << "D1::D1(int, int)" << endl;
    d2 = j;
  }

  int main() {
    D obj(1, 2);
  }
```

In the above example, the compiler will automatically call the default constructor for B2().

Note that you must declare constructors as public or protected to enable a derived class to call them. For example:

```cpp
class B {
  B() { }
};

class D : public B {

  // error: implicit call to private B() not allowed
  D() { }
};
```

The compiler does not allow the definition of D::D() because this constructor cannot access the private constructor B::B().

You must initialize the following with an initializer list: base classes with no default constructors, reference data members, non-static const data members, or a class type which contains a constant data member. The following example demonstrates this:

```cpp
class A {
  public:
    A(int) { }
};

class B : public A {
```

Chapter 14. Special member functions (C++ only)  309
The data members \( j \) and \( k \), as well as the base class \( A \) must be initialized in the initializer list of the constructor of \( B \).

You can use data members when initializing members of a class. The following example demonstrate this:

```cpp
struct A {
  int k;
  A(int i) : k(i) {}  
};
struct B : A {
  int x;
  int i;
  int j;
  int& r;
  B(int i) : r(x), A(i), j(this->i), i(i) {}  
};
```

The constructor \( \text{B}(\text{i}) \) initializes the following:
- \( \text{B}::\text{r} \) to refer to \( \text{B}::\text{x} \)
- \( \text{Class A with the value of the argument to B(int i)} \)
- \( \text{B}::\text{j} \) with the value of \( \text{B}::\text{i} \)
- \( \text{B}::\text{i} \) with the value of the argument to \( \text{B}(\text{i}) \)

You can also call member functions (including virtual member functions) or use the operators \text{typeid} or \text{dynamic_cast} when initializing members of a class. However if you perform any of these operations in a member initialization list before all base classes have been initialized, the behavior is undefined. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;
struct A {
  int i;
  A(int arg) : i(arg) {
    cout << "Value of i: " << i << endl;
  }
};
struct B : A {
  int j;
  int f() { return i; }
  B();
};
B::B() : A(f()), j(1234) {
  cout << "Value of j: " << j << endl;
}
```
int main() {
    B obj;
}

The output of the above example would be similar to the following:
Value of i: 8
Value of j: 1234

The behavior of the initializer \texttt{A(f())} in the constructor of \texttt{B} is undefined. The runtime will call \texttt{B::f()} and try to access \texttt{A::i} even though the base \texttt{A} has not been initialized.

The following example is the same as the previous example except that the initializers of \texttt{B::B()} have different arguments:

```cpp
#include <iostream>
using namespace std;

struct A {
    int i;
    A(int arg) : i(arg) {
        cout << "Value of i: " << i << endl;
    }
};

struct B : A {
    int j;
    int f() { return i; }
    B();
};

B::B() : A(5678), j(f()) {
    cout << "Value of j: " << j << endl;
}

int main() {
    B obj;
}
```

The following is the output of the above example:
Value of i: 5678
Value of j: 5678

The behavior of the initializer \texttt{j(f())} in the constructor of \texttt{B} is well-defined. The base class \texttt{A} is already initialized when \texttt{B::j} is initialized.

Related information
- "The \texttt{typeid} operator (C++ only)" on page 125
- "The \texttt{dynamic_cast} operator (C++ only)" on page 129

**Construction order of derived class objects**

When a derived class object is created using constructors, it is created in the following order:
1. Virtual base classes are initialized, in the order they appear in the base list.
2. Nonvirtual base classes are initialized, in declaration order.
3. Class members are initialized in declaration order (regardless of their order in the initialization list).
4. The body of the constructor is executed.
The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;
struct V {
    V() { cout << "V()" << endl; }
};
struct V2 {
    V2() { cout << "V2()" << endl; }
};
struct A {
    A() { cout << "A()" << endl; }
};
struct B : virtual V {
    B() { cout << "B()" << endl; }
};
struct C : B, virtual V2 {
    C() { cout << "C()" << endl; }
};
struct D : C, virtual V {
    A obj_A;
    D() { cout << "D()" << endl; }
};
int main() {
    D d;
}
```

The following is the output of the above example:

```
V()
V2()
B()
C()
A()
D()
```

The above output lists the order in which the C++ run time calls the constructors to create an object of type D.

**Related information**
- "Virtual base classes" on page 289

---

**Destructors**

Destructors are usually used to deallocate memory and do other cleanup for a class object and its class members when the object is destroyed. A destructor is called for a class object when that object passes out of scope or is explicitly deleted.

A destructor is a member function with the same name as its class prefixed by a ~ (tilde). For example:

```cpp
class X {
public:
    // Constructor for class X
    X();
    // Destructor for class X
    ~X();
};
```

A destructor takes no arguments and has no return type. Its address cannot be taken. Destructors cannot be declared const, volatile, const volatile or static. A destructor can be declared virtual or pure virtual.
If no user-defined destructor exists for a class and one is needed, the compiler implicitly declares a destructor. This implicitly declared destructor is an inline public member of its class.

The compiler will implicitly define an implicitly declared destructor when the compiler uses the destructor to destroy an object of the destructor’s class type. Suppose a class \( A \) has an implicitly declared destructor. The following is equivalent to the function the compiler would implicitly define for \( A \):

\[
A::\text{~}A() \ { } \{}
\]

The compiler first implicitly defines the implicitly declared destructors of the base classes and nonstatic data members of a class \( A \) before defining the implicitly declared destructor of \( A \).

A destructor of a class \( A \) is *trivial* if all the following are true:
- It is implicitly defined
- All the direct base classes of \( A \) have trivial destructors
- The classes of all the nonstatic data members of \( A \) have trivial destructors

If any of the above are false, then the destructor is *nontrivial*.

A union member cannot be of a class type that has a nontrivial destructor.

Class members that are class types can have their own destructors. Both base and derived classes can have destructors, although destructors are not inherited. If a base class \( A \) or a member of \( A \) has a destructor, and a class derived from \( A \) does not declare a destructor, a default destructor is generated.

The default destructor calls the destructors of the base class and members of the derived class.

The destructors of base classes and members are called in the reverse order of the completion of their constructor:
1. The destructor for a class object is called before destructors for members and bases are called.
2. Destructors for nonstatic members are called before destructors for base classes are called.
3. Destructors for nonvirtual base classes are called before destructors for virtual base classes are called.

When an exception is thrown for a class object with a destructor, the destructor for the temporary object thrown is not called until control passes out of the catch block.

Destructor are implicitly called when an automatic object (a local object that has been declared auto or register, or not declared as static or extern) or temporary object passes out of scope. They are implicitly called at program termination for constructed external and static objects. Destructors are invoked when you use the `delete` operator for objects created with the `new` operator.

For example:

```
#include <string>

class Y {
private:
```
char * string;
int number;
public:
    // Constructor
    Y(const char*, int);
    // Destructor
    ~Y() { delete[] string; }
};

// Define class Y constructor
Y::Y(const char* n, int a) {
    string = strcpy(new char[strlen(n) + 1], n);
    number = a;
}

int main () {
    // Create and initialize
    // object of class Y
    Y yobj = Y("somestring", 10);
    // ...

    // Destructor ~Y is called before
    // control returns from main()
}

You can use a destructor explicitly to destroy objects, although this practice is not recommended. However to destroy an object created with the placement new operator, you can explicitly call the object’s destructor. The following example demonstrates this:

```cpp
#include <string>
#include <iostream>
using namespace std;

class A {
public:
    A() { cout << "A::A()" << endl; }
    ~A() { cout << "A::~A()" << endl; }
};

int main () {
    char* p = new char[sizeof(A)];
    A* ap = new (p) A;
ap->~A();
delete [] p;
}
```

The statement `A* ap = new (p) A;` dynamically creates a new object of type `A` not in the free store but in the memory allocated by `p`. The statement `delete [] p` will delete the storage allocated by `p`, but the run time will still believe that the object pointed to by `ap` still exists until you explicitly call the destructor of `A` (with the statement `ap->~A();`).

**Related information**
- “The constructor and destructor function attributes” on page 209
- “Free store” on page 315
- “Temporary objects” on page 320

**Pseudo-destructors**

A pseudo-destructor is a destructor of a nonclass type.
Pseudo-destructor syntax

The following example calls the pseudo destructor for an integer type:

typedef int I;
int main() {
  I x = 10;
  x.I::"I()"
  x = 20;
}

The call to the pseudo destructor, x.I::"I()", has no effect at all. Object x has not been destroyed; the assignment x = 20 is still valid. Because pseudo destructors require the syntax for explicitly calling a destructor for a nonclass type to be valid, you can write code without having to know whether or not a destructor exists for a given type.

Related information

- Chapter 12, “Class members and friends (C++ only),” on page 255
- “Scope of class names” on page 249

Free store

Free store is a pool of memory available for you to allocate (and deallocate) storage for objects during the execution of your program. The new and delete operators are used to allocate and deallocate free store, respectively.

You can define your own versions of new and delete for a class by overloading them. You can declare the new and delete operators with additional parameters. When new and delete operate on class objects, the class member operator functions new and delete are called, if they have been declared.

If you create a class object with the new operator, one of the operator functions operator new() or operator new[]( ) (if they have been declared) is called to create the object. An operator new() or operator new[]( ) for a class is always a static class member, even if it is not declared with the keyword static. It has a return type void* and its first parameter must be the size of the object type and have type std::size_t. It cannot be virtual.

Type std::size_t is an implementation-dependent unsigned integral type defined in the standard library header <cstdlib>. When you overload the new operator, you must declare it as a class member, returning type void*, with its first parameter of type std::size_t. You can declare additional parameters in the declaration of operator new() or operator new[]( ). Use the placement syntax to specify values for these parameters in an allocation expression.

The following example overloads two operator new functions:

- X::operator new(size_t sz): This overloads the default new operator by allocating memory with the C function malloc(), and throwing a string (instead of std::bad_alloc) if malloc() fails.
**X::operator new(size_t sz, int location):** This function takes an additional integer parameter, location. This function implements a very simplistic "memory manager" that manages the storage of up to three X objects.

Static array X::buffer holds three Node objects. Each Node object contains a pointer to an X object named data and a Boolean variable named filled. Each X object stores an integer called number.

When you use this new operator, you pass the argument location which indicates the array location of buffer where you want to "create" your new X object. If the array location is not "filled" (the data member of filled is equal to false at that array location), the new operator returns a pointer pointing to the X object located at buffer[location].

```cpp
#include <new>
#include <iostream>

using namespace std;

class X;

struct Node {
   X* data;
   bool filled;
   Node() : filled(false) { }
};

class X {
   static Node buffer[size];

public:
   int number;

   enum { size = 3};

   void* operator new(size_t sz) throw (const char*) {
      void* p = malloc(sz);
      if (sz == 0) throw "Error: malloc() failed";
      cout << "X::operator new(size_t)" << endl;
      return p;
   }

   void* operator new(size_t sz, int location) throw (const char*) {
      cout << "X::operator new(size_t, " << location << ")" << endl;
      void* p = 0;
      if (location < 0 || location >= size || buffer[location].filled == true) {
         throw "Error: buffer location occupied";
      } else {
         p = malloc(sizeof(X));
         if (p == 0) throw "Error: Creating X object failed";
         buffer[location].filled = true;
         buffer[location].data = (X*)p;
      }
      return p;
   }

   static void printbuffer() {
      for (int i = 0; i < size; i++) {
         cout << buffer[i].data->number << endl;
      }
   }
};

Node X::buffer[size];
```
int main() {
    try {
        X* ptr1 = new X;
        X* ptr2 = new(0) X;
        X* ptr3 = new(1) X;
        X* ptr4 = new(2) X;
        ptr2->number = 10000;
        ptr3->number = 10001;
        ptr4->number = 10002;
        X::printbuffer();
        X* ptr5 = new(0) X;
    } catch (const char* message) {
        cout << message << endl;
    }
}

The following is the output of the above example:
X::operator new(size_t)
X::operator new(size_t, 0)
X::operator new(size_t, 1)
X::operator new(size_t, 2)
10000
10001
10002
X::operator new(size_t, 0)
Error: buffer location occupied

The statement X* ptr1 = new X calls X::operator new(sizeof(X)). The statement
X* ptr2 = new(0) X calls X::operator new(sizeof(X),0).

The delete operator destroys an object created by the new operator. The operand of
delete must be a pointer returned by new. If delete is called for an object with a
destructor, the destructor is invoked before the object is deallocated.

If you destroy a class object with the delete operator, the operator function
operator delete() or operator delete[]() (if they have been declared) is called to
destroy the object. An operator delete() or operator delete[]() for a class is
always a static member, even if it is not declared with the keyword static. Its first
parameter must have type void*. Because operator delete() and operator
delete[]() have a return type void, they cannot return a value.

The following example shows the declaration and use of the operator functions
operator new() and operator delete():
#include <cstdlib>
#include <iostream>
using namespace std;

class X {
public:
    void* operator new(size_t sz) throw (const char*) {
        void* p = malloc(sz);
        if (p == 0) throw "malloc() failed";
        return p;
    }

    // single argument
    void operator delete(void* p) {
        cout << "X::operator delete(void*)" << endl;
        free(p);
    }
};
class Y {
    int filler[100];
public:
    // two arguments
    void operator delete(void* p, size_t sz) throw (const char*) {
        cout << "Freeing " << sz << " byte(s)" << endl;
        free(p);
    }
};

int main() {
    X* ptr = new X;

    // call X::operator delete(void*)
    delete ptr;

    Y* yptr = new Y;

    // call Y::operator delete(void*, size_t)
    // with size of Y as second argument
    delete yptr;
}

The above example will generate output similar to the following:
X::operator delete(void*)
Freeing 400 byte(s)

The statement delete ptr calls X::operator delete(void*). The statement delete yptr calls Y::operator delete(void*, size_t).

The result of trying to access a deleted object is undefined because the value of the object can change after deletion.

If new and delete are called for a class object that does not declare the operator functions new and delete, or they are called for a nonclass object, the global operators new and delete are used. The global operators new and delete are provided in the C++ library.

The C++ operators for allocating and deallocating arrays of class objects are operator new[ ]() and operator delete[ ]().

You cannot declare the delete operator as virtual. However you can add polymorphic behavior to your delete operators by declaring the destructor of a base class as virtual. The following example demonstrates this:
#include <iostream>
using namespace std;

struct A {
    virtual ~A() { cout << "~A()" << endl; }
    void operator delete(void* p) {
        cout << "A::operator delete" << endl;
        free(p);
    }
};

struct B : A {
    void operator delete(void* p) {
        cout << "B::operator delete" << endl;
    }
};
free(p);
}

int main() {
    A* ap = new B;
    delete ap;
}

The following is the output of the above example:
"~A()"
B::operator delete

The statement delete ap uses the delete operator from class B instead of class A
because the destructor of A has been declared as virtual.

Although you can get polymorphic behavior from the delete operator, the delete
operator that is statically visible must still be accessible even though another
delete operator might be called. For example, in the above example, the function
A::operator delete(void*) must be accessible even though the example calls
B::operator delete(void*) instead.

Virtual destructors do not have any affect on deallocation operators for arrays
(operator delete[]()). The following example demonstrates this:
#include <iostream>
using namespace std;

struct A {
  virtual "~A()" { cout << "~A()" << endl; }
  void operator delete[](void* p, size_t) {
    cout << "A::operator delete[]" << endl;
    ::delete [] p;
  }
};

struct B : A {
  void operator delete[](void* p, size_t) {
    cout << "B::operator delete[]" << endl;
    ::delete [] p;
  }
};

int main() {
    A* bp = new B[3];
    delete[] bp;
}

The behavior of the statement delete[] bp is undefined.

When you overload the delete operator, you must declare it as class member,
returning type void, with the first parameter having type void*, as described
above. You can add a second parameter of type size_t to the declaration. You can
only have one operator delete() or operator delete[]() for a single class.

Related information
- "The new operator (C++ only)" on page 140
- "The delete operator (C++ only)" on page 142
Temporary objects

Temporary objects are unnamed objects created on the stack by the compiler. They are used during reference initialization and during evaluation of expressions including standard type conversions, argument passing, function returns, and evaluation of the `throw` expression.

When a temporary object is created to initialize a reference variable, the name of the temporary object has the same scope as that of the reference variable. When a temporary object is created during the evaluation of a full-expression (an expression that is not a subexpression of another expression), it is destroyed as the last step in its evaluation that lexically contains the point where it was created.

There are two exceptions in the destruction of full-expressions:

- The expression appears as an initializer for a declaration defining an object: the temporary object is destroyed when the initialization is complete.
- A reference is bound to a temporary object: the temporary object is destroyed at the end of the reference’s lifetime.

If a temporary object is created for a class with constructors, the compiler calls the appropriate (matching) constructor to create the temporary object.

When a temporary object is destroyed and a destructor exists, the compiler calls the destructor to destroy the temporary object. When you exit from the scope in which the temporary object was created, it is destroyed. If a reference is bound to a temporary object, the temporary object is destroyed when the reference passes out of scope unless it is destroyed earlier by a break in the flow of control. For example, a temporary object created by a constructor initializer for a reference member is destroyed on leaving the constructor.

In cases where such temporary objects are redundant, the compiler does not construct them, in order to create more efficient optimized code. This behavior could be a consideration when you are debugging your programs, especially for memory problems.

Related information

- "Arguments of catch blocks” on page 369
- "Initialization of references (C++ only)” on page 100
- “Cast expressions” on page 144
- “Function return values” on page 202

User-defined conversions

User-defined conversions allow you to specify object conversions with constructors or with conversion functions. User-defined conversions are implicitly used in addition to standard conversions for conversion of initializers, functions arguments, function return values, expression operands, expressions controlling iteration, selection statements, and explicit type conversions.

There are two types of user-defined conversions:

- Conversion by constructor
- Conversion functions
The compiler can use only one user-defined conversion (either a conversion constructor or a conversion function) when implicitly converting a single value. The following example demonstrates this:

class A {
    int x;
public:
    operator int() { return x; }
};

class B {
    A y;
public:
    operator A() { return y; }
};

int main () {
    B b_obj;
    // int i = b_obj;
    int j = A(b_obj);
}

The compiler would not allow the statement int i = b_obj. The compiler would have to implicitly convert b_obj into an object of type A (with B::operator A()), then implicitly convert that object to an integer (with A::operator int()). The statement int j = A(b_obj) explicitly converts b_obj into an object of type A, then implicitly converts that object to an integer.

User-defined conversions must be unambiguous, or they are not called. A conversion function in a derived class does not hide another conversion function in a base class unless both conversion functions convert to the same type. Function overload resolution selects the most appropriate conversion function. The following example demonstrates this:

class A {
    int a_int;
    char* a_carp;
public:
    operator int() { return a_int; }
    operator char*() { return a_carp; }
};

class B : public A {
    float b_float;
    char* b_carp;
public:
    operator float() { return b_float; }
    operator char*() { return b_carp; }
};

int main () {
    B b_obj;
    // long a = b_obj;
    char* c_p = b_obj;
}

The compiler would not allow the statement long a = b_obj. The compiler could either use A::operator int() or B::operator float() to convert b_obj into a long. The statement char* c_p = b_obj uses B::operator char*() to convert b_obj into a char* because B::operator char*() hides A::operator char*().

When you call a constructor with an argument and you have not defined a constructor accepting that argument type, only standard conversions are used to convert the argument to another argument type acceptable to a constructor for that
class. No other constructors or conversions functions are called to convert the argument to a type acceptable to a constructor defined for that class. The following example demonstrates this:

class A {
public:
    A() {}
    A(int) {}
};

int main() {
    A a1 = 1.234;
    // A moocow = "text string";
}

The compiler allows the statement A a1 = 1.234. The compiler uses the standard conversion of converting 1.234 into an int, then implicitly calls the converting constructor A(int). The compiler would not allow the statement A moocow = "text string"; converting a text string to an integer is not a standard conversion.

Related information
• Chapter 5, “Type conversions,” on page 107

Conversion by constructor

A converting constructor is a single-parameter constructor that is declared without the function specifier explicit. The compiler uses converting constructors to convert objects from the type of the first parameter to the type of the converting constructor’s class. The following example demonstrates this:

class Y {
    int a, b;
    char* name;
public:
    Y(int i) {}
    Y(const char* n, int j = 0) {}
};

void add(Y) {};

int main() {
    // equivalent to
    // obj1 = Y(2)
    Y obj1 = 2;

    // equivalent to
    // obj2 = Y("somestring",0)
    Y obj2 = "somestring";

    // equivalent to
    // obj1 = Y(10)
    obj1 = 10;

    // equivalent to
    // add(Y(5))
    add(5);
}

The above example has the following two converting constructors:
• Y(int i) which is used to convert integers to objects of class Y.
• Y(const char* n, int j = 0) which is used to convert pointers to strings to objects of class Y.
The compiler will not implicitly convert types as demonstrated above with constructors declared with the `explicit` keyword. The compiler will only use explicitly declared constructors in `new` expressions, the `static_cast` expressions and explicit casts, and the initialization of bases and members. The following example demonstrates this:

```cpp
class A {
public:
    explicit A() {}
    explicit A(int);
};

int main() {
    A z;
    // A y = 1;
    A x = A(1);
    A w(1);
    A* v = new A(1);
    A u = (A)1;
    A t = static_cast<A>(1);
}
```

The compiler would not allow the statement `A y = 1` because this is an implicit conversion; class `A` has no conversion constructors.

A copy constructor is a converting constructor.

**Related information**
- [“The new operator (C++ only)” on page 140](#)
- [“The static_cast operator (C++ only)” on page 126](#)

**The explicit specifier**

The explicit function specifier controls unwanted implicit type conversions. It can only be used in declarations of constructors within a class declaration. For example, except for the default constructor, the constructors in the following class are converting constructors.

```cpp
class A {
public:
    A();
    A(int);
    A(const char*, int = 0);
};
```

The following declarations are legal.

```
A c = 1;
A d = "Venditti";
```

The first declaration is equivalent to `A c = A(1)`.

If you declare the constructor of the class with the `explicit` keyword, the previous declarations would be illegal.

For example, if you declare the class as:

```cpp
class A {
public:
    explicit A();
    explicit A(int);
    explicit A(const char*, int = 0);
};
```
You can only assign values that match the values of the class type.

For example, the following statements will be legal:

```cpp
A a1;
A a2 = A(1);
A a3(1);
A a4 = A("Venditti");
A* p = new A(1);
A a5 = (A)1;
A a6 = static_cast<A>(1);
```

**Related information**

- “Conversion by constructor” on page 322

### Conversion functions

You can define a member function of a class, called a conversion function, that converts from the type of its class to another specified type.

#### Conversion function syntax

```cpp
operator conversion_type class::operator const\ volatile

```

A conversion function that belongs to a class X specifies a conversion from the class type X to the type specified by the conversion_type. The following code fragment shows a conversion function called operator int():

```cpp
class Y {
  int b;
public:
  operator int();
};
Y::operator int() { return b; }
void f(Y obj) {
  int i = int(obj);
  int j = (int)obj;
  int k = i + obj;
}
```

All three statements in function f(Y) use the conversion function Y::operator int().

Classes, enumerations, typedef names, function types, or array types cannot be declared or defined in the conversion_type. You cannot use a conversion function to convert an object of type A to type A, to a base class of A, or to void.

Conversion functions have no arguments, and the return type is implicitly the conversion type. Conversion functions can be inherited. You can have virtual conversion functions but not static ones.

**Related information**
Copy constructors

The copy constructor lets you create a new object from an existing one by initialization. A copy constructor of a class A is a non-template constructor in which the first parameter is of type A&, const A&, volatile A&, or const volatile A&, and the rest of its parameters (if there are any) have default values.

If you do not declare a copy constructor for a class A, the compiler will implicitly declare one for you, which will be an inline public member.

The following example demonstrates implicitly defined and user-defined copy constructors:

```cpp
#include <iostream>
using namespace std;

struct A {
    int i;
    A() : i(10) {} 
};

struct B {
    int j;
    B() : j(20) {
        cout << "Constructor B(), j = " << j << endl;
    }

    B(B& arg) : j(arg.j) {
        cout << "Copy constructor B(B&), j = " << j << endl;
    }

    B(const B&, int val = 30) : j(val) {
        cout << "Copy constructor B(const B&, int), j = " << j << endl;
    }
};

struct C {
    C() {} 
    C(C&) {} 
};

int main() {
    A a;
    A a1(a);
    B b;
    const B b_const;
    B b1(b);
    B b2(b_const);
    const C c_const;
    // C c1(c_const);
}
```

The following is the output of the above example:

Constructor B(), j = 20
Constructor B(), j = 20
Copy constructor B(B&), j = 20
Copy constructor B(const B&, int), j = 30
The statement `A a1(a)` creates a new object from `a` with an implicitly defined copy constructor. The statement `B b1(b)` creates a new object from `b` with the user-defined copy constructor `B::B(B&)`. The statement `B b2(b_const)` creates a new object with the copy constructor `B::B(const B&, int)`. The compiler would not allow the statement `C c1(c_const)` because a copy constructor that takes as its first parameter an object of type `const C&` has not been defined.

The implicitly declared copy constructor of a class `A` will have the form `A::A(const A&)` if the following are true:

- The direct and virtual bases of `A` have copy constructors whose first parameters have been qualified with `const` or `const volatile`
- The nonstatic class type or array of class type data members of `A` have copy constructors whose first parameters have been qualified with `const` or `const volatile`

If the above are not true for a class `A`, the compiler will implicitly declare a copy constructor with the form `A::A(A&)`.

The compiler cannot allow a program in which the compiler must implicitly define a copy constructor for a class `A` and one or more of the following are true:

- Class `A` has a nonstatic data member of a type which has an inaccessible or ambiguous copy constructor.
- Class `A` is derived from a class which has an inaccessible or ambiguous copy constructor.

The compiler will implicitly define an implicitly declared constructor of a class `A` if you initialize an object of type `A` or an object derived from class `A`.

An implicitly defined copy constructor will copy the bases and members of an object in the same order that a constructor would initialize the bases and members of the object.

**Related information**

- "Overview of constructors and destructors" on page 303

---

**Copy assignment operators**

The *copy assignment operator* lets you create a new object from an existing one by initialization. A copy assignment operator of a class `A` is a nonstatic non-template member function that has one of the following forms:

- `A::operator=(A)`
- `A::operator=(A&)`
- `A::operator=(const A&)`
- `A::operator=(volatile A&)`
- `A::operator=(const volatile A&)`

If you do not declare a copy assignment operator for a class `A`, the compiler will implicitly declare one for you which will be inline public.

The following example demonstrates implicitly defined and user-defined copy assignment operators:

```cpp
#include <iostream>
using namespace std;
```
struct A {
    A& operator=(const A&) {
        cout << "A::operator=(const A&)" << endl;
        return *this;
    }

    A& operator=(A&) {
        cout << "A::operator=(A&)" << endl;
        return *this;
    }
};
class B {
    A a;
};

struct C {
    C& operator=(C&) {
        cout << "C::operator=(C&)" << endl;
        return *this;
    }

    C() { }
};

int main() {
    B x, y;
    x = y;

    A w, z;
    w = z;

    C i;
    const C j();
    // i = j;
}

The following is the output of the above example:
A::operator=(const A&)
A::operator=(A&)

The assignment x = y calls the implicitly defined copy assignment operator of B, which calls the user-defined copy assignment operator A::operator=(const A&). The assignment w = z calls the user-defined operator A::operator=(A&). The compiler will not allow the assignment i = j because an operator C::operator=(const C&) has not been defined.

The implicitly declared copy assignment operator of a class A will have the form A& A::operator=(A&) if the following are true:
- A direct or virtual base B of class A has a copy assignment operator whose parameter is of type const B&, const volatile B&, or B.
- A non-static class type data member of type X that belongs to class A has a copy constructor whose parameter is of type const X&, const volatile X&, or X.

If the above are not true for a class A, the compiler will implicitly declare a copy assignment operator with the form A& A::operator=(A&).

The implicitly declared copy assignment operator returns a reference to the operator's argument.

The copy assignment operator of a derived class hides the copy assignment operator of its base class.
The compiler cannot allow a program in which the compiler must implicitly define a copy assignment operator for a class A and one or more of the following are true:

- Class A has a nonstatic data member of a const type or a reference type
- Class A has a nonstatic data member of a type which has an inaccessible copy assignment operator
- Class A is derived from a base class with an inaccessible copy assignment operator.

An implicitly defined copy assignment operator of a class A will first assign the direct base classes of A in the order that they appear in the definition of A. Next, the implicitly defined copy assignment operator will assign the nonstatic data members of A in the order of their declaration in the definition of A.

Related information

- "Assignment expressions" on page 158
Chapter 15. Templates (C++ only)

A template describes a set of related classes or set of related functions in which a list of parameters in the declaration describe how the members of the set vary. The compiler generates new classes or functions when you supply arguments for these parameters; this process is called template instantiation, and is described in detail in “Template instantiation” on page 348. This class or function definition generated from a template and a set of template parameters is called a specialization, as described in “Template specialization” on page 351.

Template declaration syntax

The compiler accepts and silently ignores the export keyword on a template.

The template_parameter_list is a comma-separated list of template parameters, which are described in “Template parameters” on page 330.

The declaration is one of the following:

• a declaration or definition of a function or a class
• a definition of a member function or a member class of a class template
• a definition of a static data member of a class template
• a definition of a static data member of a class nested within a class template
• a definition of a member template of a class or class template

The identifier of a type is defined to be a type_name in the scope of the template declaration. A template declaration can appear as a namespace scope or class scope declaration.

The following example demonstrates the use of a class template:

```cpp
template<class T> class Key
{
    T k;
    T* kptr;
    int length;
public:
    Key(T);
    // ...
};
```

Suppose the following declarations appear later:

```cpp
Key<int> i;
Key<char*> c;
Key<mytype> m;
```

The compiler would create three instances of class Key. The following table shows the definitions of these three class instances if they were written out in source form as regular classes, not as templates:
class Key<int> i;  
public:
    Key(int);
// ...

class Key<char*> c;  
public:
    Key(char*);
// ...

class Key<mytype> m;  
public:
    Key(mytype);
// ...

Note that these three classes have different names. The arguments contained within the angle braces are not just the arguments to the class names, but part of the class names themselves. Key<int> and Key<char*> are class names.

Template parameters

There are three kinds of template parameters:

- **Type template parameters**
- **Non-type template parameters**
- **Template template parameters**

You can interchange the keywords class and typename in a template parameter declaration. You cannot use storage class specifiers (static and auto) in a template parameter declaration.

Type template parameters

**Type template parameter declaration syntax**

```cpp
class typename identifier = type
```

The `identifier` is the name of a type.

Related information

- "The typename keyword" on page 359

Non-type template parameters

The syntax of a non-type template parameter is the same as a declaration of one of the following types:

- integral or enumeration
- pointer to object or pointer to function
- reference to object or reference to function
- pointer to member

Non-type template parameters that are declared as arrays or functions are converted to pointers or pointers to functions, respectively. The following example demonstrates this:

```cpp
template<int a[4]> struct A {};
template<int f(int)> struct B {};

int i;
```
int g(int) { return 0;}
A<&i> x;
B<&g> y;

The type deduced from &i is int *, and the type deduced from &g is int (*)(int).
You may qualify a non-type template parameter with const or volatile.
You cannot declare a non-type template parameter as a floating point, class, or void type.

Non-type template parameters are not lvalues.

Related information
• “Type qualifiers” on page 68
• “Lvalues and rvalues” on page 115

Template template parameters

Template template parameter declaration syntax

```cpp
<template-parameter-list> identifier <id-expression>
```

The following example demonstrates a declaration and use of a template template parameter:
```cpp
template<template <class T> class X> class A { };
template<class T> class B { };
A<B> a;
```

Related information
• “Template parameters” on page 330

Default arguments for template parameters

Template parameters may have default arguments. The set of default template arguments accumulates over all declarations of a given template. The following example demonstrates this:
```cpp
template<class T, class U = int> class A;
template<class T = float, class U> class A;

template<class T, class U> class A { 
    public:
        T x;
        U y;
    
};
A<> a;
```

The type of member a.x is float, and the type of a.y is int.

You cannot give default arguments to the same template parameters in different declarations in the same scope. For example, the compiler will not allow the following:
template<class T = char> class X;
template<class T = char> class X {};

If one template parameter has a default argument, then all template parameters following it must also have default arguments. For example, the compiler will not allow the following:
template<class T = char, class U, class V = int> class X {};

Template parameter U needs a default argument or the default for T must be removed.

The scope of a template parameter starts from the point of its declaration to the end of its template definition. This implies that you may use the name of a template parameter in other template parameter declarations and their default arguments. The following example demonstrates this:
template<class T = int> class A;
template<class T = float> class B;
template<class V, V obj> class C;
// a template parameter (T) used as the default argument
// to another template parameter (U)
template<class T, class U = T> class D {};

Related information
- "Template parameters" on page 330

Template arguments

There are three kinds of template arguments corresponding to the three types of template parameters:
- Template type arguments
- Template non-type arguments
- Template template arguments

A template argument must match the type and form specified by the corresponding parameter declared in the template.

To use the default value of a template parameter, you omit the corresponding template argument. However, even if all template parameters have defaults, you still must use the <> brackets. For example, the following will yield a syntax error:
template<class T = int> class X {};
X<> a;
X b;

The last declaration, X b, will yield an error.

Related information

Template type arguments

You cannot use one of the following as a template argument for a type template parameter:
- a local type
- a type with no linkage
- an unnamed type
- a type compounded from any of the above types
If it is ambiguous whether a template argument is a type or an expression, the template argument is considered to be a type. The following example demonstrates this:

```cpp
template<class T> void f() { }
template<int i> void f() { }

int main()
{
    f<int()>();
}
```

The function call `f<int()>()` calls the function with `T` as a template argument – the compiler considers `int()` as a type – and therefore implicitly instantiates and calls the first `f()`.

**Related information**
- “Block/local scope” on page 2
- “No linkage” on page 9
- “Bit field members” on page 58
- “typedef definitions” on page 66

### Template non-type arguments

A non-type template argument provided within a template argument list is an expression whose value can be determined at compile time. Such arguments must be constant expressions, addresses of functions or objects with external linkage, or addresses of static class members. Non-type template arguments are normally used to initialize a class or to specify the sizes of class members.

For non-type integral arguments, the instance argument matches the corresponding template parameter as long as the instance argument has a value and sign appropriate to the parameter type.

For non-type address arguments, the type of the instance argument must be of the form `identifier` or `&identifier`, and the type of the instance argument must match the template parameter exactly, except that a function name is changed to a pointer to function type before matching.

The resulting values of non-type template arguments within a template argument list form part of the template class type. If two template class names have the same template name and if their arguments have identical values, they are the same class.

In the following example, a class template is defined that requires a non-type template `int` argument as well as the type argument:

```cpp
template<class T, int size> class Myfilebuf
{
    T* filepos;
    static int array[size];
public:
    Myfilebuf() { /* ... */
    "Myfilebuf();
    advance(); // function defined elsewhere in program
};
```

In this example, the template argument `size` becomes a part of the template class name. An object of such a template class is created with both the type argument `T` of the class and the value of the non-type template argument `size`.
An object \( x \), and its corresponding template class with arguments \( \text{double} \) and \( \text{size}=200 \), can be created from this template with a value as its second template argument:

\[
\text{Myfilebuf<double,200> } x;
\]

\( x \) can also be created using an arithmetic expression:

\[
\text{Myfilebuf<double,10*20> } x;
\]

The objects created by these expressions are identical because the template arguments evaluate identically. The value 200 in the first expression could have been represented by an expression whose result at compile time is known to be equal to 200, as shown in the second construction.

**Note:** Arguments that contain the < symbol or the > symbol must be enclosed in parentheses to prevent either symbol from being parsed as a template argument list delimiter when it is in fact being used as a relational operator. For example, the arguments in the following definition are valid:

\[
\text{Myfilebuf<double, (75>25)> } x; \quad // \text{ valid}
\]

The following definition, however, is not valid because the greater than operator (>) is interpreted as the closing delimiter of the template argument list:

\[
\text{Myfilebuf<double, 75>25> } x; \quad // \text{ error}
\]

If the template arguments do not evaluate identically, the objects created are of different types:

\[
\text{Myfilebuf<double,200> } x; \quad // \text{ create object } x \text{ of class } \text{Myfilebuf<double,200>}
\]

\[
\text{Myfilebuf<double,200.0> } y; \quad // \text{ error, 200.0 is a double, not an int}
\]

The instantiation of \( y \) fails because the value 200.0 is of type \( \text{double} \), and the template argument is of type \( \text{int} \).

The following two objects:

\[
\text{Myfilebuf<double,128> } x \\
\text{Myfilebuf<double,512> } y
\]

are objects of separate template specializations. Referring either of these objects later with \( \text{Myfilebuf<double>} \) is an error.

A class template does not need to have a type argument if it has non-type arguments. For example, the following template is a valid class template:

```
\text{template<int } i\text{> class } C \\
\{ \\
\text{public:} \\
\text{ int } k; \\
\text{ C() } \{ \text{ } k = i; \} \\
\};
```

This class template can be instantiated by declarations such as:

\[
\text{class } C<100> \\
\text{class } C<200>
\]

Again, these two declarations refer to distinct classes because the values of their non-type arguments differ.
Template template arguments

A template argument for a template template parameter is the name of a class template.

When the compiler tries to find a template to match the template template argument, it only considers primary class templates. (A primary template is the template that is being specialized.) The compiler will not consider any partial specialization even if their parameter lists match that of the template template parameter. For example, the compiler will not allow the following code:

```cpp
template<class T, int i> class A {
  int x;
};

template<class T> class A<T, 5> {
  short x;
};

template<template<class T> class U> class B1 {
};

B1<A> c;
```

The compiler will not allow the declaration `B1<A> c`. Although the partial specialization of `A` seems to match the template template parameter `U` of `B1`, the compiler considers only the primary template of `A`, which has different template parameters than `U`.

The compiler considers the partial specializations based on a template template argument once you have instantiated a specialization based on the corresponding template template parameter. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

class A {
  int x;
};

class A<int> {
  short x;
};

class A<template<class T> class U> {
  U<int, char> i;
  U<char, char> j;
};

B<A> c;

int main() {
  cout << typeid(c.i.x).name() << endl;
  cout << typeid(c.j.x).name() << endl;
}
```

The following is the output of the above example:
The declaration `V<int, char> i` uses the partial specialization while the declaration `V<char, char> j` uses the primary template.

**Related information**
- "Partial specialization" on page 356
- "Template instantiation" on page 348

---

### Class templates

The relationship between a class template and an individual class is like the relationship between a class and an individual object. An individual class defines how a group of objects can be constructed, while a class template defines how a group of classes can be generated.

Note the distinction between the terms *class template* and *template class*:

**Class template**

is a template used to generate template classes. You cannot declare an object of a class template.

**Template class**

is an instance of a class template.

A template definition is identical to any valid class definition that the template might generate, except for the following:

- The class template definition is preceded by
  
  ```
  template< template-parameter-list >
  ```

  where `template-parameter-list` is a comma-separated list of one or more of the following kinds of template parameters:
  - type
  - non-type
  - template

- Types, variables, constants and objects within the class template can be declared using the template parameters as well as explicit types (for example, `int` or `char`).

A class template can be declared without being defined by using an elaborated type specifier. For example:

```
template<class L, class T> class Key;
```

This reserves the name as a class template name. All template declarations for a class template must have the same types and number of template arguments. Only one template declaration containing the class definition is allowed.

**Note:** When you have nested template argument lists, you must have a separating space between the `>` at the end of the inner list and the `>` at the end of the outer list. Otherwise, there is an ambiguity between the extraction operator `>>` and two template list delimiters `>`.  

```
template<class L, class T> class Key { /* ... */};
template<class L> class Vector { /* ... */};
```
int main()
{
  class Key<int, Vector<int>> my_key_vector;
  // implicitly instantiates template
}

Objects and function members of individual template classes can be accessed by any of the techniques used to access ordinary class member objects and functions. Given a class template:

```cpp
template<class T> class Vehicle
{
public:
  Vehicle() { /* ... */ }  // constructor
  ~Vehicle() { }           // destructor
  T kind[16];
  T* drive();
  static void roadmap();
  // ...
};
```

and the declaration:

```cpp
Vehicle<char> bicycle; // instantiates the template
```

the constructor, the constructed object, and the member function drive() can be accessed with any of the following (assuming the standard header file string.h is included in the program file):

<table>
<thead>
<tr>
<th>constructor</th>
<th>Vehicle&lt;char&gt; bicycle;</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>// constructor called automatically,</td>
</tr>
<tr>
<td></td>
<td>// object bicycle created</td>
</tr>
</tbody>
</table>

| object bicycle | strcpy (bicycle.kind, "10 speed"); |
|               | bicycle.kind[0] = '2'; |

| function drive() | char* n = bicycle.drive(); |

| function roadmap() | Vehicle<char>::roadmap(); |

Related information

- “Declaring class types” on page 246
- “Scope of class names” on page 249

**Class template declarations and definitions**

A class template must be declared before any instantiation of a corresponding template class. A class template definition can only appear once in any single translation unit. A class template must be defined before any use of a template class that requires the size of the class or refers to members of the class.

In the following example, the class template Key is declared before it is defined. The declaration of the pointer keyiptr is valid because the size of the class is not needed. The declaration of keyi, however, causes an error.

```cpp
template <class L> class Key;  // class template declared,
                              // not defined yet
class Key<int> *keyiptr;      // declaration of pointer
class Key<int> keyi;          // error, cannot declare keyi
```
// without knowing size
//
template <class L> class Key { /* ... */};

If a template class is used before the corresponding class template is defined, the compiler issues an error. A class name with the appearance of a template class name is considered to be a template class. In other words, angle brackets are valid in a class name only if that class is a template class.

The previous example uses the elaborated type specifier class to declare the class template key and the pointer keyiptr. The declaration of keyiptr can also be made without the elaborated type specifier.

```
template <class L> class Key;  // class template declared,
// not defined yet
//
Key<int>* keyiptr;  // declaration of pointer
//
Key<int> keyi;  // error, cannot declare keyi
// without knowing size
//
template <class L> class Key { /* ... */};
```

**Related information**

- "Class templates” on page 336

### Static data members and templates

Each class template instantiation has its own copy of any static data members. The static declaration can be of template argument type or of any defined type.

You must separately define static members. The following example demonstrates this:

```
template <class T> class K
{
    public:
        static T x;
    }

template <class T> T K<T>::x;

int main()
{
    K<int>::x = 0;
}
```

The statement template T K::x defines the static member of class K, while the statement in the main() function assigns a value to the data member for K <int>.

**Related information**

- "Static members” on page 265

### Member functions of class templates

You may define a template member function outside of its class template definition.

When you call a member function of a class template specialization, the compiler will use the template arguments that you used to generate the class template. The following example demonstrates this:
template<class T> class X {
    public:
        T operator+(T);
};

template<class T> T X<T>::operator+(T arg1) {
    return arg1;
};

int main() {
    X<char> a;
    X<int> b;
    a + 'z';
    b + 4;
}

The overloaded addition operator has been defined outside of class X. The statement a + 'z' is equivalent to a.operator+( 'z' ). The statement b + 4 is equivalent to b.operator+(4).

**Related information**
- [“Member functions” on page 257](#)

**Friends and templates**

There are four kinds of relationships between classes and their friends when templates are involved:

- **One-to-many**: A non-template function may be a friend to all template class instantiations.
- **Many-to-one**: All instantiations of a template function may be friends to a regular non-template class.
- **One-to-one**: A template function instantiated with one set of template arguments may be a friend to one template class instantiated with the same set of template arguments. This is also the relationship between a regular non-template class and a regular non-template friend function.
- **Many-to-many**: All instantiations of a template function may be a friend to all instantiations of the template class.

The following example demonstrates these relationships:

class B{
    template<class V> friend int j();
}

template<class S> g();

template<class T> class A {
    friend int e();
    friend int f(T);
    friend int g<T>();
    template<class U> friend int h();
};

- Function e() has a one-to-many relationship with class A. Function e() is a friend to all instantiations of class A.
- Function f() has a one-to-one relationship with class A. The compiler will give you a warning for this kind of declaration similar to the following:

  The friend function declaration "f" will cause an error when the enclosing template class is instantiated with arguments that declare a friend function that does not match an existing definition. The function declares only one function because it is not a template but the function type depends on one or more template parameters.
• Function \( g() \) has a one-to-one relationship with class \( A \). Function \( g() \) is a function template. It must be declared before here or else the compiler will not recognize \( g<T> \) as a template name. For each instantiation of \( A \) there is one matching instantiation of \( g() \). For example, \( g<int> \) is a friend of \( A<int> \).

• Function \( h() \) has a many-to-many relationship with class \( A \). Function \( h() \) is a function template. For all instantiations of \( A \) all instantiations of \( h() \) are friends.

• Function \( j() \) has a many-to-one relationship with class \( B \).

These relationships also apply to friend classes.

Related information
• "Friends" on page 272

Function templates

A function template defines how a group of functions can be generated.

A non-template function is not related to a function template, even though the non-template function may have the same name and parameter profile as those of a specialization generated from a template. A non-template function is never considered to be a specialization of a function template.

The following example implements the QuickSort algorithm with a function template named quicksort:

```c++
#include <iostream>
#include <cstdlib>
using namespace std;

template<class T>
void quicksort(T a[], const int& leftarg, const int& rightarg)
{
    if (leftarg < rightarg)
    {
        T pivotvalue = a[leftarg];
        int left = leftarg - 1;
        int right = rightarg + 1;

        for(;;) {
            while (a[--right] > pivotvalue);
            while (a[++left] < pivotvalue);
            if (left >= right) break;

            T temp = a[right];
            a[right] = a[left];
            a[left] = temp;
        }

        int pivot = right;
        quicksort(a, leftarg, pivot);
        quicksort(a, pivot + 1, rightarg);
    }
}

int main(void) {
    int sortme[10];

    for (int i = 0; i < 10; i++) {
        sortme[i] = rand();
        cout << sortme[i] << " ";
    };
    cout << endl;
```
quicksort<int>(sortme, 0, 10 - 1);
for (int i = 0; i < 10; i++) cout << sortme[i] << " 
";
    cout << endl;
return 0;
}

The above example will have output similar to the following:
16838 5758 10113 17515 31051 5627 23010 7419 16212 4086
4086 5627 5758 7419 10113 16212 16838 17515 23010 31051

This QuickSort algorithm will sort an array of type T (whose relational and assignment operators have been defined). The template function takes one template argument and three function arguments:
- the type of the array to be sorted, T
- the name of the array to be sorted, a
- the lower bound of the array, leftarg
- the upper bound of the array, rightarg

In the above example, you can also call the quicksort() template function with the following statement:
quicksort(sortme, 0, 10 - 1);

You may omit any template argument if the compiler can deduce it by the usage and context of the template function call. In this case, the compiler deduces that sortme is an array of type int.

**Template argument deduction**

When you call a template function, you may omit any template argument that the compiler can determine or deduce by the usage and context of that template function call.

The compiler tries to deduce a template argument by comparing the type of the corresponding template parameter with the type of the argument used in the function call. The two types that the compiler compares (the template parameter and the argument used in the function call) must be of a certain structure in order for template argument deduction to work. The following lists these type structures:

- T
- const T
- volatile T
- T&
- T*
- T[10]
- A<T>
- C(*)(T)
- T(*)(T)
- T(*)(U)
- T C::*
- C T::*
- T U::*
- T (C::*)(T)
- C (T::*)(U)
- D (C::*)(T)
- C (T::*)(U)
- T (C::*)(T)
- T (U::*)(T)
- T (U::*)(V)
The following example demonstrates the use of each of these type structures. The example declares a template function using each of the above structures as an argument. These functions are then called (without template arguments) in order of declaration. The example outputs the same list of type structures:

```cpp
#include <iostream>
using namespace std;

template<class T> class A
{ 
};
template<int i> class B
{ 
};
class C
{ 
    public:
    int x;
};
class D
{ 
    public:
    C y;
    int z;
};

template<class T> void f(T) { cout << "T" << endl; };
template<class T> void f1(const T) { cout << "const T" << endl; };
template<class T> void f2(volatile T) { cout << "volatile T" << endl; };
template<class T> void g(T*) { cout << "T*" << endl; };
template<class T> void g(T&) { cout << "T&" << endl; };
template<class T> void g1(T[10]) { cout << "T[10]" << endl; };
template<class T> void h1(A<T>) { cout << "A<T>" << endl; };

void test_1()
{ 
    A<char> a;
    C c;
    f(c); f1(c); f2(c);
    g(c); g(&c); g1(&c);
    h1(a);
}

template<class T> void j(C(*)(T)) { cout << "C(*)(T)" << endl; };
```
template<class T> void j(T(*)()) { cout << "T(*)()" << endl; }
template<class T, class U> void j(T(*)(U)) { cout << "T(*) (U)" << endl; }

void test_2() {
    C (*c_pfunct1)(int);
    C (*c_pfunct2)(void);
    int (*c_pfunct3)(int);
    j(c_pfunct1);
    j(c_pfunct2);
    j(c_pfunct3);
}

void test_3() {
    k(&C::x);
    k(&D::y);
    k(&D::z);
}

template<class T>
void m(T (C::*)(T))
{ cout << "T (C::*)(T)" << endl; }
template<class T>
void m(C (T::*)(T))
{ cout << "C (T::*)(T)" << endl; }
template<class T, class U>
void m(D (C::*)(T))
{ cout << "D (C::*)(T)" << endl; }

void test_4() {
    int (C::*f_membp1)(void);
    C (D::*f_membp2)(void);
    D (C::*f_membp3)(int);
    m(f_membp1);
    m(f_membp2);
    m(f_membp3);
}

void test_5() {
    C array[10][20];
    n(array);
    B<20> b;
    n(b);
}

template<template<class> class TT, class T> void p1(TT<T>)

Chapter 15. Templates (C++ only) 343
template<template<int> class TT, int i> void p2(TT<i>)
{
    cout << "TT<i>" << endl;
}

template<template<class> class TT> void p3(TT<C>)
{
    cout << "TT<C>" << endl;
}

int main()
{
    test_6();
    test_1(); test_2(); test_3(); test_4(); test_5(); test_6();
}

Deducing type template arguments

The compiler can deduce template arguments from a type composed of several of the listed type structures. The following example demonstrates template argument deduction for a type composed of several type structures:

template<class T> class Y {
};

template<class T, int i> class X {
    public:
        Y<T> f(char[20][i]) { return x; }
        Y<T> x;
};

template<template<class> class T> void g(T<U>(V::*)(W*[20][i]))
{
}

int main()
{
    Y<int>::Z<int>b;
    Y<int>c;
    g(b);
}

The type Y<int> (X<int, 20>::*p)(char[20][20])T<U>(V::*)(W*[20][i]) is based on the type structure T (U::*)(V):
• T is Y<int>
• U is X<int, 20>
• V is char[20][20]

If you qualify a type with the class to which that type belongs, and that class (a nested name specifier) depends on a template parameter, the compiler will not deduce a template argument for that parameter. If a type contains a template argument that cannot be deduced for this reason, all template arguments in that type will not be deduced. The following example demonstrates this:

template<class T, class U, class V>
void h(typename Y<T>::template Z<U>, Y<T>, Y<V>)
{
}

int main()
{
    Y<int>::Z<char> a;
    Y<int>b;
    Y<float>c;

    h<int, char, float>(a, b, c);
    h<int, char>(a, b, c);
    // h<int>(a, b, c);
}
The compiler will not deduce the template arguments \( T \) and \( U \) in `typedef Y<T>::template Z<U>` (but it will deduce the \( T \) in `Y<T>`). The compiler would not allow the template function call `h<int>(a, b, c)` because \( U \) is not deduced by the compiler.

The compiler can deduce a function template argument from a pointer to function or pointer to member function argument given several overloaded function names. However, none of the overloaded functions may be function templates, nor can more than one overloaded function match the required type. The following example demonstrates this:

```cpp
template<class T> void f(void(*) (T,int)) { }

template<class T> void g1(T, int) { }

void g2(int, int) { }
void g2(char, int) { }

void g3(int, int, int) { }
void g3(float, int) { }

int main() {  
  // f(&g1);  
  // f(&g2);  
  f(&g3);  
}
```

The compiler would not allow the call `f(&g1)` because `g1()` is a function template. The compiler would not allow the call `f(&g2)` because both functions named `g2()` match the type required by `f()`.

The compiler cannot deduce a template argument from the type of a default argument. The following example demonstrates this:

```cpp
template<class T> void f(T = 2, T = 3) { }

int main() {  
  f(6);  
  // f();  
  f<int>();  
}
```

The compiler allows the call `f(6)` because the compiler deduces the template argument (int) by the value of the function call's argument. The compiler would not allow the call `f()` because the compiler cannot deduce template argument from the default arguments of `f()`.

The compiler cannot deduce a template type argument from the type of a non-type template argument. For example, the compiler will not allow the following:

```cpp
template<class T, T i> void f(int[20][i]) { }

int main() {  
  int a[20][30];  
  f(a);  
}
```

The compiler cannot deduce the type of template parameter \( T \).

**Deducing non-type template arguments**

The compiler cannot deduce the value of a major array bound unless the bound refers to a reference or pointer type. Major array bounds are not part of function parameter types. The following code demonstrates this:
template<int i> void f(int a[10][i]) { };
template<int i> void g(int a[i]) { };
template<int i> void h(int (&a)[i]) { };

int main () {
  int b[10][20];
  int c[10];
  f(b);
  // g(c);
  h(c);
}

The compiler would not allow the call g(c); the compiler cannot deduce template argument i.

The compiler cannot deduce the value of a non-type template argument used in an expression in the template function's parameter list. The following example demonstrates this:
template<int i> class X { };
template<int i> void f(X<i - 1>) { };

int main () {
  X<0> a;
  f<1>(a);
  // f(a);
}

In order to call function f() with object a, the function must accept an argument of type X<0>. However, the compiler cannot deduce that the template argument i must be equal to 1 in order for the function template argument type X<i - 1> to be equivalent to X<0>. Therefore the compiler would not allow the function call f(a).

If you want the compiler to deduce a non-type template argument, the type of the parameter must match exactly the type of value used in the function call. For example, the compiler will not allow the following:
template<int i> class A { };
template<short d> void f(A<d>) { };

int main() {
  A<1> a;
  f(a);
}

The compiler will not convert int to short when the example calls f().

However, deduced array bounds may be of any integral type.

**Overloading function templates**

You may overload a function template either by a non-template function or by another function template.

If you call the name of an overloaded function template, the compiler will try to deduce its template arguments and check its explicitly declared template arguments. If successful, it will instantiate a function template specialization, then add this specialization to the set of candidate functions used in overload resolution. The compiler proceeds with overload resolution, choosing the most appropriate function from the set of candidate functions. Non-template functions take precedence over template functions. The following example describes this:
```cpp
#include <iostream>
using namespace std;

template<class T> void f(T x, T y) { cout << "Template" << endl; }

void f(int w, int z) { cout << "Non-template" << endl; }

int main() {
    f(1, 2);
    f('a', 'b');
    f(1, 'b');
}
```

The following is the output of the above example:

Non-template
Template
Non-template

The function call f(1, 2) could match the argument types of both the template function and the non-template function. The non-template function is called because a non-template function takes precedence in overload resolution.

The function call f('a', 'b') can only match the argument types of the template function. The template function is called.

Argument deduction fails for the function call f(1, 'b'); the compiler does not generate any template function specialization and overload resolution does not take place. The non-template function resolves this function call after using the standard conversion from char to int for the function argument 'b'.

Related information
  • [“Overload resolution” on page 241](#)

**Partial ordering of function templates**

A function template specialization might be ambiguous because template argument deduction might associate the specialization with more than one of the overloaded definitions. The compiler will then choose the definition that is the most specialized. This process of selecting a function template definition is called *partial ordering*.

A template X is more specialized than a template Y if every argument list that matches the one specified by X also matches the one specified by Y, but not the other way around. The following example demonstrates partial ordering:

```cpp
template<class T> void f(T) { }
template<class T> void f(T*) { }
template<class T> void f(const T*) { }

template<class T> void g(T) { }
template<class T> void g(T&) { }

int main() {
    const int *p;
    f(p);
}
```
int q;
// g(q);
// h(q);
}

The declaration \texttt{template<class T> void f(const T*)} is more specialized than \texttt{template<class T> void f(T*)}. Therefore, the function call \texttt{f(p)} calls \texttt{template<class T> void f(const T*)}. However, neither \texttt{void g(T)} nor \texttt{void g(T&)} is more specialized than the other. Therefore, the function call \texttt{g(q)} would be ambiguous.

Ellipses do not affect partial ordering. Therefore, the function call \texttt{h(q)} would also be ambiguous.

The compiler uses partial ordering in the following cases:

- Calling a function template specialization that requires overload resolution.
- Taking the address of a function template specialization.
- When a friend function declaration, an explicit instantiation, or explicit specialization refers to a function template specialization.
- Determining the appropriate deallocation function that is also a function template for a given placement operator \texttt{new}.

Related information
- “Template specialization” on page 351
- “The new operator (C++ only)” on page 140

\section*{Template instantiation}

The act of creating a new definition of a function, class, or member of a class from a template declaration and one or more template arguments is called \textit{template instantiation}. The definition created from a template instantiation to handle a specific set of template arguments is called a \textit{specialization}.

\begin{center}
\textbf{IBM extension}
\end{center}

A forward declaration of a template instantiation has the form of an explicit template instantiation preceded by the \texttt{extern} keyword.

\begin{center}
\textbf{End of IBM extension}
\end{center}

\subsection*{Template instantiation declaration syntax}

\begin{verbatim}
extern—template—template_declaration
\end{verbatim}

The language feature is an orthogonal extension to Standard C++ for compatibility with GNU C++, and is described further in “Explicit instantiation” on page 350.

Related information
- “Template specialization” on page 351

\section*{Implicit instantiation}

Unless a template specialization has been explicitly instantiated or explicitly specialized, the compiler will generate a specialization for the template only when it needs the definition. This is called \textit{implicit instantiation}. 
If the compiler must instantiate a class template specialization and the template is declared, you must also define the template.

For example, if you declare a pointer to a class, the definition of that class is not needed and the class will not be implicitly instantiated. The following example demonstrates when the compiler instantiates a template class:

```cpp
template<class T> class X {
    public:
    X* p;
    void f();
    void g();
};
X<int>* q;
X<int> r;
X<float>* s;
r.f();
s->g();
```

The compiler requires the instantiation of the following classes and functions:
- `X<int>` when the object `r` is declared
- `X<int>::f()` at the member function call `r.f()`
- `X<float>` and `X<float>::g()` at the class member access function call `s->g()`

Therefore, the functions `X<T>::f()` and `X<T>::g()` must be defined in order for the above example to compile. (The compiler will use the default constructor of class `X` when it creates object `r`.) The compiler does not require the instantiation of the following definitions:
- class `X` when the pointer `p` is declared
- `X<int>` when the pointer `q` is declared
- `X<float>` when the pointer `s` is declared

The compiler will implicitly instantiate a class template specialization if it is involved in pointer conversion or pointer to member conversion. The following example demonstrates this:

```cpp
template<class T> class B { };
template<class T> class D : public B<T> { };
void g(B<double>** p, D<int>** q)
{
    B<double>** r = p;
    delete q;
}
```

The assignment `B<double>** r = p` converts `p` of type `B<double>**` to a type of `B<double>**`; the compiler must instantiate `B<double>`. The compiler must instantiate `D<int>` when it tries to delete `q`.

If the compiler implicitly instantiates a class template that contains static members, those static members are not implicitly instantiated. The compiler will instantiate a static member only when the compiler needs the static member’s definition. Every instantiated class template specialization has its own copy of static members. The following example demonstrates this:

```cpp
template<class T> class X {
    public:
    static T v;
};
template<class T> T X<T>::v = 0;
```
Object a has a static member variable \( v \) of type char*. Object b has a static variable \( v \) of type float. Objects b and c share the single static data member \( v \).

An implicitly instantiated template is in the same namespace where you defined the template.

If a function template or a member function template specialization is involved with overload resolution, the compiler implicitly instantiates a declaration of the specialization.

**Related information**
- "Template instantiation" on page 348

**Explicit instantiation**
You can explicitly tell the compiler when it should generate a definition from a template. This is called *explicit instantiation*.

**Explicit instantiation declaration syntax**

```
template---template_declaration---template_declaration---template_declaration---template_declaration---
```

The following are examples of explicit instantiations:

```c
template<class T> class Array { void mf(); }
template class Array<char>; // explicit instantiation
template void Array<int>::mf(); // explicit instantiation

template<class T> void sort(Array<T>& v)
{ }
template void sort(Array<char>&); // explicit instantiation

namespace N {
    template<class T> void f(T&)
    { }
}
template void N::f<int>(int&);
// The explicit instantiation is in namespace N.

int* p = 0;
template<class T> T g(T = &p);
template char g(char); // explicit instantiation

template <class T> class X {
    private:
        T v(T arg) { return arg; }
    };
template int X<int>::v(int); // explicit instantiation

template<class T> T g(T val) { return val;}
template<class T> void Array<T>::mf()
{ }
```

A template declaration must be in scope at the point of instantiation of the template’s explicit instantiation. An explicit instantiation of a template specialization is in the same namespace where you defined the template.
Access checking rules do not apply to names in explicit instantiations. Template arguments and names in a declaration of an explicit instantiation may be private types or objects. In the above example, the compiler allows the explicit instantiation template int X<int>::v(int) even though the member function is declared private.

The compiler does not use default arguments when you explicitly instantiate a template. In the above example the compiler allows the explicit instantiation template char g(char) even though the default argument is an address of type int.

---

**IBM extension**

An extern-qualified template declaration does not instantiate the class or function. For both classes and functions, the extern template instantiation prevents the instantiation of parts of the template, provided that the instantiation has not already been triggered by code prior to the extern template instantiation. For classes, the members (both static and nonstatic) are not instantiated. The class itself is instantiated if required to map the class. For functions, the prototype is instantiated, but the body of the template function is not instantiated.

The following examples show template instantiation using extern:

```cpp
template<class T>class C {
    static int i;
    void f(T) {}
};
template<class U>int C<U>::i = 0;
extern template C<int>; // extern explicit template instantiation
C<int>::i; // does not cause instantiation of C<int>::i
            // or C<int>::f(int) in this file,
            // but the class is instantiated for mapping
C<char>d; // normal instantiations

template<class C> C foo(C c) { return c; }
extern template int foo<int>(int); // extern explicit template instantiation
int i = foo(i); // does not cause instantiation of the body of foo<int>
```

---

**Template specialization**

The act of creating a new definition of a function, class, or member of a class from a template declaration and one or more template arguments is called *template instantiation*. The definition created from a template instantiation is called a *specialization*. A *primary template* is the template that is being specialized.

**Related information**

- “Template instantiation” on page 348

**Explicit specialization**

When you instantiate a template with a given set of template arguments the compiler generates a new definition based on those template arguments. You can override this behavior of definition generation. You can instead specify the definition the compiler uses for a given set of template arguments. This is called *explicit specialization*. You can explicitly specialize any of the following:

- Function or class template
• Member function of a class template
• Static data member of a class template
• Member class of a class template
• Member function template of a class template
• Member class template of a class template

Explicit specialization declaration syntax

```
<template>declaration_name[template_argument_list]declaration_body
```

The `template<>` prefix indicates that the following template declaration takes no template parameters. The `declaration_name` is the name of a previously declared template. Note that you can forward-declare an explicit specialization so the `declaration_body` is optional, at least until the specialization is referenced.

The following example demonstrates explicit specialization:

```
using namespace std;

template<class T = float, int i = 5> class A
{
    public:
    A();
    int value;
};

template<> class A<> { public: A(); };
template<> class A<double, 10> { public: A(); };

template<class T, int i> A<T, i>::A() : value(i) {
    cout << "Primary template, "
    << "non-type argument is " << value << endl;
}

A<>::A() {
    cout << "Explicit specialization "
    << "default arguments" << endl;
}

A<double, 10>::A() {
    cout << "Explicit specialization "
    << "<double, 10>" << endl;
}

int main()
{
    A<int,6> x;
    A<> y;
    A<double, 10> z;
}
```

The following is the output of the above example:

```
Primary template non-type argument is: 6
Explicit specialization default arguments
Explicit specialization <double, 10>
```

This example declared two explicit specializations for the primary template (the template which is being specialized) class A. Object x uses the constructor of the primary template. Object y uses the explicit specialization A<>::A(). Object z uses the explicit specialization A<double, 10>::A().
Related information

- "Function templates" on page 340
- "Class templates" on page 336
- "Member functions of class templates" on page 338
- "Static data members and templates" on page 338

Definition and declaration of explicit specializations

The definition of an explicitly specialized class is unrelated to the definition of the primary template. You do not have to define the primary template in order to define the specialization (nor do you have to define the specialization in order to define the primary template). For example, the compiler will allow the following:

```cpp
template<class T> class A;
template<> class A<int>;

template<> class A<int> { /* ... */ };
```

The primary template is not defined, but the explicit specialization is.

You can use the name of an explicit specialization that has been declared but not defined the same way as an incompletely defined class. The following example demonstrates this:

```cpp
template<class T> class X { };
template<> class X<char>;
X<char>* p;
X<int> i;
// X<char> j;
```

The compiler does not allow the declaration `X<char> j` because the explicit specialization of `X<char>` is not defined.

Explicit specialization and scope

A declaration of a primary template must be in scope at the point of declaration of the explicit specialization. In other words, an explicit specialization declaration must appear after the declaration of the primary template. For example, the compiler will not allow the following:

```cpp
template<> class A<int>;
template<class T> class A;
```

An explicit specialization is in the same namespace as the definition of the primary template.

Class members of explicit specializations

A member of an explicitly specialized class is not implicitly instantiated from the member declaration of the primary template. You have to explicitly define members of a class template specialization. You define members of an explicitly specialized template class as you would normal classes, without the `template<>` prefix. In addition, you can define the members of an explicit specialization inline; no special template syntax is used in this case. The following example demonstrates a class template specialization:

```cpp
template<class T> class A {
    public:
        void f(T);
};

template<> class A<int> {
    public:
        int g(int);
```
int A<int>::g(int arg) { return 0; }

int main() {
    A<int> a;
    a.g(1234);
}

The explicit specialization A<int> contains the member function g(), which the primary template does not.

If you explicitly specialize a template, a member template, or the member of a class template, then you must declare this specialization before that specialization is implicitly instantiated. For example, the compiler will not allow the following code:

template<class T> class A {};

void f() { A<int> x; }
template<> class A<int> {};

int main() { f(); }

The compiler will not allow the explicit specialization template<> class A<int> {}; because function f() uses this specialization (in the construction of x) before the specialization.

**Explicit specialization of function templates**

In a function template specialization, a template argument is optional if the compiler can deduce it from the type of the function arguments. The following example demonstrates this:

template<class T> class X {};
template<class T> void f(X<T>);
template<> void f(X<int>);

The explicit specialization template<> void f(X<int>) is equivalent to template<> void f<int>(X<int>).

You cannot specify default function arguments in a declaration or a definition for any of the following:

- Explicit specialization of a function template
- Explicit specialization of a member function template

For example, the compiler will not allow the following code:

template<class T> void f(T a) { }
template<> void f<int>({ a = 5 }) { }

template<class T> class X {
    void f(T a) { }
};
template<> void X<int>::f(int a = 10) { }

**Related information**

- "Function templates” on page 340

**Explicit specialization of members of class templates**

Each instantiated class template specialization has its own copy of any static members. You may explicitly specialize static members. The following example demonstrates this:
template<class T> class X {
    public:
        static T v;
        static void f(T);
    
    template<class T> T X<T>::v = 0;
    template<class T> void X<T>::f(T arg) { v = arg; }

    template<> char* X<char*>::v = "Hello";
    template<> void X<float>::f(float arg) { v = arg * 2; }

    int main() {
        X<char*> a, b;
        X<float> c;
        c.f(10);
    }

This code explicitly specializes the initialization of static data member X::v to point to the string "Hello" for the template argument char*. The function X::f() is explicitly specialized for the template argument float. The static data member v in objects a and b point to the same string, "Hello". The value of c.v is equal to 20 after the call function call c.f(10).

You can nest member templates within many enclosing class templates. If you explicitly specialize a template nested within several enclosing class templates, you must prefix the declaration with template<> for every enclosing class template you specialize. You may leave some enclosing class templates unspecialized, however you cannot explicitly specialize a class template unless its enclosing class templates are also explicitly specialized. The following example demonstrates explicit specialization of nested member templates:

#include <iostream>
using namespace std;

template<class T> class X {
    public:
        template<class U> class Y {
            public:
                template<class V> void f(U,V);
                void g(U);
            
        
    
    template<class T> template<class U> template<class V> void X<T>::Y<U>::f(U,V) { cout << "Template 1" << endl; }
    template<class T> template<class U> void X<T>::Y<U>::g(U) { cout << "Template 2" << endl; }
    template<> template<>
        void X<int>::Y<int>::g(int) { cout << "Template 3" << endl; }
    template<> template<class V> template<class V> void X<int>::Y<int>::f(int, V) { cout << "Template 4" << endl; }
    template<> template<>
        void X<int>::Y<int>::f<int>(int, int) { cout << "Template 5" << endl; }
    // template<> template<class U> template<class V>
    //        void X<char>::Y<U>::f(U, V) { cout << "Template 6" << endl; }
    // template<class T> template<>
    //        void X<T>::Y<float>::g(float) { cout << "Template 7" << endl; }

Chapter 15. Templates (C++ only) 355
int main() {
    X<int>::Y<int> a;
    X<char>::Y<char> b;
    a.f(1, 2);
    a.f(3, 'x');
    a.g(3);
    b.f('x', 'y');
    b.g('z');
}

The following is the output of the above program:
Template 5
Template 4
Template 3
Template 1
Template 2

• The compiler would not allow the template specialization definition that would output "Template 6" because it is attempting to specialize a member (function f()) without specialization its containing class (Y).

• The compiler would not allow the template specialization definition that would output "Template 7" because the enclosing class of class Y (which is class X) is not explicitly specialized.

A friend declaration cannot declare an explicit specialization.

Related information
• "Static data members and templates" on page 338

Partial specialization

When you instantiate a class template, the compiler creates a definition based on the template arguments you have passed. Alternatively, if all those template arguments match those of an explicit specialization, the compiler uses the definition defined by the explicit specialization.

A partial specialization is a generalization of explicit specialization. An explicit specialization only has a template argument list. A partial specialization has both a template argument list and a template parameter list. The compiler uses the partial specialization if its template argument list matches a subset of the template arguments of a template instantiation. The compiler will then generate a new definition from the partial specialization with the rest of the unmatched template arguments of the template instantiation.

You cannot partially specialize function templates.

Partial specialization syntax

```
➤---template---template_parameter_list---declaration_name-------------------------➤
➤<template_argument_list>---declaration_body-------------------------------------➤
```

The declaration_name is a name of a previously declared template. Note that you can forward-declare a partial specialization so that the declaration_body is optional.

The following demonstrates the use of partial specializations:
```
template<class T, class U, int I> struct X
{ void f() { cout << "Primary template" << endl; } };

template<class T, int I> struct X<T, T*, I>
{ void f() { cout << "Partial specialization 1" << endl; } };

template<class T, class U, int I> struct X<T*, U, I>
{ void f() { cout << "Partial specialization 2" << endl; } };

template<class T> struct X<int, T*, 10>
{ void f() { cout << "Partial specialization 3" << endl; } };

template<class T, class U, int I> struct X<T, U*, I>
{ void f() { cout << "Partial specialization 4" << endl; } };

int main()
{ X<int, int, 10> a;
  X<int, int*, 5> b;
  X<int*, float, 10> c;
  X<int, char*, 10> d;
  X<float, int*, 10> e;
  // X<int, int*, 10> f;
  a.f(); b.f(); c.f(); d.f(); e.f();
}

The following is the output of the above example:
Primary template
Partial specialization 1
Partial specialization 2
Partial specialization 3
Partial specialization 4

The compiler would not allow the declaration X<int, int*, 10> f because it can match template struct X<T, T*, I>, template struct X<int, T*, 10>, or template struct X<T, U*, I>, and none of these declarations are a better match than the others.

Each class template partial specialization is a separate template. You must provide definitions for each member of a class template partial specialization.

**Template parameter and argument lists of partial specializations**
Primary templates do not have template argument lists; this list is implied in the template parameter list.

Template parameters specified in a primary template but not used in a partial specialization are omitted from the template parameter list of the partial specialization. The order of a partial specialization's argument list is the same as the order of the primary template's implied argument list.

In a template argument list of a partial template parameter, you cannot have an expression that involves non-type arguments unless that expression is only an identifier. In the following example, the compiler will not allow the first partial specialization, but will allow the second one:

template<int I, int J> class X { };

// Invalid partial specialization
template<int I> class X <I * 4, I + 3> { };  

// Valid partial specialization  
template <int I> class X <I, I> { };  

The type of a non-type template argument cannot depend on a template parameter of a partial specialization. The compiler will not allow the following partial specialization:

```
template<class T, T i> class X { };  
```

// Invalid partial specialization  
template<class T> class X<T, 25> { };  

A partial specialization’s template argument list cannot be the same as the list implied by the primary template.

You cannot have default values in the template parameter list of a partial specialization.

Related information  
• “Template parameters” on page 330  
• “Template arguments” on page 332

Matching of class template partial specializations
The compiler determines whether to use the primary template or one of its partial specializations by matching the template arguments of the class template specialization with the template argument lists of the primary template and the partial specializations:

• If the compiler finds only one specialization, then the compiler generates a definition from that specialization.

• If the compiler finds more than one specialization, then the compiler tries to determine which of the specializations is the most specialized. A template X is more specialized than a template Y if every argument list that matches the one specified by X also matches the one specified by Y, but not the other way around. If the compiler cannot find the most specialized specialization, then the use of the class template is ambiguous; the compiler will not allow the program.

• If the compiler does not find any matches, then the compiler generates a definition from the primary template.

Name binding and dependent names

Name binding is the process of finding the declaration for each name that is explicitly or implicitly used in a template. The compiler may bind a name in the definition of a template, or it may bind a name at the instantiation of a template.

A dependent name is a name that depends on the type or the value of a template parameter. For example:

```
template<class T> class U : A<T>  
{  
    typename T::B x;  
    void f(A<T>& y)  
    {  
        *y++;  
    }  
};  
```
The dependent names in this example are the base class `A<T>`, the type name `T::B`, and the variable `y`.

The compiler binds dependent names when a template is instantiated. The compiler binds non-dependent names when a template is defined. For example:

```cpp
void f(double) { cout << "Function f(double)" << endl; }

template<class T> void g(T a) {
    f(123);
    h(a);
}

void f(int) { cout << "Function f(int)" << endl; }
void h(double) { cout << "Function h(double)" << endl; }

void i() {
    extern void h(int);
    g<int>(234);
}

void h(int) { cout << "Function h(int)" << endl; }
```

The following is the output if you call function `i()`:

```
Function f(double)
Function h(double)
```

The point of definition of a template is located immediately before its definition. In this example, the point of definition of the function template `g(T)` is located immediately before the keyword `template`. Because the function call `f(123)` does not depend on a template argument, the compiler will consider names declared before the definition of function template `g(T)`. Therefore, the call `f(123)` will call `f(double)`. Although `f(int)` is a better match, it is not in scope at the point of definition of `g(T)`.

The point of instantiation of a template is located immediately before the declaration that encloses its use. In this example, the point of instantiation of the call to `g<int>(234)` is located immediately before `i()`. Because the function call `h(a)` depends on a template argument, the compiler will consider names declared before the instantiation of function template `g(T)`. Therefore, the call `h(a)` will call `h(double)`. It will not consider `h(int)`, because this function was not in scope at the point of instantiation of `g<int>(234)`.

Point of instantiation binding implies the following:

- A template parameter cannot depend on any local name or class member.
- An unqualified name in a template cannot depend on a local name or class member.

**Related information**

- [“Template instantiation” on page 348](#)

---

**The typename keyword**

Use the keyword `typename` if you have a qualified name that refers to a type and depends on a template parameter. Only use the keyword `typename` in template declarations and definitions. The following example illustrates the use of the keyword `typename`:

```cpp
void f(T::B a) {
    T x;
    a;
}

template<class T> void g(T a) {
    T::B b;
    a;
}
```

The `typename` keyword is used in the declaration of `b` in the template declaration of `g(T)` to indicate that `b` is a type name.

---

Chapter 15. Templates (C++ only)  359
template<class T> class A
{
    T::x(y);
    typedef char C;
    A::C d;
}

The statement T::x(y) is ambiguous. It could be a call to function x() with a nonlocal argument y, or it could be a declaration of variable y with type T::x. C++ will interpret this statement as a function call. In order for the compiler to interpret this statement as a declaration, you would add the keyword typename to the beginning of it. The statement A::C d; is ill-formed. The class A also refers to A<T> and thus depends on a template parameter. You must add the keyword typename to the beginning of this declaration:

typename A::C d;

You can also use the keyword typename in place of the keyword class in template parameter declarations.

**Related information**

- "Template parameters" on page 330

---

**The template keyword as qualifier**

Use the keyword template as a qualifier to distinguish member templates from other names. The following example illustrates when you must use template as a qualifier:

```cpp
class A
{
    public:
        template<class T> T function_m() { }
    }

template<class U> void function_n(U argument)
{    
    char object_x = argument.function_m<char>();
}
```

The declaration char object_x = argument.function_m<char>(); is ill-formed. The compiler assumes that the < is a less-than operator. In order for the compiler to recognize the function template call, you must add the template quantifier:

```
char object_x = argument.template function_m<char>();
```

If the name of a member template specialization appears after a ., ->, or :: operator, and that name has explicitly qualified template parameters, prefix the member template name with the keyword template. The following example demonstrates this use of the keyword template:

```cpp
#include <iostream>
using namespace std;

class X
{
    public:
        template <int j> struct S
        {
            void h() {    
                cout << "member template's member function: " << j << endl;
            }
        }
        template <int i> void f() {
            cout << "Primary: " << i << endl;
        }
```
template<> void X::f<20>() {
    cout << "Specialized, non-type argument = 20" << endl;
}

template<class T> void g(T* p) {
    p->template f<100>();
    p->template f<20>();
    typename T::template S<40> s; // use of scope operator on a member template
    s.h();
}

int main()
{
    X temp;
    g(&temp);
}

The following is the output of the above example:
Primary: 100
Specialized, non-type argument = 20
member template's member function: 40

If you do not use the keyword template in these cases, the compiler will interpret the < as a less-than operator. For example, the following line of code is ill-formed:
p->f<100>();

The compiler interprets f as a non-template member, and the < as a less-than operator.
Chapter 16. Exception handling (C++ only)

Exception handling is a mechanism that separates code that detects and handles exceptional circumstances from the rest of your program. Note that an exceptional circumstance is not necessarily an error.

When a function detects an exceptional situation, you represent this with an object. This object is called an exception object. In order to deal with the exceptional situation you throw the exception. This passes control, as well as the exception, to a designated block of code in a direct or indirect caller of the function that threw the exception. This block of code is called a handler. In a handler, you specify the types of exceptions that it may process. The C++ runtime, together with the generated code, will pass control to the first appropriate handler that is able to process the exception thrown. When this happens, an exception is caught. A handler may rethrow an exception so it can be caught by another handler.

The exception handling mechanism is made up of the following elements:

- **try blocks**
- **catch blocks**
- **throw expressions**
- **Exception specifications**

### try blocks

You use a try block to indicate which areas in your program that might throw exceptions you want to handle immediately. You use a function try block to indicate that you want to detect exceptions in the entire body of a function.

**try block syntax**

```
try-{statements}->handler
```

**Function try block syntax**

```
try::{member_initializer_list}function_body->handler
```

The following is an example of a function try block with a member initializer, a function try block and a try block:

```cpp
#include <iostream>
using namespace std;

class E
{
    public:
        const char* error;
        E(const char* arg) : error(arg) { }
};
```
class A {
public:
    int i;
    // A function try block with a member
    // initializer
    A() try : i(0) {
        throw E("Exception thrown in A()");
    }
    catch (E& e) {
        cout << e.error << endl;
    }
};

// A function try block
void f() try {
    throw E("Exception thrown in f()");
} catch (E& e) {
    cout << e.error << endl;
}

void g() {
    throw E("Exception thrown in g()");
}

int main() {
    f();
    // A try block
    try {
        g();
    } catch (E& e) {
        cout << e.error << endl;
    } try {
        A x;
    } catch(...) { }
}

The following is the output of the above example:
Exception thrown in f()
Exception thrown in g()
Exception thrown in A()

The constructor of class A has a function try block with a member initializer. Function f() has a function try block. The main() function contains a try block.

Related information
- "Initializing base classes and members" on page 308

**Nested try blocks**

When try blocks are nested and a throw occurs in a function called by an inner try block, control is transferred outward through the nested try blocks until the first catch block is found whose argument matches the argument of the throw expression.
For example:

```c
try
{
    func1();
    try
    {
        func2();
    }
    catch (spec_err) { /* ... */ }
    func3();
}
catch (type_err) { /* ... */ }
// if no throw is issued, control resumes here.
```

In the above example, if `spec_err` is thrown within the inner try block (in this case, from `func2()`), the exception is caught by the inner catch block, and, assuming this catch block does not transfer control, `func3()` is called. If `spec_err` is thrown after the inner try block (for instance, by `func3()`), it is not caught and the function `terminate()` is called. If the exception thrown from `func2()` in the inner try block is `type_err`, the program skips out of both try blocks to the second catch block without invoking `func3()`, because no appropriate catch block exists following the inner try block.

You can also nest a try block within a catch block.

---

**catch blocks**

**catch block syntax**

```c
catch(---exception_declaration---){---statements---}
```

You can declare a handler to catch many types of exceptions. The allowable objects that a function can catch are declared in the parentheses following the catch keyword (the `exception_declaration`). You can catch objects of the fundamental types, base and derived class objects, references, and pointers to all of these types. You can also catch `const` and `volatile` types. The `exception_declaration` cannot be an incomplete type, or a reference or pointer to an incomplete type other than one of the following:

- `void*`
- `const void*`
- `volatile void*`
- `const volatile void*`

You cannot define a type in an `exception_declaration`.

You can also use the `catch(...)` form of the handler to catch all thrown exceptions that have not been caught by a previous catch block. The ellipsis in the catch argument indicates that any exception thrown can be handled by this handler.

If an exception is caught by a `catch(...)` block, there is no direct way to access the object thrown. Information about an exception caught by `catch(...)` is very limited.

You can declare an optional variable name if you want to access the thrown object in the catch block.
A catch block can only catch accessible objects. The object caught must have an accessible copy constructor.

Related information
- “Type qualifiers” on page 68
- “Member access” on page 269

Function try block handlers
The scope and lifetime of the parameters of a function or constructor extend into the handlers of a function try block. The following example demonstrates this:

```c++
void f(int &x) try {
    throw 10;
} catch (const int &i) {
    x = i;
}

int main() {
    int v = 0;
    f(v);
}
```

The value of v after f() is called is 10.

A function try block on main() does not catch exceptions thrown in destructors of objects with static storage duration, or constructors of namespace scope objects.

The following example throws an exception from a destructor of a static object:

```c++
#include <iostream>
using namespace std;

class E {
public:
    const char* error;
    E(const char* arg) : error(arg) { }
};

class A {
public: "A() { throw E("Exception in "A()); }
};

class B {
public: "B() { throw E("Exception in "B()); }
};

int main() try {
    cout << "In main" << endl;
    static A cow;
    B bull;
} catch (E& e) {
    cout << e.error << endl;
}
```

The following is the output of the above example:

```
In main
Exception in "B()
```
The run time will not catch the exception thrown when object cow is destroyed at the end of the program.

The following example throws an exception from a constructor of a namespace scope object:

```cpp
#include <iostream>
using namespace std;

class E {
public:
    const char* error;
    E(const char* arg) : error(arg) { }
};

namespace N {
    class C {
public:
        C() {
            cout << "In C()" << endl;
            throw E("Exception in C()");
        }
    };

    C calf;
};

int main() try {
    cout << "In main" << endl;
} catch (E& e) {
    cout << e.error << endl;
}
```

The following is the output of the above example:

```
In C()
The compiler will not catch the exception thrown when object calf is created.
```

In a function try block’s handler, you cannot have a jump into the body of a constructor or destructor.

A return statement cannot appear in a function try block’s handler of a constructor.

When the function try block’s handler of an object’s constructor or destructor is entered, fully constructed base classes and members of that object are destroyed. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

class E {
public:
    const char* error;
    E(const char* arg) : error(arg) { }
};

class B {
public:
    B() { }
    "B()" { cout << "B() called" << endl; }
};

class D : public B {
public:
```
D();
~D() { cout << "D() called" << endl; }
;
D::D() try : B() {
    throw E("Exception in D()");
} catch(E& e) {
    cout << "Handler of function try block of D(): " << e.error << endl;
};

int main() {
    try {
        D val;
    } catch(...) {
    }
}

The following is the output of the above example:
"B() called
Handler of function try block of D(): Exception in D()

When the function try block’s handler of D() is entered, the run time first calls the destructor of the base class of D, which is B. The destructor of D is not called because val is not fully constructed.

The run time will rethrow an exception at the end of a function try block’s handler of a constructor or destructor. All other functions will return once they have reached the end of their function try block’s handler. The following example demonstrates this:

#include <iostream>
using namespace std;

class E {
    public:
        const char* error;
        E(const char* arg) : error(arg) {
    }
};

class A {
    public:
        A() try { throw E("Exception in A()"); } catch(E& e) {
            cout << "Handler in A(): " << e.error << endl;
        }
    }

int f() try {
    throw E("Exception in f()");
    return 0;
} catch(E& e) {
    cout << "Handler in f(): " << e.error << endl;
    return 1;
}

int main() {
    try {
        A cow;
    } catch(E& e) {
        cout << "Handler in main(): " << e.error << endl;
    }

    try {
        i = f();
    } catch(E& e) {
        cout << "Another handler in main(): " << e.error << endl;
    }
}
cout << "Returned value of f(): " << i << endl;

The following is the output of the above example:
Handler in A(): Exception in A()
Handler in main(): Exception in A()
Handler in f(): Exception in f()
Returned value of f(): 1

Related information
- “The main() function” on page 212
- “The static storage class specifier” on page 44
- Chapter 9, “Namespaces (C++ only),” on page 221
- “Destructors” on page 312

Arguments of catch blocks
If you specify a class type for the argument of a catch block (the
exception_declaration), the compiler uses a copy constructor to initialize that
argument. If that argument does not have a name, the compiler initializes a
temporary object and destroys it when the handler exits.

The ISO C++ specifications do not require the compiler to construct temporary
objects in cases where they are redundant. The compiler takes advantage of this
rule to create more efficient, optimized code. Take this into consideration when
debugging your programs, especially for memory problems.

Related information
- “Temporary objects” on page 320

Matching between exceptions thrown and caught
An argument in the catch argument of a handler matches an argument in the
assignment_expression of the throw expression (throw argument) if any of the
following conditions is met:

- The catch argument type matches the type of the thrown object.
- The catch argument is a public base class of the thrown class object.
- The catch specifies a pointer type, and the thrown object is a pointer type that
can be converted to the pointer type of the catch argument by standard pointer
conversion.

Note: If the type of the thrown object is const or volatile, the catch argument
must also be a const or volatile for a match to occur. However, a const,
volatile, or reference type catch argument can match a nonconstant,
nonvolatile, or nonreference object type. A nonreference catch argument type
matches a reference to an object of the same type.

Related information
- “Pointer conversions” on page 110
- “Type qualifiers” on page 68
- “References (C++ only)” on page 88
Order of catching

If the compiler encounters an exception in a try block, it will try each handler in order of appearance.

If a catch block for objects of a base class precedes a catch block for objects of a class derived from that base class, the compiler issues a warning and continues to compile the program despite the unreachable code in the derived class handler.

A catch block of the form catch(...) must be the last catch block following a try block or an error occurs. This placement ensures that the catch(...) block does not prevent more specific catch blocks from catching exceptions intended for them.

If the run time cannot find a matching handler in the current scope, the run time will continue to find a matching handler in a dynamically surrounding try block. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

class E {
public:
    const char* error;
    E(const char* arg) : error(arg) {};
};

class F : public E {
public:
    F(const char* arg) : E(arg) {};
};

void f() {
    try {
        cout << "In try block of f()" << endl;
        throw E("Class E exception");
    }
    catch (F& e) {
        cout << "In handler of f()";
        cout << e.error << endl;
    }
};

int main() {
    try {
        cout << "In main" << endl;
        f();
    }
    catch (E& e) {
        cout << "In handler of main: ";
        cout << e.error << endl;
    }
    cout << "Resume execution in main" << endl;
}
```

The following is the output of the above example:

```
In main
In try block of f()
In handler of main: Class E exception
Resume execution in main
```

In function f(), the run time could not find a handler to handle the exception of type E thrown. The run time finds a matching handler in a dynamically surrounding try block: the try block in the main() function.
If the run time cannot find a matching handler in the program, it calls the terminate() function.

**Related information**
- "try blocks" on page 363

### throw expressions

You use a *throw expression* to indicate that your program has encountered an exception.

#### throw expression syntax

\[
\text{throw } \text{assignment_expression}
\]

The type of *assignment_expression* cannot be an incomplete type, or a pointer to an incomplete type other than one of the following:
- `void*`
- `const void*`
- `volatile void*`
- `const volatile void*`

The *assignment_expression* is treated the same way as a function argument in a call or the operand of a return statement.

If the *assignment_expression* is a class object, the copy constructor and destructor of that object must be accessible. For example, you cannot throw a class object that has its copy constructor declared as private.

**Related information**
- "Incomplete types" on page 40

### Rethrowing an exception

If a catch block cannot handle the particular exception it has caught, you can rethrow the exception. The rethrow expression (*throw* without *assignment_expression*) causes the originally thrown object to be rethrown.

Because the exception has already been caught at the scope in which the rethrow expression occurs, it is rethrown out to the next dynamically enclosing try block. Therefore, it cannot be handled by catch blocks at the scope in which the rethrow expression occurred. Any catch blocks for the dynamically enclosing try block have an opportunity to catch the exception.

The following example demonstrates rethrowing an exception:

```cpp
#include <iostream>
using namespace std;

struct E {  
    const char* message;  
    E() : message("Class E") { }  
};

struct E1 : E {  
    const char* message;  
    E1() : message("Class E1") { }  
}
```

```cpp
Chapter 16. Exception handling (C++ only) 371
```
struct E2 : E {
    const char* message;
    E2() : message("Class E2") { }
};

void f() {
    try {
        cout << "In try block of f()" << endl;
        cout << "Throwing exception of type E1" << endl;
        E1 myException;
        throw myException;
    } catch (E2& e) {
        cout << "In handler of f(), catch (E2& e)" << endl;
        cout << "Exception: " << e.message << endl;
        throw;
    } catch (E1& e) {
        cout << "In handler of f(), catch (E1& e)" << endl;
        cout << "Exception: " << e.message << endl;
        throw;
    } catch (E& e) {
        cout << "In handler of f(), catch (E& e)" << endl;
        cout << "Exception: " << e.message << endl;
        throw;
    }
}

int main() {
    try {
        cout << "In try block of main()" << endl;
        f();
    } catch (E2& e) {
        cout << "In handler of main(), catch (E2& e)" << endl;
        cout << "Exception: " << e.message << endl;
    } catch (...) {
        cout << "In handler of main(), catch (...)" << endl;
    }
}

The following is the output of the above example:
In try block of main()
In try block of f()
Throwing exception of type E1
In handler of f(), catch (E1& e)
Exception: Class E1
In handler of main(), catch (...)

The try block in the main() function calls function f(). The try block in function f() throws an object of type E1 named myException. The handler catch (E1 & e) catches myException. The handler then rethrows myException with the statement throw to the next dynamically enclosing try block: the try block in the main() function. The handler catch(...) catches myException.
Stack unwinding

When an exception is thrown and control passes from a try block to a handler, the C++ run time calls destructors for all automatic objects constructed since the beginning of the try block. This process is called stack unwinding. The automatic objects are destroyed in reverse order of their construction. (Automatic objects are local objects that have been declared auto or register, or not declared static or extern. An automatic object x is deleted whenever the program exits the block in which x is declared.)

If an exception is thrown during construction of an object consisting of subobjects or array elements, destructors are only called for those subobjects or array elements successfully constructed before the exception was thrown. A destructor for a local static object will only be called if the object was successfully constructed.

If during stack unwinding a destructor throws an exception and that exception is not handled, the terminate() function is called. The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct E {
    const char* message;
    E(const char* arg) : message(arg) { }
};

void my_terminate() {
    cout << "Call to my_terminate" << endl;
}

struct A {
    A() { cout << "In constructor of A" << endl; }
    ~A() {
        cout << "In destructor of A" << endl;
        throw E("Exception thrown in ~A()");
    }
};

struct B {
    B() { cout << "In constructor of B" << endl; }
    ~B() { cout << "In destructor of B" << endl; }
};

int main() {
    set_terminate(my_terminate);

    try {
        cout << "In try block" << endl;
        A a;
        B b;
        throw("Exception thrown in try block of main()");
    } catch (const char* e) {
        cout << "Exception: " << e << endl;
    } catch (...) {
        cout << "Some exception caught in main()" << endl;
    }

    cout << "Resume execution of main()" << endl;
}
```

Chapter 16. Exception handling (C++ only) 373
The following is the output of the above example:

In try block
In constructor of A
In constructor of B
In destructor of B
In destructor of A
Call to my_terminate

In the try block, two automatic objects are created: a and b. The try block throws an exception of type const char*. The handler catch (const char* e) catches this exception. The C++ run time unwinds the stack, calling the destructors for a and b in reverse order of their construction. The destructor for a throws an exception. Since there is no handler in the program that can handle this exception, the C++ run time calls terminate(). (The function terminate() calls the function specified as the argument to set_terminate(). In this example, terminate() has been specified to call my_terminate().)

**Exception specifications**

C++ provides a mechanism to ensure that a given function is limited to throwing only a specified list of exceptions. An exception specification at the beginning of any function acts as a guarantee to the function’s caller that the function will throw only the exceptions contained in the exception specification.

For example, a function:

```c
void translate() throw(unknown_word, bad_grammar) { /* ... */ }
```

explicitly states that it will only throw exception objects whose types are `unknown_word` or `bad_grammar`, or any type derived from `unknown_word` or `bad_grammar`.

**Exception specification syntax**

```c
throw(type_id_list)
```

The `type_id_list` is a comma-separated list of types. In this list you cannot specify an incomplete type, a pointer or a reference to an incomplete type, other than a pointer to `void`, optionally qualified with `const` and/or `volatile`. You cannot define a type in an exception specification.

A function with no exception specification allows all exceptions. A function with an exception specification that has an empty `type_id_list`, `throw()`, does not allow any exceptions to be thrown.

An exception specification is not part of a function’s type.

An exception specification may only appear at the end of a function declarator of a function, pointer to function, reference to function, pointer to member function declaration, or pointer to member function definition. An exception specification cannot appear in a `typedef` declaration. The following declarations demonstrate this:

```c
void f() throw(int);
void (*g)() throw(int);
void h(void i() throw(int));
// typedef int (*j)() throw(int); This is an error.
```
The compiler would not allow the last declaration, typedef int (*j)() throw(int).

Suppose that class A is one of the types in the type_id_list of an exception specification of a function. That function may throw exception objects of class A, or any class publicly derived from class A. The following example demonstrates this:

```cpp
class A { }
class B : public A { }
class C { }

void f(int i) throw (A) {
    switch (i) {
    case 0: throw A();
    case 1: throw B();
    default: throw C();
    }
}

void g(int i) throw (A*) {
    A* a = new A();
    B* b = new B();
    C* c = new C();
    switch (i) {
    case 0: throw a;
    case 1: throw b;
    default: throw c;
    }
}
```

Function f() can throw objects of types A or B. If the function tries to throw an object of type C, the compiler will call unexpected() because type C has not been specified in the function’s exception specification, nor does it derive publicly from A. Similarly, function g() cannot throw pointers to objects of type C; the function may throw pointers of type A or pointers of objects that derive publicly from A.

A function that overrides a virtual function can only throw exceptions specified by the virtual function. The following example demonstrates this:

```cpp
class A {
    public:
        virtual void f() throw (int, char);  
};

class B : public A{
    public: void f() throw (int) { }
};

/* The following is not allowed. */
/*
class C : public A {  
    public: void f() { }
};

class D : public A {
    public: void f() throw (int, char, double) { }
}*/
```

The compiler allows B::f() because the member function may throw only exceptions of type int. The compiler would not allow C::f() because the member function may throw any kind of exception. The compiler would not allow D::f() because the member function can throw more types of exceptions (int, char, and double) than A::f().
Suppose that you assign or initialize a pointer to function named x with a function or pointer to function named y. The pointer to function x can only throw exceptions specified by the exception specifications of y. The following example demonstrates this:

```c
void (*f)();
void (*g)();
void (*h)() throw (int);

void i() {
    f = h;
    // h = g; This is an error.
}
```

The compiler allows the assignment `f = h` because `f` can throw any kind of exception. The compiler would not allow the assignment `h = g` because `h` can only throw objects of type `int`, while `g` can throw any kind of exception.

Implicitly declared special member functions (default constructors, copy constructors, destructors, and copy assignment operators) have exception specifications. An implicitly declared special member function will have in its exception specification the types declared in the functions’ exception specifications that the special function invokes. If any function that a special function invokes allows all exceptions, then that special function allows all exceptions. If all the functions that a special function invokes allow no exceptions, then that special function will allow no exceptions. The following example demonstrates this:

```c
class A {
    public:
    A() throw (int);
    A(const A&) throw (float);
    "A"() throw();
};

class B {
    public:
    B() throw (char);
    B(const A&);
    "B"() throw();
};

class C : public B, public A {
};
```

The following special functions in the above example have been implicitly declared:

```c
C::C() throw (int, char);
C::C(const C&); // Can throw any type of exception, including float
C::"C"() throw();
```

The default constructor of `C` can throw exceptions of type `int` or `char`. The copy constructor of `C` can throw any kind of exception. The destructor of `C` cannot throw any exceptions.

**Related information**

- “Incomplete types” on page 40
- “Function declarations and definitions” on page 191
- “Pointers to functions” on page 219
- Chapter 14, “Special member functions (C++ only),” on page 303
Special exception handling functions

Not all thrown errors can be caught and successfully dealt with by a catch block. In some situations, the best way to handle an exception is to terminate the program. Two special library functions are implemented in C++ to process exceptions not properly handled by catch blocks or exceptions thrown outside of a valid try block. These functions are:

- The unexpected() function
- The terminate() function

The unexpected() function

When a function with an exception specification throws an exception that is not listed in its exception specification, the C++ run time does the following:

1. The unexpected() function is called.
2. The unexpected() function calls the function pointed to by unexpected_handler. By default, unexpected_handler points to the function terminate().

You can replace the default value of unexpected_handler with the function set_unexpected().

Although unexpected() cannot return, it may throw (or rethrow) an exception. Suppose the exception specification of a function f() has been violated. If unexpected() throws an exception allowed by the exception specification of f(), then the C++ run time will search for another handler at the call of f(). The following example demonstrates this:

```cpp
#include <iostream>
using namespace std;

struct E {
    const char* message;
    E(const char* arg) : message(arg) { }
};

void my_unexpected() {
    cout << "Call to my_unexpected" << endl;
    throw E("Exception thrown from my_unexpected");
}

void f() throw(E) {
    cout << "In function f(), throw const char* object" << endl;
    throw("Exception, type const char*, thrown from f()");
}

int main() {
    set_unexpected(my_unexpected);
    try{
        f();
    } catch (E& e) {
        cout << "Exception in main(): " << e.message << endl;
    }
}
```

The following is the output of the above example:

In function f(), throw const char* object
Call to my_unexpected
Exception in main(): Exception thrown from my_unexpected
The `main()` function’s try block calls function `f()`. Function `f()` throws an object of type `const char*`. However the exception specification of `f()` allows only objects of type `E` to be thrown. The function `unexpected()` is called. The function `unexpected()` calls `my_unexpected()`. The function `my_unexpected()` throws an object of type `E`. Since `unexpected()` throws an object allowed by the exception specification of `f()`, the handler in the `main()` function may handle the exception.

If `unexpected()` did not throw (or rethrow) an object allowed by the exception specification of `f()`, then the C++ run time does one of two things:

- If the exception specification of `f()` included the class `std::bad_exception`, `unexpected()` will throw an object of type `std::bad_exception`, and the C++ run time will search for another handler at the call of `f()`.
- If the exception specification of `f()` did not include the class `std::bad_exception`, the function `terminate()` is called.

**Related information**

- “Special exception handling functions” on page 377
- “The `set_unexpected()` and `set_terminate()` functions” on page 379

**The `terminate()` function**

In some cases, the exception handling mechanism fails and a call to `void terminate()` is made. This `terminate()` call occurs in any of the following situations:

- The exception handling mechanism cannot find a handler for a thrown exception. The following are more specific cases of this:
  - During stack unwinding, a destructor throws an exception and that exception is not handled.
  - The expression that is thrown also throws an exception, and that exception is not handled.
  - The constructor or destructor of a nonlocal static object throws an exception, and the exception is not handled.
  - A function registered with `atexit()` throws an exception, and the exception is not handled. The following demonstrates this:

```c
extern "C" printf(char* ...);
#include <exception>
#include <cstdlib>
using namespace std;

void f() {
    printf("Function f()\n");
    throw "Exception thrown from f()";
}

void g() { printf("Function g()\n"); }
void h() { printf("Function h()\n"); }

void my_terminate() {
    printf("Call to my_terminate\n");
    abort();
}

int main() {
    set_terminate(my_terminate);
    atexit(f);
    atexit(g);
    atexit(h);
    printf("In main\n");
}```
The following is the output of the above example:

In main
Function h()
Function g()
Function f()
Call to my_terminate

To register a function with atexit(), you pass a parameter to atexit() a pointer to the function you want to register. At normal program termination, atexit() calls the functions you have registered with no arguments in reverse order. The atexit() function is in the <cstdlib> library.

- A throw expression without an operand tries to rethrow an exception, and no exception is presently being handled.
- A function f() throws an exception that violates its exception specification. The unexpected() function then throws an exception which violates the exception specification of f(), and the exception specification of f() did not include the class std::bad_exception.
- The default value of unexpected_handler is called.

The terminate() function calls the function pointed to by terminate_handler. By default, terminate_handler points to the function abort(), which exits from the program. You can replace the default value of terminate_handler with the function set_terminate().

A terminate function cannot return to its caller, either by using return or by throwing an exception.

Related information
- “The set_unexpected() and set_terminate() functions”

The set_unexpected() and set_terminate() functions

The function unexpected(), when invoked, calls the function most recently supplied as an argument to set_unexpected(). If set_unexpected() has not yet been called, unexpected() calls terminate().

The function terminate(), when invoked, calls the function most recently supplied as an argument to set_terminate(). If set_terminate() has not yet been called, terminate() calls abort(), which ends the program.

You can use set_unexpected() and set_terminate() to register functions you define to be called by unexpected() and terminate(). The functions set_unexpected() and set_terminate() are included in the standard header files. Each of these functions has as its return type and its argument type a pointer to function with a void return type and no arguments. The pointer to function you supply as the argument becomes the function called by the corresponding special function: the argument to set_unexpected() becomes the function called by unexpected(), and the argument to set_terminate() becomes the function called by terminate().

Both set_unexpected() and set_terminate() return a pointer to the function that was previously called by their respective special functions (unexpected() and terminate()). By saving the return values, you can restore the original special functions later so that unexpected() and terminate() will once again call terminate() and abort.
If you use `set_terminate()` to register your own function, the function should no return to its caller but terminate execution of the program.

**Example using the exception handling functions**

The following example shows the flow of control and special functions used in exception handling:

```c++
#include <iostream>
#include <exception>
using namespace std;

class X {}
class Y {}
class A {}

// pfv type is pointer to function returning void
typedef void (*pfv)();

void my_terminate() {
    cout << "Call to my terminate" << endl;
    abort();
}

void my_unexpected() {
    cout << "Call to my_unexpected()" << endl;
    throw;
}

void f() throw(X, Y, bad_exception) {
    throw A();
}

void g() throw(X, Y) {
    throw A();
}

int main() {
    pfv old_term = set_terminate(my_terminate);
    pfv old_unex = set_unexpected(my_unexpected);
    try {
        cout << "In first try block" << endl;
        f();
    } catch(X) {
        cout << "Caught X" << endl;
    } catch(Y) {
        cout << "Caught Y" << endl;
    } catch(bad_exception& e1) {
        cout << "Caught bad_exception" << endl;
    } catch(...) {
        cout << "Caught some exception" << endl;
    }
    cout << endl;
    try {
        cout << "In second try block" << endl;
        g();
    } catch(X) {
        cout << "Caught X" << endl;
    }
    return 0;
}
```

---

380 XL C/C++ Language Reference
catch(Y) {
    cout << "Caught Y" << endl;
}
catch (bad_exception& e2) {
    cout << "Caught bad_exception" << endl;
}
catch (...) {
    cout << "Caught some exception" << endl;
}
}

The following is the output of the above example:
In first try block
Call to my_unexpected()
Caught bad_exception

In second try block
Call to my_unexpected()
Call to my_terminate

At run time, this program behaves as follows:
1. The call to set_terminate() assigns to old_term the address of the function last passed to set_terminate() when set_terminate() was previously called.
2. The call to set_unexpected() assigns to old_unex the address of the function last passed to set_unexpected() when set_unexpected() was previously called.
3. Within the first try block, function f() is called. Because f() throws an unexpected exception, a call to unexpected() is made. unexpected() in turn calls my_unexpected(), which prints a message to standard output. The function my_unexpected() tries to rethrow the exception of type A. Because class A has not been specified in the exception specification of function f(), my_unexpected() throws an exception of type bad_exception.
4. Because bad_exception has been specified in the exception specification of function f(), the handler catch(bad_exception& e1) is able to handle the exception.
5. Within the second try block, function g() is called. Because g() throws an unexpected exception, a call to unexpected() is made. The unexpected() throws an exception of type bad_exception. Because bad_exception has not been specified in the exception specification of g(), unexpected() calls terminate(), which calls the function my_terminate().
6. my_terminate() displays a message then calls abort(), which terminates the program.

Note that the catch blocks following the second try block are not entered, because the exception was handled by my_unexpected() as an unexpected throw, not as a valid exception.
Chapter 17. Preprocessor Directives

The preprocessor is a program that is invoked by the compiler to process code before compilation. Commands for that program, known as directives, are lines of the source file beginning with the character #, which distinguishes them from lines of source program text. The effect of each preprocessor directive is a change to the text of the source code, and the result is a new source code file, which does not contain the directives. The preprocessed source code, an intermediate file, must be a valid C or C++ program, because it becomes the input to the compiler.

Preprocessor directives consist of the following:

- **Macro definition directives**, which replace tokens in the current file with specified replacement tokens
- **File inclusion directives**, which imbed files within the current file
- **Conditional compilation directives**, which conditionally compile sections of the current file
- **Message generation directives**, which control the generation of diagnostic messages
- **Assertion directives**, which specify attributes of the system the program is to run on
- **The null directive (#)**, which performs no action
- **Pragma directives**, which apply compiler-specific rules to specified sections of code

Preprocessor directives begin with the # token followed by a preprocessor keyword. The # token must appear as the first character that is not white space on a line. The # is not part of the directive name and can be separated from the name with white spaces.

A preprocessor directive ends at the new-line character unless the last character of the line is the \ (backslash) character. If the \ character appears as the last character in the preprocessor line, the preprocessor interprets the \ and the new-line character as a continuation marker. The preprocessor deletes the \ (and the following new-line character) and splices the physical source lines into continuous logical lines. White space is allowed between backslash and the end of line character or the physical end of record. However, this white space is usually not visible during editing.

Except for some #pragma directives, preprocessor directives can appear anywhere in a program.

### Macro definition directives

Macro definition directives include the following directives and operators:

- **The #define directive**, which defines a macro
- **The #undef directive**, which removes a macro definition

“Standard predefined macro names” on page 390 describes the macros that are predefined by the ISO C standard. Macros that are predefined for XL C/C++ are described in “Predefined macros” in the XL C/C++ Compiler Reference
The **define directive**

A *preprocessor define directive* directs the preprocessor to replace all subsequent occurrences of a macro with specified replacement tokens.

### define directive syntax

```
#define identifier (identifier character)
```

The `define` directive can contain:

- **Object-like macros**
- **Function-like macros**

The following are some differences between `define` and the `const` type qualifier:

- The `define` directive can be used to create a name for a numerical, character, or string constant, whereas a `const` object of any type can be declared.
- A `const` object is subject to the scoping rules for variables, whereas a constant created using `define` is not.
- Unlike a `const` object, the value of a macro does not appear in the intermediate source code used by the compiler because they are expanded inline. The inline expansion makes the macro value unavailable to the debugger.
- A macro can be used in a constant expression, such as an array bound, whereas a `const` object cannot.
- The compiler does not type-check a macro, including macro arguments.

**Related information**

- [“The const type qualifier” on page 69](#)

### Object-like macros

An *object-like macro definition* replaces a single identifier with the specified replacement tokens. The following object-like definition causes the preprocessor to replace all subsequent instances of the identifier `COUNT` with the constant 1000:

```
#define COUNT 1000
```

If the statement

```
int arry[COUNT];
```

appears after this definition and in the same file as the definition, the preprocessor would change the statement to

```
int arry[1000];
```

in the output of the preprocessor.

Other definitions can make reference to the identifier `COUNT`:

```
#define MAX_COUNT COUNT + 100
```
The preprocessor replaces each subsequent occurrence of \texttt{MAX\_COUNT} with 
\texttt{COUNT + 100}, which the preprocessor then replaces with \texttt{1000 + 100}.

If a number that is partially built by a macro expansion is produced, the 
preprocessor does not consider the result to be a single value. For example, the 
following will not result in the value 10.2 but in a syntax error.
\begin{verbatim}
#define a 10
a.2
\end{verbatim}

Identifiers that are partially built from a macro expansion may not be produced. 
Therefore, the following example contains two identifiers and results in a syntax 
error:
\begin{verbatim}
#define d efg
abcd
\end{verbatim}

\textbf{Function-like macros}

More complex than object-like macros, a function-like macro definition declares the 
names of formal parameters within parentheses, separated by commas. An empty 
formal parameter list is legal: such a macro can be used to simulate a function that 
takes no arguments. C99 adds support for function-like macros with a variable 
number of arguments. XL C++ supports function-like macros with a variable 
number of arguments, as a language extension for compatibility with C.

\textbf{Function-like macro definition:}
An identifier followed by a parameter list in parentheses and the 
replacement tokens. The parameters are imbedded in the replacement code. 
White space cannot separate the identifier (which is the name of the 
macro) and the left parenthesis of the parameter list. A comma must 
separate each parameter.

For portability, you should not have more than 31 parameters for a macro. 
The parameter list may end with an ellipsis (...) in this case, the identifier 
\texttt{\_VA\_ARGS\_} may appear in the replacement list.

\textbf{Function-like macro invocation:}
An identifier followed by a comma-separated list of arguments in 
parentheses. The number of arguments should match the number of 
parameters in the macro definition, unless the parameter list in the 
definition ends with an ellipsis. In this latter case, the number of 
arguments in the invocation should exceed the number of parameters in 
the definition. The excess are called \textit{trailing arguments}. Once the 
preprocessor identifies a function-like macro invocation, argument 
substitution takes place. A parameter in the replacement code is replaced 
by the corresponding argument. If trailing arguments are permitted by the 
macro definition, they are merged with the intervening commas to replace the 
identifier \texttt{\_VA\_ARGS\_}, as if they were a single argument. Any macro 
invocations contained in the argument itself are completely replaced before 
the argument replaces its corresponding parameter in the replacement 
code.

A macro argument can be empty (consisting of zero preprocessing tokens). 
For example,
\begin{verbatim}
#define SUM(a,b,c) a + b + c
SUM(1,,3) /* No error message. 
 1 is substituted for a, 3 is substituted for c. */
\end{verbatim}

If the identifier list does not end with an ellipsis, the number of arguments in a 
macro invocation must be the same as the number of parameters in the
corresponding macro definition. During parameter substitution, any arguments remaining after all specified arguments have been substituted (including any separating commas) are combined into one argument called the variable argument. The variable argument will replace any occurrence of the identifier __VA_ARGS__ in the replacement list. The following example illustrates this:

```c
#define debug(...) fprintf(stderr, __VA_ARGS__)
```

dbg("flag"); /* Becomes fprintf(stderr, "flag"); */

Commas in the macro invocation argument list do not act as argument separators when they are:

- In character constants
- In string literals
- Surrounded by parentheses

The following line defines the macro SUM as having two parameters a and b and the replacement tokens \((a + b)\):

```c
#define SUM(a,b) (a + b)
```

This definition would cause the preprocessor to change the following statements (if the statements appear after the previous definition):

```c
c = SUM(x,y);
c = d * SUM(x,y);
```

In the output of the preprocessor, these statements would appear as:

```c
c = (x + y);
c = d * (x + y);
```

Use parentheses to ensure correct evaluation of replacement text. For example, the definition:

```c
#define SQR(c) ((c) * (c))
```

requires parentheses around each parameter \(c\) in the definition in order to correctly evaluate an expression like:

```c
y = SQR(a + b);
```

The preprocessor expands this statement to:

```c
y = ((a + b) * (a + b));
```

Without parentheses in the definition, the correct order of evaluation is not preserved, and the preprocessor output is:

```c
y = (a + b * a + b);
```

Arguments of the \# and ## operators are converted before replacement of parameters in a function-like macro.

Once defined, a preprocessor identifier remains defined and in scope independent of the scoping rules of the language. The scope of a macro definition begins at the definition and does not end until a corresponding #undef directive is encountered. If there is no corresponding #undef directive, the scope of the macro definition lasts until the end of the translation unit.

A recursive macro is not fully expanded. For example, the definition

```c
#define x(a,b) x(a+1,b+1) + 4
```
expands
\[ x(20,10) \]
to
\[ x(20+1,10+1) + 4 \]
rather than trying to expand the macro \( x \) over and over within itself. After the macro \( x \) is expanded, it is a call to function \( x() \).

A definition is not required to specify replacement tokens. The following definition removes all instances of the token \( \text{debug} \) from subsequent lines in the current file:
\[
\#define \text{debug}
\]
You can change the definition of a defined identifier or macro with a second preprocessor \#define directive only if the second preprocessor \#define directive is preceded by a preprocessor \#undef directive. The \#undef directive nullifies the first definition so that the same identifier can be used in a redefinition.

Within the text of the program, the preprocessor does not scan character constants or string constants for macro invocations.

The following example program contains two macro definitions and a macro invocation that refers to both of the defined macros:

/\
  ** This example illustrates \#define directives.  **\
/\n#include <stdio.h>
#define SQR(s) ((s) * (s))
#define PRNT(a,b) \n printf("value 1 = \%d\n", a); \n printf("value 2 = \%d\n", b); 

int main(void)
{
  int x = 2;
  int y = 3;

  PRNT(SQR(x),y);

  return(0);
}

After being interpreted by the preprocessor, this program is replaced by code equivalent to the following:
#include <stdio.h>

int main(void)
{
  int x = 2;
  int y = 3;

  printf("value 1 = \%d\n", ( (x) * (x) ) );
  printf("value 2 = \%d\n", y);

  return(0);
}

This program produces the following output:
Variadic macro extensions:

Variadic macro extensions refer to two extensions to C99 and Standard C++ related to macros with variable number of arguments. One extension is a mechanism for renaming the variable argument identifier from __VA_ARGS__ to a user-defined identifier. The other extension provides a way to remove the dangling comma in a variadic macro when no variable arguments are specified. Both extensions have been implemented to facilitate porting programs developed with GNU C and C++.

The following examples demonstrate the use of an identifier in place of __VA_ARGS__. The first definition of the macro debug exemplifies the usual usage of __VA_ARGS__. The second definition shows the use of the identifier args in place of __VA_ARGS__.

```c
#define debug1(format, ...) printf(format, __VA_ARGS__)
#define debug2(format, args ...) printf(format, args)
```

<table>
<thead>
<tr>
<th>Invocation</th>
<th>Result of Macro Expansion</th>
</tr>
</thead>
<tbody>
<tr>
<td>debug1(&quot;Hello %s\n&quot;,&quot;World&quot;);</td>
<td>printf(&quot;Hello %s\n&quot;,&quot;World&quot;);</td>
</tr>
<tr>
<td>debug2(&quot;Hello %s\n&quot;,&quot;World&quot;);</td>
<td>printf(&quot;Hello %s\n&quot;,&quot;World&quot;);</td>
</tr>
</tbody>
</table>

The preprocessor removes the trailing comma if the variable arguments to a function macro are omitted or empty and the comma followed by ## precedes the variable argument identifier in the function macro definition.

The #undef directive

A preprocessor undef directive causes the preprocessor to end the scope of a preprocessor definition.

#undef directive syntax

```c
#-define identifier
```

If the identifier is not currently defined as a macro, #undef is ignored.

The following directives define BUFFER and SQR:

```c
#define BUFFER 512
#define SQR(x) ((x) * (x))
```

The following directives nullify these definitions:

```c
#-define BUFFER
#-define SQR
```
Any occurrences of the identifiers BUFFER and SQR that follow these #undef directives are not replaced with any replacement tokens. Once the definition of a macro has been removed by an #undef directive, the identifier can be used in a new #define directive.

The # operator

The # (single number sign) operator converts a parameter of a function-like macro into a character string literal. For example, if macro ABC is defined using the following directive:

```
#define ABC(x) #x
```

all subsequent invocations of the macro ABC would be expanded into a character string literal containing the argument passed to ABC. For example:

<table>
<thead>
<tr>
<th>Invocation</th>
<th>Result of Macro Expansion</th>
</tr>
</thead>
<tbody>
<tr>
<td>ABC(1)</td>
<td>&quot;1&quot;</td>
</tr>
<tr>
<td>ABC(Hello there)</td>
<td>&quot;Hello there&quot;</td>
</tr>
</tbody>
</table>

The # operator should not be confused with the null directive.

Use the # operator in a function-like macro definition according to the following rules:

- A parameter following # operator in a function-like macro is converted into a character string literal containing the argument passed to the macro.
- White-space characters that appear before or after the argument passed to the macro are deleted.
- Multiple white-space characters imbedded within the argument passed to the macro are replaced by a single space character.
- If the argument passed to the macro contains a string literal and if a \ (backslash) character appears within the literal, a second \ character is inserted before the original \ when the macro is expanded.
- If the argument passed to the macro contains a " (double quotation mark) character, a \ character is inserted before the " when the macro is expanded.
- The conversion of an argument into a string literal occurs before macro expansion on that argument.
- If more than one ## operator or # operator appears in the replacement list of a macro definition, the order of evaluation of the operators is not defined.
- If the result of the macro expansion is not a valid character string literal, the behavior is undefined.

The following examples demonstrate the use of the # operator:

```
#define STR(x) #x
#define XSTR(x) STR(x)
#define ONE 1
```

<table>
<thead>
<tr>
<th>Invocation</th>
<th>Result of Macro Expansion</th>
</tr>
</thead>
<tbody>
<tr>
<td>STR\n &quot;\n &quot;\n \n \n\n)</td>
<td>&quot;\n &quot;\n \n &quot;\n \n\n</td>
</tr>
<tr>
<td>STR(ONE)</td>
<td>&quot;ONE&quot;</td>
</tr>
<tr>
<td>XSTR(ONE)</td>
<td>&quot;1&quot;</td>
</tr>
<tr>
<td>XSTR(&quot;hello&quot;)</td>
<td>&quot;\hello&quot;</td>
</tr>
</tbody>
</table>
Related information
• “The null directive (#)” on page 401

The ## operator
The ## (double number sign) operator concatenates two tokens in a macro invocation (text and/or arguments) given in a macro definition.

If a macro XY was defined using the following directive:

```
#define XY(x, y) x##y
```

the last token of the argument for x is concatenated with the first token of the argument for y.

Use the ## operator according to the following rules:
• The ## operator cannot be the very first or very last item in the replacement list of a macro definition.
• The last token of the item in front of the ## operator is concatenated with first token of the item following the ## operator.
• Concatenation takes place before any macros in arguments are expanded.
• If the result of a concatenation is a valid macro name, it is available for further replacement even if it appears in a context in which it would not normally be available.
• If more than one ## operator and/or # operator appears in the replacement list of a macro definition, the order of evaluation of the operators is not defined.

The following examples demonstrate the use of the ## operator:

```
#define ArgArg(x, y) x##y
#define ArgText(x) x##TEXT
#define TextArg(x) TEXT##x
#define TextText TEXT##text
#define Jitter 1
#define bug 2
#define Jitterbug 3
```

<table>
<thead>
<tr>
<th>Invocation</th>
<th>Result of Macro Expansion</th>
</tr>
</thead>
<tbody>
<tr>
<td>ArgArg(lady, bug)</td>
<td>&quot;ladybug&quot;</td>
</tr>
<tr>
<td>ArgText(con)</td>
<td>&quot;conTEXT&quot;</td>
</tr>
<tr>
<td>TextArg(book)</td>
<td>&quot;TEXTbook&quot;</td>
</tr>
<tr>
<td>TextText</td>
<td>&quot;TEXTtext&quot;</td>
</tr>
<tr>
<td>ArgArg(Jitter, bug)</td>
<td>3</td>
</tr>
</tbody>
</table>

Standard predefined macro names
Both C and C++ provide the following predefined macro names as specified in the ISO C language standard. Except for __FILE__ and __LINE__, the value of the predefined macros remains constant throughout the translation unit.

<table>
<thead>
<tr>
<th><strong>DATE</strong></th>
<th>A character string literal containing the date when the source file was compiled.</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>The value of <strong>DATE</strong> changes as the compiler processes any include files that are part of your source program. The date is in the form:</td>
</tr>
<tr>
<td></td>
<td>&quot;Mmm dd yyyy&quot;</td>
</tr>
</tbody>
</table>
where:

Mmm  Represents the month in an abbreviated form (Jan, Feb, Mar, Apr, May, Jun, Jul, Aug, Sep, Oct, Nov, or Dec).

dd    Represents the day. If the day is less than 10, the first d is a blank character.

yyyy  Represents the year.

_FILE_  A character string literal containing the name of the source file.

The value of _FILE_ changes as the compiler processes include files that are part of your source program. It can be set with the #line directive.

_LINE_  An integer representing the current source line number.

The value of _LINE_ changes during compilation as the compiler processes subsequent lines of your source program. It can be set with the #line directive.

__STDC__  For C, the integer 1 (one) indicates that the C compiler supports the ISO standard. (When a macro is undefined, it behaves as if it had the integer value 0 when used in a #if statement.)

For C++, this macro is predefined to have the value 0 (zero). This indicates that the C++ language is not a proper superset of C, and that the compiler does not conform to ISO C.

_________________________ C only __________________________

__STDC_HOSTED__  The value of this C99 macro is 1, indicating that the C compiler is a hosted implementation.

_________________________ End of C only __________________________

_________________________ C only __________________________

__STDC_VERSION__  The integer constant of type long int: 199409L for the C89 language level, 199901L for C99.

_________________________ End of C only __________________________

__TIME__  A character string literal containing the time when the source file was compiled.

The value of _TIME_ changes as the compiler processes any include files that are part of your source program. The time is in the form:

"hh:mm:ss"

where:

hh    Represents the hour.

mm    Represents the minutes.

ss    Represents the seconds.
C++ only

__cplusplus  For C++ programs, this macro expands to the long integer literal 199711L, indicating that the compiler is a C++ compiler. For C programs, this macro is not defined. Note that this macro name has no trailing underscores.

Related information
- "The #line directive" on page 398
- "Object-like macros" on page 384

File inclusion directives
File inclusion directives consist of:
- The #include directive which inserts text from another source file
- The #include_next directive

The #include directive
A preprocessor include directive causes the preprocessor to replace the directive with the contents of the specified file.

#include directive syntax

```
#include "file_name"
#include <file_name>
#include <header_name> identifiers
```

In all C and C++ implementations, the preprocessor resolves macros contained in an #include directive. After macro replacement, the resulting token sequence must consist of a file name enclosed in either double quotation marks or the characters < and >.

For example:
```
define MONTH <july.h>
#include MONTH
```

If the file name is enclosed in double quotation marks, for example:
```
#include "payroll.h"
```

the preprocessor treats it as a user-defined file, and searches for the file in a manner defined by the preprocessor.

If the file name is enclosed in angle brackets, for example:
```
#include <stdio.h>
```

it is treated as a system-defined file, and the preprocessor searches for the file in a manner defined by the preprocessor.
The new-line and > characters cannot appear in a file name delimited by < and >. The new-line and " (double quotation marks) character cannot appear in a file name delimited by " and ", although > can.

Declarations that are used by several files can be placed in one file and included with #include in each file that uses them. For example, the following file defs.h contains several definitions and an inclusion of an additional file of declarations:

```c
/* defs.h */
#define TRUE 1
#define FALSE 0
#define BUFFERSIZE 512
#define MAX_ROW 66
#define MAX_COLUMN 80
int hour;
int min;
int sec;
#include "mydefs.h"
```

You can embed the definitions that appear in defs.h with the following directive:

```
#include "defs.h"
```

In the following example, a #define combines several preprocessor macros to define a macro that represents the name of the C standard I/O header file. A #include makes the header file available to the program.

```
#define C_IO_HEADER <stdio.h>

/* The following is equivalent to: */
/* #include <stdio.h> */

#include C_IO_HEADER
```

### The #include_next directive

The preprocessor directive #include_next instructs the preprocessor to continue searching for the specified file name, and to include the subsequent instance encountered after the current directory. The syntax of the directive is similar to that of #include.

The language feature is an extension to C and C++. It extends the techniques available to address the issue of duplicate file names among applications and shared libraries.

---

**Conditional compilation directives**

A *preprocessor conditional compilation directive* causes the preprocessor to conditionally suppress the compilation of portions of source code. These directives test a constant expression or an identifier to determine which tokens the preprocessor should pass on to the compiler and which tokens should be bypassed during preprocessing. The directives are:

- [The #if and #elif directives](#) which conditionally include or suppress portions of source code, depending on the result of a constant expression
The #ifdef directive which conditionally includes source text if a macro name is defined

The #ifndef directive which conditionally includes source text if a macro name is not defined

The #else directive which conditionally includes source text if the previous #if, #ifdef, #ifndef, or #elif test fails

The #endif directive which ends conditional text

The preprocessor conditional compilation directive spans several lines:

- The condition specification line (beginning with #if, #ifdef, or #ifndef)
- Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to a nonzero value (optional)
- The #elif line (optional)
- Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to a nonzero value (optional)
- The #else line (optional)
- Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to zero (optional)
- The preprocessor #endif directive

For each #if, #ifdef, and #ifndef directive, there are zero or more #elif directives, zero or one #else directive, and one matching #endif directive. All the matching directives are considered to be at the same nesting level.

You can nest conditional compilation directives. In the following directives, the first #else is matched with the #if directive.

```c
#ifdef MACNAME
  /* tokens added if MACNAME is defined */
#endif

if TEST <=10
  /* tokens added if MACNAME is defined and TEST <= 10 */
#else
  /* tokens added if MACNAME is defined and TEST > 10 */
#endif
#else /* tokens added if MACNAME is not defined */
#endif
```

Each directive controls the block immediately following it. A block consists of all the tokens starting on the line following the directive and ending at the next conditional compilation directive at the same nesting level.

Each directive is processed in the order in which it is encountered. If an expression evaluates to zero, the block following the directive is ignored.

When a block following a preprocessor directive is to be ignored, the tokens are examined only to identify preprocessor directives within that block so that the conditional nesting level can be determined. All tokens other than the name of the directive are ignored.

Only the first block whose expression is nonzero is processed. The remaining blocks at that nesting level are ignored. If none of the blocks at that nesting level has been processed and there is a #else directive, the block following the #else directive is processed. If none of the blocks at that nesting level has been processed and there is no #else directive, the entire nesting level is ignored.
The #if and #elif directives

The #if and #elif directives compare the value of constant_expression to zero:

### #if and #elif directive syntax

```
#define constant_expression token_sequence
```

If the constant expression evaluates to a nonzero value, the lines of code that immediately follow the condition are passed on to the compiler.

If the expression evaluates to zero and the conditional compilation directive contains a preprocessor #elif directive, the source text located between the #elif and the next #elif or preprocessor #else directive is selected by the preprocessor to be passed on to the compiler. The #elif directive cannot appear after the preprocessor #else directive.

All macros are expanded, any `defined()` expressions are processed and all remaining identifiers are replaced with the token 0.

The constant_expression that is tested must be integer constant expressions with the following properties:

- No casts are performed.
- Arithmetic is performed using `long int` values.
- The constant_expression can contain defined macros. No other identifiers can appear in the expression.
- The constant_expression can contain the unary operator `defined`. This operator can be used only with the preprocessor keyword #if or #elif. The following expressions evaluate to 1 if the identifier is defined in the preprocessor, otherwise to 0:
  ```
  defined identifier
defined( identifier)
  ```

For example:

```
#define TEST1
#define TEST2

if defined(TEST1) || defined(TEST2)
```

**Note:** If a macro is not defined, a value of 0 (zero) is assigned to it. In the following example, TEST must be a macro identifier:

```
#define TEST

#if TEST >= 1
  printf("i = %d\n", i);
  printf("array[i] = %d\n", array[i]);
#elif TEST < 0
  printf("array subscript out of bounds \n");
#endif
```

The #ifdef directive

The #ifdef directive checks for the existence of macro definitions.

If the identifier specified is defined as a macro, the lines of code that immediately follow the condition are passed on to the compiler.
The following example defines MAX_LEN to be 75 if EXTENDED is defined for the preprocessor. Otherwise, MAX_LEN is defined to be 50.

```c
#define EXTENDED
#define MAX_LEN 75
#else
#define MAX_LEN 50
#endif
```

**The #ifndef directive**

The #ifndef directive checks whether a macro is not defined.

If the identifier specified is not defined as a macro, the lines of code immediately follow the condition are passed on to the compiler.

```c
#ifndef identifier
token_sequence
#endif
```

An identifier must follow the #ifndef keyword. The following example defines MAX_LEN to be 50 if EXTENDED is not defined for the preprocessor. Otherwise, MAX_LEN is defined to be 75.

```c
#define EXTENDED
#define MAX_LEN 50
#else
#define MAX_LEN 75
#endif
```

**The #else directive**

If the condition specified in the #if, #ifdef, or #ifndef directive evaluates to 0, and the conditional compilation directive contains a preprocessor #else directive, the lines of code located between the preprocessor #else directive and the preprocessor #endif directive is selected by the preprocessor to be passed on to the compiler.

```c
#else
token_sequence
#endif
```

**The #endif directive**

The preprocessor #endif directive ends the conditional compilation directive.
#endif directive syntax

Examples of conditional compilation directives

The following example shows how you can nest preprocessor conditional compilation directives:

```c
#if defined(TARGET1)
  # define SIZEOF_INT 16
  #ifdef PHASE2
    # define MAX_PHASE 2
  # else
    # define MAX_PHASE 8
  # endif
  #endif defined(TARGET2)
  # define SIZEOF_INT 32
  # define MAX_PHASE 16
#else
  # define SIZEOF_INT 32
  # define MAX_PHASE 32
#endif
```

The following program contains preprocessor conditional compilation directives:

```c
/**
** This example contains preprocessor conditional compilation directives.
**/
#include <stdio.h>

int main(void)
{  
  static int array[ ] = { 1, 2, 3, 4, 5 };  
  int i;

  for (i = 0; i <= 4; i++)
  {  
    array[i] *= 2;
    #if TEST >= 1
      printf("i = %d\n", i);
      printf("array[i] = %d\n", array[i]);
      #endif
  }
  return(0);
}
```

Message generation directives

Message generation directives include the following:

- **The #error directive** which defines text for a compile-time error message
- **The #warning directive** which defines text for a compile-time warning message
- **The #line directive** which supplies a line number for compiler messages

Related information

- "Conditional compilation directives" on page 393
The `#error` directive

A preprocessor error directive causes the preprocessor to generate an error message and causes the compilation to fail.

`#error` directive syntax

```
#error preprocessor_token
```

The argument `preprocessor_token` is not subject to macro expansion.

The `#error` directive is often used in the `#else` portion of a `#if-#elif-#else` construct, as a safety check during compilation. For example, `#error` directives in the source file can prevent code generation if a section of the program is reached that should be bypassed.

For example, the directive

```
#define BUFFER_SIZE 255

#if BUFFER_SIZE < 256
#error "BUFFER_SIZE is too small."
#endif
```

generates the error message:

`BUFFER_SIZE is too small.`

The `#warning` directive

```
IBM extension
```

A preprocessor warning directive causes the preprocessor to generate a warning message but allows compilation to continue. The argument to `#warning` is not subject to macro expansion.

`#warning` directive syntax

```
#warning preprocessor_token
```

The preprocessor `#warning` directive is a language extension provided to facilitate handling programs developed with GNU C. The IBM implementation preserves multiple white spaces.

```
End of IBM extension
```

The `#line` directive

A preprocessor line control directive supplies line numbers for compiler messages. It causes the compiler to view the line number of the next source line as the specified number.
#line directive syntax

```
#line decimal_constant ["file_name"]
```

In order for the compiler to produce meaningful references to line numbers in preprocessed source, the preprocessor inserts #line directives where necessary (for example, at the beginning and after the end of included text).

A file name specification enclosed in double quotation marks can follow the line number. If you specify a file name, the compiler views the next line as part of the specified file. If you do not specify a file name, the compiler views the next line as part of the current source file.

At the C99 language level, the maximum value of the #line preprocessing directive is 2147483647.

In all C and C++ implementations, the token sequence on a #line directive is subject to macro replacement. After macro replacement, the resulting character sequence must consist of a decimal constant, optionally followed by a file name enclosed in double quotation marks.

You can use #line control directives to make the compiler provide more meaningful error messages. The following example program uses #line control directives to give each function an easily recognizable line number:

```
/**
 ** This example illustrates #line directives.
 **/

#include <stdio.h>
#define LINE200 200

int main(void)
{
    func_1();
    func_2();
}

#line 100
func_1()
{
    printf("Func_1 - the current line number is %d\n", _LINE_);
}

#line LINE200
func_2()
{
    printf("Func_2 - the current line number is %d\n", _LINE_);
}

This program produces the following output:
Func_1 - the current line number is 102
Func_2 - the current line number is 202
```
Assertion directives

An assertion directive is an alternative to a macro definition, used to define the computer or system the compiled program will run on. Assertions are usually predefined, but you can define them with the #assert preprocessor directive.

### #assert directive syntax

```
#—assert—predicate—(—answer—)
```

The predicate represents the assertion entity you are defining. The answer represents a value you are assigning to the assertion. You can make several assertions using the same predicate and different answers. All the answers for any given predicate are simultaneously true. For example, the following directives create assertions regarding font properties:

```c
#assert font(arial)
#assert font(blue)
```

Once an assertion has been defined, the assertion predicate can be used in conditional directives to test the current system. The following directive tests whether arial or blue is asserted for font:

```c
#if #font(arial) || #font(blue)
```

You can test whether any answer is asserted for a predicate by omitting the answer in the conditional:

```c
#if #font
```

Assertions can be cancelled with the #unassert directive. If you use the same syntax as the #assert directive, the directive cancels only the answer you specify. For example, the following directive cancels the arial answer for the font predicate:

```c
#unassert font(arial)
```

An entire predicate is cancelled by omitting the answer from the #unassert directive. The following directive cancels the font directive altogether:

```c
#unassert font
```

### Related information

- “Conditional compilation directives” on page 393

### Predefined assertions

The following assertions are predefined for the Linux platform:

Table 26. Predefined assertions for Linux

<table>
<thead>
<tr>
<th>#machine(powerpc)</th>
<th>#system(unix)</th>
<th>#cpu(powerpc)</th>
</tr>
</thead>
<tbody>
<tr>
<td>#machine(bigendian)</td>
<td>#system(posix)</td>
<td></td>
</tr>
</tbody>
</table>

---

End of IBM extension
The null directive (#)

The null directive performs no action. It consists of a single # on a line of its own.

The null directive should not be confused with the # operator or the character that starts a preprocessor directive.

In the following example, if MINVAL is a defined macro name, no action is performed. If MINVAL is not a defined identifier, it is defined 1.

```c
#ifdef MINVAL
#
#else
#define MINVAL 1
#endif
```

Related information
- "The # operator" on page 389

Pragma directives

A pragma is an implementation-defined instruction to the compiler. It has the general form:

```
#pragma directive syntax
```

The `character_sequence` is a series of characters giving a specific compiler instruction and arguments, if any. The token `STDC` indicates a standard pragma; consequently, no macro substitution takes place on the directive. The `new-line` character must terminate a pragma directive.

The `character_sequence` on a pragma is subject to macro substitutions. For example,

```c
#define XX_ISO_DATA isolated_call(LG_ISO_DATA)
// ...
#pragma XX_ISO_DATA
```

Note: You can also use the _Pragma operator syntax to specify a pragma directive; for details, see "The _Pragma preprocessing operator" on page 143.

More than one pragma construct can be specified on a single pragma directive. The compiler ignores unrecognized pragmas.

Standard C pragmas are described in "Standard pragmas." IBM Pragmas available for XL C/C++ are described in "General purpose pragmas" in the XL C/C++ Compiler Reference

Standard pragmas

A standard pragma is a pragma preprocessor directive for which the C Standard defines the syntax and semantics and for which no macro replacement is performed. A standard pragma must be one of the following:
The FP_CONTRACT and FENV_ACCESS pragmas are recognized and ignored.

CX_LIMITED_RANGE is described below.

\texttt{\textbf{pragma \texttt{STDC CX\_LIMITED\_RANGE}}}  

The usual mathematical formulas for complex multiplication, division, and absolute value are problematic because of their treatment of infinities and because of undue overflow and underflow. The usual formulas are as follows:

\[(x + iy) \times (u + iv) = (xu - yv) + i(yu + xv)\]

\[(x + iy)/(u + iv) = [(xu + yv) + i(yu - xv)]/(u^2 + v^2)\]

\[|x + iy| = \sqrt{x^2 + y^2}\]

By default, the compiler uses slightly more complex but mathematically safer algorithms to implement these calculations. Where you determine that the usual mathematical formulas are safe, you can use the \texttt{STDC CX\_LIMITED\_RANGE} pragma to inform the compiler that, when the state is "on", the formulas are acceptable. In doing so, you allow the compiler to generate faster code for these computations. When the state is "off", the compiler will continue to use the safer algorithms. For details on the implementation of this pragma, see \texttt{#pragma \texttt{STDC CX\_LIMITED\_RANGE}} in the \textit{XL C/C++ Compiler Reference}.  

---

402 XL C/C++ Language Reference
## Appendix A. The IBM XL C language extensions

This appendix presents the IBM XL C extensions in the following categories:

- “C99 features as extensions to C89”
- “Extensions for Unicode support” on page 405
- “Extensions for GNU C compatibility” on page 405
- “Extensions for VMX support” on page 407

### C99 features as extensions to C89

The following features are enabled by default when you compile with the xlc invocation command. They are also enabled with the options `-qlanglvl=extc89` (the default in xlc), `-qlanglvl=stdc99`, `-qlanglvl=extc99` and `-qlanglvl=extended`. For more information on these options, see the `qlanglvl` option in the XL C/C++ Compiler Reference.

Table 27. Default C99 features as extensions to C89

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hexadecimal floating-point constants</td>
<td>“Hexadecimal floating-point literals” on page 23</td>
</tr>
<tr>
<td><code>__func__</code> predefined identifier</td>
<td>“The <code>__func__</code> predefined identifier” on page 16</td>
</tr>
<tr>
<td>Concatenation of wide and non-wide character strings</td>
<td>“String concatenation” on page 29</td>
</tr>
<tr>
<td>Mixed declarations and code</td>
<td>“Overview of data declarations and definitions” on page 41</td>
</tr>
<tr>
<td><code>long long</code> data type</td>
<td>“Integral types” on page 50</td>
</tr>
<tr>
<td>Complex data type</td>
<td>“Complex floating-point types” on page 52</td>
</tr>
<tr>
<td><code>__bool__</code> data type</td>
<td>“Boolean types” on page 50</td>
</tr>
<tr>
<td>Trailing comma allowed in enum declaration</td>
<td>“Enumeration type definition” on page 63</td>
</tr>
<tr>
<td>Duplicate type qualifiers</td>
<td>“Type qualifiers” on page 68</td>
</tr>
<tr>
<td>Variable length arrays</td>
<td>“Variable length arrays” on page 87</td>
</tr>
<tr>
<td>Non-lvalue array subscripts</td>
<td>“Array subscripting operator [ ]” on page 122</td>
</tr>
<tr>
<td>Flexible array members at the end of a structure or union</td>
<td>“Flexible array members” on page 57</td>
</tr>
<tr>
<td>Non-constant expression in initializer for structure or union</td>
<td>“Initialization of structures and unions” on page 94</td>
</tr>
<tr>
<td>Designated initializers</td>
<td>“Designated initializers for aggregate types (C only)” on page 91</td>
</tr>
<tr>
<td>Removal of implicit function declaration</td>
<td>“Function declarations” on page 192</td>
</tr>
<tr>
<td>Removal of implicit <code>int</code> return type in function declarations</td>
<td>“Function return type specifiers” on page 201</td>
</tr>
<tr>
<td>Static arrays as function parameters</td>
<td>“Static array indices in function parameter declarations (C only)” on page 205</td>
</tr>
<tr>
<td>Variable arguments in function-like macros</td>
<td>“Function-like macros” on page 385</td>
</tr>
</tbody>
</table>
Table 27. Default C99 features as extensions to C89 (continued)

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Empty arguments in function-like macros</td>
<td>“Function-like macros” on page 385</td>
</tr>
<tr>
<td>Additional predefined macro names</td>
<td>“Standard predefined macro names” on page 390</td>
</tr>
<tr>
<td>Compound literals</td>
<td>“Compound literals” on page 124</td>
</tr>
<tr>
<td>_Pragma operator</td>
<td>“The _Pragma preprocessing operator” on page 143</td>
</tr>
<tr>
<td>Standard pragmas</td>
<td>“Standard pragmas” on page 401</td>
</tr>
<tr>
<td>New limit for #line directive</td>
<td>“The #line directive” on page 398</td>
</tr>
</tbody>
</table>

The following feature is enabled by default when you compile with the `xlc` invocation command. It is also enabled with the options -qlanglvl=extc89 (the default in `xlc`), -qlanglvl=stdc99, -qlanglvl=extc99 and -qlanglvl=extended. It is also enabled or disabled by a specific compiler option, which is listed in the table below.

Table 28. Default C99 features as extensions to C89, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td>Digraphs</td>
<td>“Digraph characters” on page 34</td>
<td>-q[no]digraph</td>
</tr>
</tbody>
</table>

The following features are enabled by default when you compile with the `c99` invocation command, or with the -qlanglvl=stdc99 or -qlanglvl=extc99 compiler options or related pragmas. They are also enabled or disabled by specific compiler options, which are listed in the table below; these compiler options are enabled in the default configuration file for the `xlc` invocation command.

Table 29. Strict C99 features as extensions to C89, with individual option controls enabled by default for `xlc` in the configuration file

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td>C++ style comments</td>
<td>“Comments” on page 35</td>
<td>-q[no]cpluscmt</td>
</tr>
<tr>
<td>The inline function specifier</td>
<td>“The inline function specifier” on page 197</td>
<td>-qkeyword=inline</td>
</tr>
</tbody>
</table>

The following features are enabled only when you compile with the `c99` invocation command, or with the -qlanglvl=stdc99 or -qlanglvl=extc99 compiler options or related pragmas. They are also enabled by specific compiler options, listed in the table below.

Table 30. Strict C99 features as extensions to C89, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td>Universal character names</td>
<td>“The Unicode standard” on page 32</td>
<td>-qlanglvl=[no]ucs</td>
</tr>
<tr>
<td>The restrict type qualifier</td>
<td>“The restrict type qualifier” on page 70</td>
<td>-qkeyword=restrict</td>
</tr>
</tbody>
</table>
The following feature is enabled only when you compile with the c99 invocation command, or with the -qlanglvl=stdc99 or -qlanglvl=extc99 compiler options or related pragmas.

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Unsuffixed long long integer literals</td>
<td>“Decimal integer literals” on page 20</td>
</tr>
</tbody>
</table>

Related information
- `-qcpluscmt` `-qkeyword` and `-qdigraph` in the XL C/C++ Compiler Reference
- “How to choose a compiler invocation” in the XL C/C++ Compiler Reference

Extensions for Unicode support

The following feature requires compilation with the use of an additional option.

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Required compilation option</th>
</tr>
</thead>
<tbody>
<tr>
<td>UTF-16, UTF-32 literals</td>
<td>“UTF literals” on page 33</td>
<td><code>-qutf</code></td>
</tr>
</tbody>
</table>

Related information
- `-qutf` in the XL C/C++ Compiler Reference

Extensions for GNU C compatibility

The following features are enabled by default when you compile with the xlc invocation command. They are also enabled with the options -qlanglvl=extc89 (the default in xlc), -qlanglvl=extc99, and -qlanglvl=extended.

Table 31. Default IBM XL C extensions for GNU C compatibility

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Alternate keywords</td>
<td>“Keywords for language extensions” on page 14</td>
</tr>
<tr>
<td><em>extension</em> keyword</td>
<td>“Keywords for language extensions” on page 14</td>
</tr>
<tr>
<td>asm labels</td>
<td>“Assembly labels” on page 17</td>
</tr>
<tr>
<td>Complex literal suffixes</td>
<td>“Complex literals” on page 24</td>
</tr>
<tr>
<td>Global register variables</td>
<td>“Global variables in specified registers (C only)” on page 48</td>
</tr>
<tr>
<td>Placement of flexible array members anywhere in structure or union</td>
<td>“Flexible array members” on page 57</td>
</tr>
<tr>
<td>Static initialization of flexible array members of aggregates</td>
<td>“Flexible array members” on page 57</td>
</tr>
<tr>
<td>Zero-extent arrays</td>
<td>“Zero-extent array members” on page 58</td>
</tr>
<tr>
<td>Type attributes</td>
<td>“Type attributes” on page 74</td>
</tr>
<tr>
<td>Variable attributes</td>
<td>“Variable attributes” on page 101</td>
</tr>
<tr>
<td>Locally declared labels</td>
<td>“Locally declared labels” on page 168</td>
</tr>
<tr>
<td>Labels as values</td>
<td>“Labels as values” on page 168</td>
</tr>
<tr>
<td><strong>alignof</strong> operator</td>
<td>“The <strong>alignof</strong> operator” on page 136</td>
</tr>
</tbody>
</table>

Appendix A. The IBM XL C language extensions 405
The following features are enabled by default when you compile with the xlc invocation command. They are also enabled with the options -qlanglvl=extc89 (the default in xlc), -qlanglvl=extc99, and -qlanglvl=extended. They can also be enabled and disabled by specific compiler options, which are listed in the table below.

Table 32. IBM XL C extensions for GNU C compatibility with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option controls</th>
</tr>
</thead>
<tbody>
<tr>
<td>typedef, asm, and __asm keywords</td>
<td>“The typeof operator” on page 139, “Assembly labels” on page 17, “Inline assembly statements” on page 186</td>
<td>-qkeyword</td>
</tr>
<tr>
<td>asm inline assembly-language statements</td>
<td>“Inline assembly statements” on page 186</td>
<td>-qasm</td>
</tr>
</tbody>
</table>

The following feature requires compilation with the use of an additional option.

Table 33. IBM XL C extensions for GNU C compatibility, requiring additional compiler options

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Required compilation option</th>
</tr>
</thead>
<tbody>
<tr>
<td>Dollar signs in identifiers</td>
<td>“Characters in identifiers” on page 16</td>
<td>-qdollar</td>
</tr>
</tbody>
</table>
• \texttt{-qkeyword}, \texttt{-qasm}, and \texttt{-qdollar} in the \textit{XL C/C++ Compiler Reference}
• “How to choose a compiler invocation” in the \textit{XL C/C++ Compiler Reference}

## Extensions for VMX support

The VMX extensions are only accepted when all of the following conditions are met:

1. The \texttt{-qarch} option is set to a target architecture that supports VMX instructions.
2. The \texttt{-qenablevmx} option is in effect.
3. The \texttt{-qaltivec} option is in effect.

For more information on these options, see the \textit{XL C/C++ Compiler Reference}.

\textit{Table 34. IBM XL C extensions to support the AltiVec Application Programming Interface specification}

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Vector programming language extensions</td>
<td>“Vector types” on page 53, Appendix C,</td>
</tr>
<tr>
<td></td>
<td>“Vector data types and literals,” on page 415</td>
</tr>
</tbody>
</table>

The following features are IBM extensions to the AltiVec Application Programming Interface specification.

\textit{Table 35. IBM XL C extensions to the AltiVec Application Programming Interface specification}

<table>
<thead>
<tr>
<th>Language extension</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Initializer lists for vector constants</td>
<td>“Initialization of vectors” on page 93</td>
</tr>
<tr>
<td>typedef definitions for vector types</td>
<td>“typedef definitions” on page 66</td>
</tr>
<tr>
<td>compound literals as initializers for static vector variables</td>
<td>“Compound literals” on page 124</td>
</tr>
<tr>
<td>vector types as arguments to the \texttt{<strong>alignof</strong>} and \texttt{typeof} operators</td>
<td>“The \texttt{<strong>alignof</strong>} operator” on page 136, “The \texttt{typeof} operator” on page 139</td>
</tr>
</tbody>
</table>
Appendix B. The IBM XL C++ language extensions

This appendix presents the IBM XL C++ extensions to Standard C++ in the following categories:

- “General IBM extensions”
- “Extensions for C99 compatibility”
- “Extensions for Unicode support” on page 410
- “Extensions for GNU C compatibility” on page 410
- “Extensions for GNU C++ compatibility” on page 412
- “Extensions for VMX support” on page 413

General IBM extensions

The following feature is enabled with the `-qlanglvl=extended` option, which is the default language level used when you compile with the xIC and xlc++ invocation commands. It can also be enabled or disabled by a specific compiler option, listed in the table below.

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option controls</th>
</tr>
</thead>
<tbody>
<tr>
<td>long long data type</td>
<td>“Integral types” on page 50</td>
<td><code>-q[no]longlong</code></td>
</tr>
</tbody>
</table>

Related information

- `-qlonglong` in the XL C/C++ Compiler Reference

Extensions for C99 compatibility

IBM XL C++ adds support for the following C99 language features. All of these features are enabled with the `-qlanglvl=extended` option, which is the default language level used when you compile with the xIC and xlc++ invocation commands. For more information, see the `-qlanglvl` option in the XL C/C++ Compiler Reference.

Table 36. Default C99 features as extensions to Standard C++

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Duplicate type qualifiers</td>
<td>“Type qualifiers” on page 68</td>
</tr>
<tr>
<td>Flexible array members at the end of a structure or union</td>
<td>“Flexible array members” on page 57</td>
</tr>
<tr>
<td>_Pragma operator</td>
<td>“The _Pragma preprocessing operator” on page 143</td>
</tr>
<tr>
<td>Additional predefined macro names</td>
<td>“Standard predefined macro names” on page 390</td>
</tr>
<tr>
<td>Empty arguments in function-like macros</td>
<td>“Function-like macros” on page 385</td>
</tr>
<tr>
<td>C standard pragmas</td>
<td>“Standard pragmas” on page 401</td>
</tr>
</tbody>
</table>

The following features are enabled with the `-qlanglvl=extended` option, which is the default language level used when you compile with the xIC and xlc++
invocation commands. They can also be enabled or disabled by specific compiler options, which are listed in the table below.

Table 37. Default C99 features as extensions to Standard C++, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>func</strong> predefined identifier</td>
<td>“The <strong>func</strong> predefined identifier” on page 16</td>
<td>-qlanglvl=[no]c99__func__</td>
</tr>
<tr>
<td>Hexadecimal floating-point literals</td>
<td>“Hexadecimal floating-point literals” on page 23</td>
<td>-qlanglvl=[no]c99hexfloat</td>
</tr>
<tr>
<td>Complex data type</td>
<td>“Complex floating-point types” on page 52</td>
<td>-qlanglvl=[no]c99complex</td>
</tr>
<tr>
<td>Trailing comma allowed in enum declaration</td>
<td>“Enumeration type definition” on page 63</td>
<td>-qlanglvl=[no]trailenum</td>
</tr>
<tr>
<td>The restrict type qualifier</td>
<td>“The restrict type qualifier” on page 70</td>
<td>-q[no]keyword=restrict</td>
</tr>
<tr>
<td>Variable length arrays</td>
<td>“Variable length arrays” on page 87</td>
<td>-qlanglvl=[no]c99vla</td>
</tr>
<tr>
<td>Compound literals</td>
<td>“Compound literals” on page 124</td>
<td>-qlanglvl=[no]c99compoundliteral</td>
</tr>
<tr>
<td>Variable arguments in function-like macros</td>
<td>“Function-like macros” on page 385</td>
<td>-qlanglvl=[no]varargmacros</td>
</tr>
</tbody>
</table>

The following feature is only enabled by a specific compiler option, listed in the table below.

Table 38. C99 features as extensions to Standard C++, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td>Universal character names</td>
<td>“The Unicode standard” on page 32</td>
<td>-qlanglvl=ucs</td>
</tr>
</tbody>
</table>

Related information
- `-qlanglvl` and `-qkeyword` in the XL C++ Compiler Reference

Extensions for Unicode support

The following feature requires compilation with the use of an additional option.

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Required compilation option</th>
</tr>
</thead>
<tbody>
<tr>
<td>UTF-16, UTF-32 literals</td>
<td>“UTF literals” on page 33</td>
<td><code>-qutf</code></td>
</tr>
</tbody>
</table>

Related information
- `-qutf` in the XL C++ Compiler Reference

Extensions for GNU C compatibility

The following subset of the GNU C language extensions is enabled with the `-qlanglvl=extended` option, which is the default language level used when you compile with the xlc and xlc++ invocation commands.
The following subset of the GNU C language extensions is enabled with the `-qlanglvl=extended` option, which is the default language level used when you compile with the `xlc` and `xlc++` invocation commands. These extensions can also be enabled or disabled by specific compiler options, which are listed in the table below.

### Table 40. Default IBM XL C++ extensions for compatibility with GNU C, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option controls</th>
</tr>
</thead>
<tbody>
<tr>
<td>Alternate keywords</td>
<td>“Keywords for language extensions” on page 14</td>
<td><code>-q[no]keyword=token</code></td>
</tr>
<tr>
<td><strong>extension</strong> keyword</td>
<td>“Keywords for language extensions” on page 14</td>
<td><code>-q[no]keyword=__extension__</code></td>
</tr>
<tr>
<td>Complex literal suffixes</td>
<td>“Complex literals” on page 24</td>
<td><code>-qlanglvl=[no]gnu_suffixij</code></td>
</tr>
<tr>
<td>typeof operator</td>
<td>“The typeof operator” on page 139</td>
<td><code>-q[no]keyword=typeof</code></td>
</tr>
<tr>
<td>Locally declared labels</td>
<td>“Locally declared labels” on page 168</td>
<td><code>-qlanglvl=[no]gnu_locallabel</code></td>
</tr>
<tr>
<td>Labels as values</td>
<td>“Labels as values” on page 168</td>
<td><code>-qlanglvl=[no]gnu_labelvalue</code></td>
</tr>
<tr>
<td>Computed goto statements</td>
<td>“Computed goto statement” on page 185</td>
<td><code>-qlanglvl=[no]gnu_computedgoto</code></td>
</tr>
<tr>
<td>inline assembly-language statements</td>
<td>“Inline assembly statements” on page 186</td>
<td><code>-qasm</code></td>
</tr>
<tr>
<td>Variadic macro extensions</td>
<td>“Variadic macro extensions” on page 388</td>
<td><code>-qlanglvl=[no]gnu_varargmacros</code></td>
</tr>
</tbody>
</table>
Table 40. Default IBM XL C++ extensions for compatibility with GNU C, with individual option controls (continued)

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option controls</th>
</tr>
</thead>
<tbody>
<tr>
<td>#warning preprocessor directive</td>
<td>“The #warning directive” on page 398</td>
<td>-qlanglvl=[no]gnu_warning</td>
</tr>
<tr>
<td>#include_next preprocessor directive</td>
<td>“The #include_next directive” on page 393</td>
<td>-qlanglvl=[no]gnu_include_next</td>
</tr>
<tr>
<td>#assert, #unassert, #cpu, #machine, #system preprocessor directives</td>
<td>“Assertion directives” on page 400</td>
<td>-qlanglvl=[no]gnu_assert</td>
</tr>
</tbody>
</table>

The following feature requires compilation with the use of an additional option, listed in the table below.

Table 41. IBM XL C++ extensions for GNU C compatibility, requiring additional compiler options

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Required compilation option</th>
</tr>
</thead>
<tbody>
<tr>
<td>Dollar signs in identifiers</td>
<td>“Characters in identifiers” on page 16</td>
<td>-qdollar</td>
</tr>
</tbody>
</table>

Related information

- **-qdollar** - **-qlanglvl** - **-qkeyword** and -**qasm** in the [XL C/C++ Compiler Reference](#)

### Extensions for GNU C++ compatibility

The following GNU C++ language extension is enabled with the **-qlanglvl=extended** option, which is the default language level used when you compile with the `xlC` and `xlc++` invocation commands.

Table 42. IBM XL C++ language extensions for compatibility with GNU C++

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>init_priority variable attribute</td>
<td>“The init_priority variable attribute (C++ only)” on page 104</td>
</tr>
</tbody>
</table>

The following GNU C++ language extension is enabled with the **-qlanglvl=extended** option, which is the default language level used when you compile with the `xlC` and `xlc++` invocation commands. It can also be enabled or disabled by a specific compiler option, listed in the table below.

Table 43. IBM XL C++ language extensions for compatibility with GNU C++, with individual option controls

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
<th>Individual option control</th>
</tr>
</thead>
<tbody>
<tr>
<td>Template instantiations declared as extern</td>
<td>“Template instantiation” on page 348</td>
<td>-qlanglvl=[no]gnuExternTemplate</td>
</tr>
</tbody>
</table>

Related information

- **-qlanglvl** in the [XL C/C++ Compiler Reference](#)
Extensions for VMX support

The VMX extensions are only accepted when all of the following conditions are met:
1. The `-qarch` option is set to a target architecture that supports VMX instructions.
2. The `-qenablevmx` option is in effect.
3. The `-qaltivec` option is in effect.

For more information on these options, see the XL C/C++ Compiler Reference.

Table 44. IBM XL C++ extensions to support the AltiVec Application Programming Interface specification

<table>
<thead>
<tr>
<th>Language feature</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Vector programming language extensions</td>
<td>&quot;Vector types” on page 53, Appendix C,</td>
</tr>
<tr>
<td></td>
<td>&quot;Vector data types and literals,” on page 415</td>
</tr>
</tbody>
</table>

The following features are IBM extensions to the AltiVec Application Programming Interface specification.

Table 45. IBM XL C++ extensions to the AltiVec Application Programming Interface specification

<table>
<thead>
<tr>
<th>Language extension</th>
<th>Discussed in:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Initializer lists for vector constants</td>
<td>&quot;Initialization of vectors” on page 93</td>
</tr>
<tr>
<td>typedef definitions for vector types</td>
<td>&quot;typedef definitions” on page 66</td>
</tr>
<tr>
<td>compound literals as initializers for static vector variables</td>
<td>&quot;Compound literals” on page 124</td>
</tr>
<tr>
<td>vector types as arguments to the <em>alignof</em> and typeof operators</td>
<td>&quot;The <em>alignof</em> operator” on page 136,</td>
</tr>
<tr>
<td></td>
<td>&quot;The typeof operator” on page 139</td>
</tr>
</tbody>
</table>
Appendix C. Vector data types and literals

This appendix presents the supported vector data types and literals. Note that these types are only recognized when all of the following conditions are met:

1. The `qarch` option is set to a target architecture that supports VMX instructions.
2. The `qenablevmx` option is in effect.
3. The `qaltivec` option is in effect.

For more information on these options, see the XL C/C++ Compiler Reference.

The following table lists the supported vector data types and the size and possible values for each type.

Table 46. Vector data types

<table>
<thead>
<tr>
<th>Type</th>
<th>Interpretation of content</th>
<th>Range of values</th>
</tr>
</thead>
<tbody>
<tr>
<td>vector unsigned char</td>
<td>16 unsigned char</td>
<td>0..255</td>
</tr>
<tr>
<td>vector signed char</td>
<td>16 signed char</td>
<td>-128..127</td>
</tr>
<tr>
<td>vector bool char</td>
<td>16 unsigned char</td>
<td>0, 255</td>
</tr>
<tr>
<td>vector unsigned short</td>
<td>8 unsigned short</td>
<td>0..65535</td>
</tr>
<tr>
<td>vector unsigned short int</td>
<td>8 signed short</td>
<td>-32768..32767</td>
</tr>
<tr>
<td>vector signed short int</td>
<td>8 unsigned short</td>
<td>0, 65535</td>
</tr>
<tr>
<td>vector bool short</td>
<td>8 unsigned short</td>
<td>0, 65535</td>
</tr>
<tr>
<td>vector bool short int</td>
<td>8 unsigned short</td>
<td>0, 65535</td>
</tr>
<tr>
<td>vector unsigned int</td>
<td>4 unsigned int</td>
<td>0..2^{32}-1</td>
</tr>
<tr>
<td>vector unsigned long</td>
<td>8 unsigned int</td>
<td>0..2^{64}-1</td>
</tr>
<tr>
<td>vector unsigned long int</td>
<td>8 unsigned int</td>
<td>0, 2^{64}-1</td>
</tr>
<tr>
<td>vector signed int</td>
<td>4 signed int</td>
<td>-2^{31}..2^{31}-1</td>
</tr>
<tr>
<td>vector signed long</td>
<td>8 signed int</td>
<td>-2^{63}..2^{63}-1</td>
</tr>
<tr>
<td>vector signed long int</td>
<td>8 signed int</td>
<td>0, 2^{64}-1</td>
</tr>
<tr>
<td>vector bool int</td>
<td>4 unsigned int</td>
<td>0, 2^{32}-1</td>
</tr>
<tr>
<td>vector bool long</td>
<td>8 unsigned int</td>
<td>0..2^{64}-1</td>
</tr>
<tr>
<td>vector bool long int</td>
<td>8 unsigned int</td>
<td>0, 2^{64}-1</td>
</tr>
<tr>
<td>vector float</td>
<td>4 float</td>
<td>IEEE-754 values</td>
</tr>
<tr>
<td>vector pixel</td>
<td>8 unsigned short</td>
<td>1/5/5/5/5 pixel</td>
</tr>
</tbody>
</table>

The compiler considers any long vector data type compatible with the corresponding `int` vector type.

Note: § The long vector types are deprecated.

The following table shows the supported vector literals and how the compiler interprets them to determine their values.
<table>
<thead>
<tr>
<th>Syntax</th>
<th>Interpreted by the compiler as</th>
</tr>
</thead>
<tbody>
<tr>
<td>(vector unsigned char)(single unsigned int value)</td>
<td>A set of 16 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector unsigned char)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 16 unsigned constants with a value specified by the 16 integer constant expressions.</td>
</tr>
<tr>
<td>(vector signed char)(single int value)</td>
<td>A set of 16 signed constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector signed char)(int value, ..., int value)</td>
<td>A set of 16 signed constants with a value specified by the 16 integer constant expressions.</td>
</tr>
<tr>
<td>(vector bool char)(single unsigned int)</td>
<td>A set of 16 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector bool char)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 16 unsigned constants with a value specified by the 16 integer constant expressions.</td>
</tr>
<tr>
<td>(vector unsigned short)(single unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector unsigned short)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the 8 integer constant expressions.</td>
</tr>
<tr>
<td>(vector signed short)(single int value)</td>
<td>A set of 8 signed constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector signed short)(int value, ..., int value)</td>
<td>A set of 8 signed constants with a value specified by the 8 integer constant expressions.</td>
</tr>
<tr>
<td>(vector bool short)(single unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector bool short)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the 8 integer constant expressions.</td>
</tr>
<tr>
<td>(vector unsigned int)(single unsigned int value)</td>
<td>A set of 4 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector unsigned int)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 4 unsigned constants with a value specified by the 4 integer constant expressions.</td>
</tr>
<tr>
<td>(vector signed int)(single int value)</td>
<td>A set of 4 signed constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector signed int)(int value, ..., int value)</td>
<td>A set of 4 signed constants with a value specified by the 4 integer constant expressions.</td>
</tr>
<tr>
<td>(vector bool int)(single unsigned int value)</td>
<td>A set of 4 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector bool int)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 4 unsigned constants with a value specified by the 4 integer constant expressions.</td>
</tr>
<tr>
<td>(vector float)(single float value)</td>
<td>A set of 4 floating-point constants with a value specified by the floating-point constant expression.</td>
</tr>
<tr>
<td>(vector float)(float value, ... float value)</td>
<td>A set of 4 floating-point constants with a value specified by the 4 floating-point constant expressions.</td>
</tr>
<tr>
<td>(vector pixel)(single unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the integer constant expression.</td>
</tr>
<tr>
<td>(vector pixel)(unsigned int value, ..., unsigned int value)</td>
<td>A set of 8 unsigned constants with a value specified by the 8 integer constant expressions.</td>
</tr>
</tbody>
</table>
Index

Special characters

<table>
<thead>
<tr>
<th>Character</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>align</strong></td>
<td>71</td>
</tr>
<tr>
<td><strong>cdecl</strong></td>
<td>219</td>
</tr>
<tr>
<td><strong>func</strong></td>
<td>16</td>
</tr>
<tr>
<td><strong>VA_ARGS</strong></td>
<td>385, 388</td>
</tr>
<tr>
<td><strong>Bool</strong></td>
<td>107, 108</td>
</tr>
<tr>
<td><strong>Pragma</strong></td>
<td>143</td>
</tr>
<tr>
<td>- (subtraction operator)</td>
<td>149</td>
</tr>
<tr>
<td>- (unary minus operator)</td>
<td>133</td>
</tr>
<tr>
<td>+= (increment operator)</td>
<td>132</td>
</tr>
<tr>
<td>&gt;&gt;= (arrow operator)</td>
<td>123</td>
</tr>
<tr>
<td>, (comma operator)</td>
<td>160</td>
</tr>
<tr>
<td>: (scope resolution operator)</td>
<td>120</td>
</tr>
<tr>
<td>!= (not equal to operator)</td>
<td>151</td>
</tr>
<tr>
<td>?: (conditional operators)</td>
<td>156</td>
</tr>
<tr>
<td>/ (division operator)</td>
<td>148</td>
</tr>
<tr>
<td>/= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td>. (dot operator)</td>
<td>123</td>
</tr>
<tr>
<td>$</td>
<td>16</td>
</tr>
<tr>
<td>* (indirection operator)</td>
<td>136</td>
</tr>
<tr>
<td>*+ (multiplication operator)</td>
<td>148</td>
</tr>
<tr>
<td>*= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td>\ continuation character</td>
<td>28, 383</td>
</tr>
<tr>
<td>\ escape character</td>
<td>31</td>
</tr>
<tr>
<td>[ ] (array subscript operator)</td>
<td>122</td>
</tr>
<tr>
<td>% (remainder)</td>
<td>148</td>
</tr>
<tr>
<td>&gt; (greater than operator)</td>
<td>150</td>
</tr>
<tr>
<td>&gt;&gt; (right-shift operator)</td>
<td>149</td>
</tr>
<tr>
<td>&gt;&gt;= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td>&gt;= (greater than or equal to operator)</td>
<td>150</td>
</tr>
<tr>
<td>&lt; (less than operator)</td>
<td>150</td>
</tr>
<tr>
<td>&lt;= (less than or equal to operator)</td>
<td>150</td>
</tr>
<tr>
<td>&amp;&amp; (logical AND operator)</td>
<td>154</td>
</tr>
<tr>
<td>&amp; (address operator)</td>
<td>134</td>
</tr>
<tr>
<td>&amp;&amp;= (bitwise inclusive OR operator)</td>
<td>153</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td>&amp;= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td># (preprocessor directive character)</td>
<td>383</td>
</tr>
<tr>
<td># (preprocessor operator)</td>
<td>389</td>
</tr>
<tr>
<td>#= (macro concatenation)</td>
<td>390</td>
</tr>
<tr>
<td>+= (addition operator)</td>
<td>148</td>
</tr>
<tr>
<td>+= (unary plus operator)</td>
<td>133</td>
</tr>
<tr>
<td>+= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td>= (assignment operator)</td>
<td>158</td>
</tr>
<tr>
<td>^= (bitwise exclusive OR operator)</td>
<td>153</td>
</tr>
<tr>
<td>^= (compound assignment operator)</td>
<td>159</td>
</tr>
<tr>
<td>~ (bitwise negation operator)</td>
<td>134</td>
</tr>
</tbody>
</table>

A

<table>
<thead>
<tr>
<th>Topic</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>aborting functions</td>
<td>379</td>
</tr>
<tr>
<td>abstract classes</td>
<td>298, 300</td>
</tr>
<tr>
<td>access rules</td>
<td>base classes</td>
</tr>
<tr>
<td></td>
<td>class types</td>
</tr>
<tr>
<td></td>
<td>friends</td>
</tr>
<tr>
<td></td>
<td>members</td>
</tr>
<tr>
<td></td>
<td>multiple access</td>
</tr>
<tr>
<td></td>
<td>protected members</td>
</tr>
<tr>
<td></td>
<td>virtual functions</td>
</tr>
<tr>
<td></td>
<td>access specifiers</td>
</tr>
<tr>
<td></td>
<td>in class derivations</td>
</tr>
<tr>
<td>accessibility</td>
<td>269, 290</td>
</tr>
<tr>
<td>addition operator</td>
<td>148</td>
</tr>
<tr>
<td>address operator</td>
<td>&amp;</td>
</tr>
<tr>
<td>GNU C extension</td>
<td>168</td>
</tr>
<tr>
<td>aggregate types</td>
<td>39, 306</td>
</tr>
<tr>
<td>initialization</td>
<td>94, 306</td>
</tr>
<tr>
<td>alias</td>
<td>88</td>
</tr>
<tr>
<td>function</td>
<td>17</td>
</tr>
<tr>
<td>type-based aliasing</td>
<td>84</td>
</tr>
<tr>
<td>alignment</td>
<td>102, 103</td>
</tr>
<tr>
<td>bit fields</td>
<td>59</td>
</tr>
<tr>
<td>structure members</td>
<td>56</td>
</tr>
<tr>
<td>structures</td>
<td>102</td>
</tr>
<tr>
<td>structures and unions</td>
<td>71</td>
</tr>
<tr>
<td>alignof operator</td>
<td>136</td>
</tr>
<tr>
<td>allocation</td>
<td>expressions</td>
</tr>
<tr>
<td>ambiguities</td>
<td>base and derived member names</td>
</tr>
<tr>
<td></td>
<td>base classes</td>
</tr>
<tr>
<td></td>
<td>resolving</td>
</tr>
<tr>
<td></td>
<td>virtual function calls</td>
</tr>
<tr>
<td>AND operator, bitwise</td>
<td>&amp;</td>
</tr>
<tr>
<td>AND operator, logical</td>
<td>&amp;</td>
</tr>
<tr>
<td>argc (argument count)</td>
<td>212</td>
</tr>
<tr>
<td>example</td>
<td>213</td>
</tr>
<tr>
<td>arguments</td>
<td>default</td>
</tr>
<tr>
<td>evaluation</td>
<td>218</td>
</tr>
<tr>
<td>macro</td>
<td>385</td>
</tr>
<tr>
<td>main function</td>
<td>212</td>
</tr>
<tr>
<td>of catch blocks</td>
<td>369</td>
</tr>
<tr>
<td>passing</td>
<td>191, 213</td>
</tr>
<tr>
<td>passing by reference</td>
<td>215</td>
</tr>
<tr>
<td>passing by value</td>
<td>214</td>
</tr>
<tr>
<td>trailing</td>
<td>385, 388</td>
</tr>
<tr>
<td>argy (argument vector)</td>
<td>212</td>
</tr>
<tr>
<td>example</td>
<td>213</td>
</tr>
<tr>
<td>arithmetic conversions</td>
<td>107</td>
</tr>
<tr>
<td>arithmetic types</td>
<td>type compatibility</td>
</tr>
<tr>
<td>arrays</td>
<td>array-to-pointer conversions</td>
</tr>
<tr>
<td>as function parameter</td>
<td>45, 205</td>
</tr>
<tr>
<td>declaration</td>
<td>45, 205, 256</td>
</tr>
<tr>
<td>arrays (continued)</td>
<td>description</td>
</tr>
<tr>
<td>flexible array member</td>
<td>56, 57</td>
</tr>
<tr>
<td>initialization</td>
<td>91</td>
</tr>
<tr>
<td>initializing</td>
<td>97</td>
</tr>
<tr>
<td>multidimensional</td>
<td>86</td>
</tr>
<tr>
<td>subscribing operator</td>
<td>122</td>
</tr>
<tr>
<td>type compatibility</td>
<td>88</td>
</tr>
<tr>
<td>variable length</td>
<td>82, 87</td>
</tr>
<tr>
<td>zero-extent</td>
<td>56, 58</td>
</tr>
<tr>
<td>ASCII character codes</td>
<td>32</td>
</tr>
<tr>
<td>asm</td>
<td>13</td>
</tr>
<tr>
<td>keyword</td>
<td>17, 48</td>
</tr>
<tr>
<td>labels</td>
<td>17</td>
</tr>
<tr>
<td>statements</td>
<td>186</td>
</tr>
<tr>
<td>assembly</td>
<td>labels</td>
</tr>
<tr>
<td>statements</td>
<td>186</td>
</tr>
<tr>
<td>assignment operator</td>
<td>=</td>
</tr>
<tr>
<td>compound</td>
<td>159</td>
</tr>
<tr>
<td>pointers</td>
<td>85</td>
</tr>
<tr>
<td>simple</td>
<td>158</td>
</tr>
<tr>
<td>associativity of operators</td>
<td>162</td>
</tr>
<tr>
<td>atoi function</td>
<td>378</td>
</tr>
<tr>
<td>auto storage class specifier</td>
<td>44</td>
</tr>
</tbody>
</table>

B

<table>
<thead>
<tr>
<th>Topic</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>base classes</td>
<td>abstract</td>
</tr>
<tr>
<td>access rules</td>
<td>283</td>
</tr>
<tr>
<td>ambiguities</td>
<td>289, 291</td>
</tr>
<tr>
<td>direct</td>
<td>288</td>
</tr>
<tr>
<td>indirect</td>
<td>278, 289</td>
</tr>
<tr>
<td>initialization</td>
<td>308</td>
</tr>
<tr>
<td>multiple access</td>
<td>290</td>
</tr>
<tr>
<td>pointers</td>
<td>281</td>
</tr>
<tr>
<td>virtual</td>
<td>289, 294</td>
</tr>
<tr>
<td>base list</td>
<td>289</td>
</tr>
<tr>
<td>best viable function</td>
<td>241</td>
</tr>
<tr>
<td>binary expressions and operators</td>
<td>147</td>
</tr>
<tr>
<td>binding</td>
<td>100</td>
</tr>
<tr>
<td>direct</td>
<td>101</td>
</tr>
<tr>
<td>dynamic</td>
<td>295</td>
</tr>
<tr>
<td>static</td>
<td>295</td>
</tr>
<tr>
<td>virtual functions</td>
<td>295</td>
</tr>
<tr>
<td>bit fields</td>
<td>58</td>
</tr>
<tr>
<td>as structure member</td>
<td>56</td>
</tr>
<tr>
<td>type name</td>
<td>139</td>
</tr>
<tr>
<td>bitwise negation operator</td>
<td>~</td>
</tr>
<tr>
<td>block statement</td>
<td>170</td>
</tr>
<tr>
<td>block visibility</td>
<td>3</td>
</tr>
<tr>
<td>bool</td>
<td>53</td>
</tr>
<tr>
<td>Boolean</td>
<td>conversions</td>
</tr>
<tr>
<td>data types</td>
<td>50</td>
</tr>
<tr>
<td>literals</td>
<td>21</td>
</tr>
<tr>
<td>break statement</td>
<td>181</td>
</tr>
<tr>
<td>built-in data types</td>
<td>39</td>
</tr>
</tbody>
</table>
compatibility (continued)
XL C++ and GCC 410, 412
compatible types
arithmetic types 53
arrays 88
in conditional expressions 156
vector types 54
complex types 52
candidate types 41
across source files 66
compound
assignment 159
expression 160
literal 121, 124
statement 170
types 39
casted goto 140, 168, 185
concatenation
macros 390
multibyte characters 29
u-literals, U-literals 54
conditional compilation directives 393
elf preprocessor directive 395
else preprocessor directive 396
endif preprocessor directive 396
elements 397
if preprocessor directive 395
ifdef preprocessor directive 395
ifndef preprocessor directive 396
conditional expressions (?) 156, 160
const 54, 69
casting away constness 216
function attribute 208
member functions 258
object 115
placement in type name 82
qualifier 68
vs. #define 384
cast, const 128, 216
constant expressions 62, 118
constant initializers 256
constructors 305
casting 322, 323
copy 325
conversion handling 373
conversion thrown in constructor 366
initialization
explicit 306
nontrivial 305, 313
overview 303
trivial 305, 313
continuation character 28, 383
continue statement 182
conversion
constructors 322
function 324
implicit conversion sequences 242
conversion sequence
ellipses 243
implicit 242
standard 242
user-defined 243
conversions
arithmetic 107
array-to-pointer 111
Boolean 107, 108
cast 144
explicit keyword 323
floating-point 108
function arguments 112
function-to-pointer 111
integral 108
long 110, 115, 242
pointer 110
pointer to derived class 294
pointer to member 261
qualification 112
references 112
standard 107
user-defined 320
void pointer 111
copy assignment operators 326
copy constructors 325
covariant virtual functions 297
CPLUSPLUS macro 392
cv-qualifier 68, 80
in parameter type specification 231
syntax 68

D
data members
description 256
scope 259
static 266
data types
aggregates 39
Boolean 50
built-in 39
character 52
compatible 41
complex 52
composite 41
compound 39
enumerated 62
floating 51
incomplete 40
integral 50
scalar 39
user-defined 39, 55
vector 53
void 53
DATE macro 390
deallocation
expressions 142
decimal integer literals 20
declaration 191
declarations
classes 246, 250
description 41
duplicate type qualifiers 68
friend specifier in member list 272
friends 276
pointers to members 260
resolving ambiguous statements 169
syntax 42, 81, 192
unsubscripted arrays 87
vector types 53
declarative region 2
declarators
description 79
reference 88
enumerations (continued)
   initialization 96
   trailing comma 63
enumerator 63
equal to operator (==) 151
error preprocessor directive 398
escape character \ 31
escape sequence 31
   alarm \a 32
   backslash \ \ 32
   backspace \b 32
   carriage return \r 32
double quotation mark " 32
form feed \f 32
horizontal tab \t 32
double quotation mark " 32
new-line \n 32
question mark \? 32
single quotation mark \' 32
vertical tab \v 32

exception handling 363
   argument matching 369
catch blocks 365
   arguments 369
constructors 373
destructors 373
   example, C++ 380
exception objects 363
function try blocks 363
   handlers 363, 365
order of catching 370
rethrowing exceptions 371
set_terminate 379
set_unexpected 379
special functions 377
stack unwinding 373
terminate function 378
throw expressions 364, 371
try blocks 363
try exceptions 366
unexpected function 377

exceptions
   declaration 365
function try block handlers 366
   specification 374
exclusive OR operator, bitwise (\^) 153
explicit
   instantiation, templates 350
keyword 322, 323
specializations, templates 351, 353
type conversions 144

exponent 22

expressions
   allocation 140
   assignment 158
   binary 147
cast 144
   comma 160
   conditional 156
deallocation 142
description 115
   integer constant 118
parenthesized 119
   pointer to member 155
   primary 117
   resolving ambiguous statements 169
statement 169
   throw 162, 371


expressions (continued)
   unary 131
external storage class specifier 9, 11, 46, 196
   with function pointers 219
   with template declaration 348, 351
   with variable length arrays 87

F
   file inclusion 392, 393
   FILE macro 391
   file scope data declarations
   unsynchronized arrays 87
   flexible array member 57
   float type specifier 51
   floating-point
   constant 23
   conversion 108
   literal 22
   promotion 109
   floating-point types 51
   for statement 179
   free store 315
   delete operator 142
   new operator 140
   friend
   access rules 276
   implicit conversion of pointers 284
   member functions 257
   nested classes 274
   relationships with classes when
   templates are involved 339
   scope 274
   specifier 272
   virtual functions 298

function
   aliases 17
   function attributes 206
   function designer 116
   function specificer
   explicit 322, 323
   function templates
   explicit specialization 354
   function try blocks 363
   handlers 366
   function-like macro 385
   functions 191
   arguments 191, 214
   conversions 112
   block 191
   body 191
   calling 213
   calls 121
   as I value 116
   class templates 338
   conversion function 324
declaration 191
   C++ 258
   examples 193
   multiple 195
   parameter names 205
   default arguments 216
evaluation 218
   restrictions 217
definition 191
   examples 194

E
   EBCDIC character codes 32
   elaborated type specifier 249
   else preprocessor directive 395
ellipsis
   conversion sequence 243
in function declaration 204
in function definition 204
in macro argument list 385
else
   preprocessor directive 396
statement 172
enclosing class 257, 274
endif preprocessor directive 396
entry point
   program 212
enum
   keyword 63
   enumerations 62
   compatibility 65, 66
   declaration 63


Index 419
functions (continued)
  exception handling 377
  exception specification 374
  friends 272
  function call operator 191
  function templates 340
  function-to-pointer conversions 111
  inline 197, 257
  library functions 191
  main 212
  name 191
    diagnostic 16
    overloading 229
  parameters 213, 214
  pointers to 219
  polymorphic 278
  predefined identifier 16
  prototype 191
  return statements 183
  return type 191, 201, 202
  return value 191, 202
  signature 203
  specifiable attributes 206
  specifiers 197
  template function
    template argument deduction 341
    type name 82
    virtual 258, 295, 298

G
  global register variables 48
  global variable 3, 9
  uninitialized 91
  goto statement 184
    computed goto 185
    restrictions 185
  greater than operator (>) 150
  greater than or equal to operator
    (>=) 150

H
  handlers 365
  hexadecimal
    floating constants 23
  hexadecimal integer literals 21
  hidden names 247, 249

I
  identifiers 15, 117
    case sensitivity 16
    id-expression 80, 118
    labels 167
    linkage 9
    namespace 6
    predefined 16
    reserved 13, 14, 16
    special characters 16, 31
  if
    preprocessor directive 395
    statement 172
  ifdef preprocessor directive 395
  ifndef preprocessor directive 396

K
  keywords 13
    exception handling 363
    language extension 14
    template 329, 359, 360
    underscore characters 14

L
  label
    as values 168
    implicit declaration 3
    in switch statement 174
    locally declared 168
    statement 167
  language extension viii, 14
    C99 403
    GNU C ix, 403
  left-shift operator (<) 149
  less than operator (<) 150
  less than or equal to operator (<=) 150
  LINE macro 391
  line preprocessor directive 398
  linkage 1, 8
    auto storage class specifier 44
    const cv-qualifier 69
    extern storage class specifier 11, 46
    external 9
    in function definition 196
    inline member functions 258
    internal 8, 44, 196
    language 10
    multiple function declarations 195
    none 9
    register storage class specifier 48
    specifications 10
    static storage class specifier 45
    weak symbols 105
    with function pointers 219
  linking to non-C++ programs 10
  literals 19
    Boolean 21
    character 26
    compound 121, 124
    floating-point 22
    integer 19, 20
      decimal 20
      hexadecimale 21
      octal 21
    string 27
    Unicode 33
  local
    classes 252
    type names 253
  logical operators
    ! (logical negation) 134
    || (logical OR) 154
    && (logical AND) 154
  long double type specifier 51
  long long type specifier 50, 54
  long type specifier 50, 54
  lvalues 68, 115, 117
    casting 144
    conversions 110, 115, 242

M
  macro
    definition 384
    typedef operator 139
    function-like 385
  invocation 385
  object-like 384
  variable argument 385, 388
promotions
  integral and floating-point 109
pseudo-destructors 314
punctuators 29
alternative representations 30
pure specifier 256, 258, 298, 300
pure virtual functions 300

Q
qualification conversions 112
qualified name 120, 251

R
references
  as return types 202
  binding 100
  conversions 112
  declarator 135
  description 88
  initialization 100
register storage class specifier 47
register variables 48
reinterpret_cast 127
remainder operator (%) 148
restrict 70
  in parameter type specification 231
return statement 183, 202
return type
  reference as 202
  size_t 137
right-shift operator (>>) 149
RTTI support 125, 130
rvalues 115

S
scalar types 39, 82
scope 1
  class 5
  class names 249
  description 2
  enclosing and nested 3
  friends 274
  function 3
  function prototype 3
  global 3
  global namespace 3
  identifiers 6
  local (block) 2
  local classes 252
  macro names 388
  member 259
  nested classes 250
  scope resolution operator
  ambiguous base classes 291
  description 120
  inheritance 280
  virtual functions 298
  sequence point 161
set_terminate function 379
set_unexpected function 377, 379
shift operators << and >> 149
short type specifier 50
side effect 70
signed type specifiers
  char 52
  int 50
  long 50
  long long 50
size_t 137
sized operator 137
  with variable length arrays 87
space character 383
special characters 31
special member functions 258
specifiers
  access control 283
  inline 197
  pure 258
  storage class 43
  splice preprocessor directive ## 390
  stack unwinding 373
standard type conversions 107
statement expression 171
statements 167
  block 170
  break 181
  compound 171
  continue 182
  do 178
  expressions 169
  for 179
  goto 184
  if 172
  labels 167
  null 186
  resolving ambiguities 169
  return 183, 202
  selection 172, 173
  switch 173
  while 177
static 54
  binding 295
  data members 266
  in array declaration 45, 205
  initialization of data members 268
  member functions 268
  members 251, 265
  storage class specifier 44, 196
  linkage 45
  with variable length arrays 87
static storage class specifier 9
static_cast 126
STDC macro 391
STDC_HOSTED macro 391
STDC_VERSION macro 391
storage class specifiers 43
  auto 44
  extern 46, 196
  mutable 47
  register 47
  static 44, 196
storage duration 1
  auto storage class specifier 44
  extern storage class specifier 46
  register storage class specifier 48
templates (continued)
  explicit specializations (continued)
    definition and declaration 353
    function templates 354
  function
    argument deduction 345
    overloadng 346
    partial ordering 347
    function templates 340
    type template argument deduction 344
  instantiation 329, 348, 351
  explicit 350
  forward declaration 348, 351
  implicit 348
  name binding 358
  parameters 330
    default arguments 331
    non-type 330
    template 331
    type 330
    partial specialization 356
    matching 358
    parameter and argument lists 357
    point of definition 359
    point of instantiation 359
    relationship between classes and their friends 339
  scope 353
  specialization 329, 348, 351
  temporary objects 320, 369
  tentative definition 43
  terminate function 363, 364, 370, 373,
  377, 378
  set terminate 379
  this pointer 261, 299
  throw expressions 162, 363, 371
  argument matching 369
  rethrowing exceptions 371
  within nested try blocks 364
  TIME macro 391
  tokens 13, 383
  alternative representations for operators and punctuators 30
  translation unit 1
  trigraph sequences 34
  truncation
    integer division 148
    try blocks 363
    nested 364
    try keyword 363
    type attributes 74
    type name 81
    local 253
    qualified 120, 251
    typename keyword 359
    typeof operator 139
  type qualifiers
    const 68, 69
    const and volatile 80
    duplicate 68
    in function parameters 231
    restrict 68, 70
    volatile 68
  type specifier
    class types 246
    elaborated 249
  type specifier (continued)
    overriding 104
    type specifiers 49
      _Bool 50
      (long) double 51
      bool 50
      char 52
      complex 52
      enumeration 62
      float 51
      int 50
      long 50
      long long 50
      short 50
      unsigned 50
      vector data types 53
      void 53
      wchar_t 50, 52
  typedef specifier 66
    class declaration 253
    local type names 253
    pointers to members 261
    qualified type name 251
    with variable length arrays 88
  typeid operator 125
  typename keyword 359
  typeof operator 139
  types
    class 246
    conversions 144
    type-based aliasing 84
    variably modified 86

U
  u-literal, U-literal 33
  unary expressions 131
  unary operators 131
    label value 140
    minus (−) 133
    plus (+) 133
  undeclared preprocessor directive 388
  underscore character 14, 16
  unexpected function 363, 377, 378
  set_unexpected 379
  Unicode 32
  unions 55
    as class type 245, 246
    cast to union type 146
    compatibility 65, 66
    designated initializer 91
    initialization 96
    specifier 56
    unnamed members 94
  universal character name 16, 26, 32
  unnamed namespaces 223
  unsubscribed arrays
    description 87, 205
  user-defined conversions 320
  user-defined data types 39, 55
  using declarations 226, 284, 292

Index 423
using declarations (continued)
   changing member access 287
   overloading member functions 285
using directive 226
UTF-16, UTF-32 33

V
variable
   in specified registers 48
variable attributes 101
variable length array 41, 87, 185
   as function parameter 88, 214, 241
   sizeof 118
   type name 82
variably modified types 86, 87, 175
vector data types 53
vector types 139, 415
   in typedef declarations 66
   literals 25, 415
virtual
   base classes 279, 289, 294
virtual functions 258, 295
   access 300
   ambiguous calls 298
   overriding 299
   pure specifier 300
visibility 1, 7
   block 3
   class members 271
VMX support ix, 407, 413
void 53
   in function definition 201, 204
   pointer 110, 111
volatile
   member functions 258
   qualifier 68, 70

W
warning preprocessor directive 398
wchar_t type specifier 26, 50, 52
weak symbol 105
while statement 177
white space 13, 35, 383, 389
wide characters
   literals 26
wide string literal 29

Z
zero-extent array 58
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